# The Fashoda Meet Theorem for Rectangles

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**Summary.** Here, the so called Fashoda Meet Theorem is proven in the case of rectangles. All cases of proper location of arcs are listed up, and it is shown that the theorem is valid in each case. Such a list of cases will be useful when one wants to apply the theorem.

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The articles [1], [6], [15], [17], [5], [2], [3], [16], [7], [14], [13], [10], [11], [8], [4], [9], and [12] provide the notation and terminology for this paper.

One can prove the following propositions:

- (1) For all real numbers a, b, d and for every point p of  $\mathcal{E}_{\mathrm{T}}^2$  such that a < b and  $p_2 = d$  and  $a \leq p_1$  and  $p_1 \leq b$  holds  $p \in \mathcal{L}([a, d], [b, d])$ .
- (2) Let *n* be a natural number, *P* be a subset of  $\mathcal{E}_{\mathrm{T}}^{n}$ , and  $p_{1}$ ,  $p_{2}$  be points of  $\mathcal{E}_{\mathrm{T}}^{n}$ . Suppose *P* is an arc from  $p_{1}$  to  $p_{2}$ . Then there exists a map *f* from I into  $\mathcal{E}_{\mathrm{T}}^{n}$  such that *f* is continuous and one-to-one and rng f = P and  $f(0) = p_{1}$  and  $f(1) = p_{2}$ .
- (3) Let  $p_1, p_2$  be points of  $\mathcal{E}^2_{\mathrm{T}}$  and b, c, d be real numbers. If  $(p_1)_1 < b$ and  $(p_1)_1 = (p_2)_1$  and  $c \leq (p_1)_2$  and  $(p_1)_2 < (p_2)_2$  and  $(p_2)_2 \leq d$ , then  $p_1 \leq_{\mathrm{Rectangle}((p_1)_1, b, c, d)} p_2$ .
- (4) Let  $p_1$ ,  $p_2$  be points of  $\mathcal{E}^2_{\mathrm{T}}$  and b, c be real numbers. Suppose  $(p_1)_1 < b$ and  $c < (p_2)_2$  and  $c \le (p_1)_2$  and  $(p_1)_2 \le (p_2)_2$  and  $(p_1)_1 \le (p_2)_1$  and  $(p_2)_1 \le b$ . Then  $p_1 \le_{\mathrm{Rectangle}((p_1)_1, b, c, (p_2)_2)} p_2$ .
- (5) Let  $p_1$ ,  $p_2$  be points of  $\mathcal{E}_{\mathrm{T}}^2$  and c, d be real numbers. Suppose  $(p_1)_1 < (p_2)_1$  and c < d and  $c \le (p_1)_2$  and  $(p_1)_2 \le d$  and  $c \le (p_2)_2$  and  $(p_2)_2 \le d$ . Then  $p_1 \le_{\mathrm{Rectangle}((p_1)_1, (p_2)_1, c, d)} p_2$ .

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- (6) Let  $p_1, p_2$  be points of  $\mathcal{E}_{\mathrm{T}}^2$  and b, d be real numbers. If  $(p_2)_2 < d$  and  $(p_2)_2 \leq (p_1)_2$  and  $(p_1)_2 \leq d$  and  $(p_1)_1 < (p_2)_1$  and  $(p_2)_1 \leq b$ , then  $p_1 \leq_{\mathrm{Rectangle}((p_1)_1, b, (p_2)_2, d)} p_2$ .
- (7) Let  $p_1, p_2$  be points of  $\mathcal{E}_{\mathrm{T}}^2$  and a, b, c, d be real numbers. Suppose a < band c < d and  $(p_1)_2 = d$  and  $(p_2)_2 = d$  and  $a \leq (p_1)_1$  and  $(p_1)_1 < (p_2)_1$ and  $(p_2)_1 \leq b$ . Then  $p_1 \leq_{\mathrm{Rectangle}(a,b,c,d)} p_2$ .
- (8) Let  $p_1, p_2$  be points of  $\mathcal{E}^2_{\mathrm{T}}$  and a, b, c, d be real numbers. Suppose a < band c < d and  $(p_1)_2 = d$  and  $(p_2)_1 = b$  and  $a \leq (p_1)_1$  and  $(p_1)_1 \leq b$  and  $c \leq (p_2)_2$  and  $(p_2)_2 \leq d$ . Then  $p_1 \leq_{\mathrm{Rectangle}(a,b,c,d)} p_2$ .
- (9) Let  $p_1, p_2$  be points of  $\mathcal{E}^2_{\mathrm{T}}$  and a, b, c, d be real numbers. Suppose a < band c < d and  $(p_1)_2 = d$  and  $(p_2)_2 = c$  and  $a \leq (p_1)_1$  and  $(p_1)_1 \leq b$  and  $a < (p_2)_1$  and  $(p_2)_1 \leq b$ . Then  $p_1 \leq_{\mathrm{Rectangle}(a,b,c,d)} p_2$ .
- (10) Let  $p_1$ ,  $p_2$  be points of  $\mathcal{E}_T^2$  and a, b, c, d be real numbers. Suppose a < band c < d and  $(p_1)_1 = b$  and  $(p_2)_1 = b$  and  $c \leq (p_2)_2$  and  $(p_2)_2 < (p_1)_2$ and  $(p_1)_2 \leq d$ . Then  $p_1 \leq_{\text{Rectangle}(a,b,c,d)} p_2$ .
- (11) Let  $p_1$ ,  $p_2$  be points of  $\mathcal{E}^2_T$  and a, b, c, d be real numbers. Suppose a < band c < d and  $(p_1)_1 = b$  and  $(p_2)_2 = c$  and  $c \le (p_1)_2$  and  $(p_1)_2 \le d$  and  $a < (p_2)_1$  and  $(p_2)_1 \le b$ . Then  $p_1 \le_{\text{Rectangle}(a,b,c,d)} p_2$ .
- (12) Let  $p_1$ ,  $p_2$  be points of  $\mathcal{E}^2_{\mathrm{T}}$  and a, b, c, d be real numbers. Suppose a < band c < d and  $(p_1)_2 = c$  and  $(p_2)_2 = c$  and  $a < (p_2)_1$  and  $(p_2)_1 < (p_1)_1$ and  $(p_1)_1 \leq b$ . Then  $p_1 \leq_{\mathrm{Rectangle}(a,b,c,d)} p_2$ .
- (13) Let  $p_1$ ,  $p_2$  be points of  $\mathcal{E}_{\mathrm{T}}^2$  and a, b, c, d be real numbers. Suppose a < band c < d and  $(p_1)_2 = d$  and  $(p_2)_1 = b$  and  $a \leq (p_1)_1$  and  $(p_1)_1 \leq b$  and  $c \leq (p_2)_2$  and  $(p_2)_2 \leq d$ . Then  $p_1 \leq_{\mathrm{Rectangle}(a,b,c,d)} p_2$ .
- (14) Let  $p_1, p_2, p_3, p_4$  be points of  $\mathcal{E}_T^2$  and a, b, c, d be real numbers. Suppose a < b and c < d and  $(p_1)_1 = a$  and  $(p_2)_1 = a$  and  $(p_3)_1 = a$  and  $(p_4)_1 = a$  and  $c \le (p_1)_2$  and  $(p_1)_2 < (p_2)_2$  and  $(p_2)_2 < (p_3)_2$  and  $(p_3)_2 < (p_4)_2$  and  $(p_4)_2 \le d$ . Then  $p_1, p_2, p_3, p_4$  are in this order on Rectangle(a, b, c, d).
- (15) Let  $p_1, p_2, p_3, p_4$  be points of  $\mathcal{E}_T^2$  and a, b, c, d be real numbers. Suppose a < b and c < d and  $(p_1)_1 = a$  and  $(p_2)_1 = a$  and  $(p_3)_1 = a$  and  $(p_4)_2 = d$  and  $c \leq (p_1)_2$  and  $(p_1)_2 < (p_2)_2$  and  $(p_2)_2 < (p_3)_2$  and  $(p_3)_2 \leq d$  and  $a \leq (p_4)_1$  and  $(p_4)_1 \leq b$ . Then  $p_1, p_2, p_3, p_4$  are in this order on Rectangle(a, b, c, d).
- (16) Let  $p_1$ ,  $p_2$ ,  $p_3$ ,  $p_4$  be points of  $\mathcal{E}_{\mathrm{T}}^2$  and a, b, c, d be real numbers. Suppose a < b and c < d and  $(p_1)_1 = a$  and  $(p_2)_1 = a$  and  $(p_3)_1 = a$  and  $(p_4)_1 = b$  and  $c \leq (p_1)_2$  and  $(p_1)_2 < (p_2)_2$  and  $(p_2)_2 < (p_3)_2$  and  $(p_3)_2 \leq d$  and  $c \leq (p_4)_2$  and  $(p_4)_2 \leq d$ . Then  $p_1$ ,  $p_2$ ,  $p_3$ ,  $p_4$  are in this order on Rectangle(a, b, c, d).
- (17) Let  $p_1, p_2, p_3, p_4$  be points of  $\mathcal{E}_T^2$  and a, b, c, d be real numbers. Suppose a < b and c < d and  $(p_1)_1 = a$  and  $(p_2)_1 = a$  and  $(p_3)_1 = a$  and

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 $(p_4)_2 = c \text{ and } c \leq (p_1)_2 \text{ and } (p_1)_2 < (p_2)_2 \text{ and } (p_2)_2 < (p_3)_2 \text{ and } (p_3)_2 \leq d$ and  $a < (p_4)_1$  and  $(p_4)_1 \leq b$ . Then  $p_1, p_2, p_3, p_4$  are in this order on Rectangle(a, b, c, d).

- (18) Let  $p_1, p_2, p_3, p_4$  be points of  $\mathcal{E}_T^2$  and a, b, c, d be real numbers. Suppose a < b and c < d and  $(p_1)_1 = a$  and  $(p_2)_1 = a$  and  $(p_3)_2 = d$  and  $(p_4)_2 = d$  and  $c \leq (p_1)_2$  and  $(p_1)_2 < (p_2)_2$  and  $(p_2)_2 \leq d$  and  $a \leq (p_3)_1$  and  $(p_3)_1 < (p_4)_1$  and  $(p_4)_1 \leq b$ . Then  $p_1, p_2, p_3, p_4$  are in this order on Rectangle(a, b, c, d).
- (19) Let  $p_1, p_2, p_3, p_4$  be points of  $\mathcal{E}_T^2$  and a, b, c, d be real numbers. Suppose a < b and c < d and  $(p_1)_1 = a$  and  $(p_2)_1 = a$  and  $(p_3)_2 = d$  and  $(p_4)_1 = b$  and  $c \leq (p_1)_2$  and  $(p_1)_2 < (p_2)_2$  and  $(p_2)_2 \leq d$  and  $a \leq (p_3)_1$  and  $(p_3)_1 \leq b$  and  $c \leq (p_4)_2$  and  $(p_4)_2 \leq d$ . Then  $p_1, p_2, p_3, p_4$  are in this order on Rectangle(a, b, c, d).
- (20) Let  $p_1$ ,  $p_2$ ,  $p_3$ ,  $p_4$  be points of  $\mathcal{E}_T^2$  and a, b, c, d be real numbers. Suppose a < b and c < d and  $(p_1)_1 = a$  and  $(p_2)_1 = a$  and  $(p_3)_2 = d$  and  $(p_4)_2 = c$  and  $c \leq (p_1)_2$  and  $(p_1)_2 < (p_2)_2$  and  $(p_2)_2 \leq d$  and  $a \leq (p_3)_1$  and  $(p_3)_1 \leq b$  and  $a < (p_4)_1$  and  $(p_4)_1 \leq b$ . Then  $p_1$ ,  $p_2$ ,  $p_3$ ,  $p_4$  are in this order on Rectangle(a, b, c, d).
- (21) Let  $p_1, p_2, p_3, p_4$  be points of  $\mathcal{E}_T^2$  and a, b, c, d be real numbers. Suppose a < b and c < d and  $(p_1)_1 = a$  and  $(p_2)_1 = a$  and  $(p_3)_1 = b$  and  $(p_4)_1 = b$  and  $c \leq (p_1)_2$  and  $(p_1)_2 < (p_2)_2$  and  $(p_2)_2 \leq d$  and  $c \leq (p_4)_2$  and  $(p_4)_2 < (p_3)_2$  and  $(p_3)_2 \leq d$ . Then  $p_1, p_2, p_3, p_4$  are in this order on Rectangle(a, b, c, d).
- (22) Let  $p_1, p_2, p_3, p_4$  be points of  $\mathcal{E}_T^2$  and a, b, c, d be real numbers. Suppose a < b and c < d and  $(p_1)_1 = a$  and  $(p_2)_1 = a$  and  $(p_3)_1 = b$  and  $(p_4)_2 = c$  and  $c \leq (p_1)_2$  and  $(p_1)_2 < (p_2)_2$  and  $(p_2)_2 \leq d$  and  $c \leq (p_3)_2$  and  $(p_3)_2 \leq d$  and  $a < (p_4)_1$  and  $(p_4)_1 \leq b$ . Then  $p_1, p_2, p_3, p_4$  are in this order on Rectangle(a, b, c, d).
- (23) Let  $p_1, p_2, p_3, p_4$  be points of  $\mathcal{E}_T^2$  and a, b, c, d be real numbers. Suppose a < b and c < d and  $(p_1)_1 = a$  and  $(p_2)_1 = a$  and  $(p_3)_2 = c$  and  $(p_4)_2 = c$  and  $c \leq (p_1)_2$  and  $(p_1)_2 < (p_2)_2$  and  $(p_2)_2 \leq d$  and  $a < (p_4)_1$  and  $(p_4)_1 < (p_3)_1$  and  $(p_3)_1 \leq b$ . Then  $p_1, p_2, p_3, p_4$  are in this order on Rectangle(a, b, c, d).
- (24) Let  $p_1, p_2, p_3, p_4$  be points of  $\mathcal{E}_T^2$  and a, b, c, d be real numbers. Suppose a < b and c < d and  $(p_1)_1 = a$  and  $(p_2)_2 = d$  and  $(p_3)_2 = d$  and  $(p_4)_2 = d$  and  $c \leq (p_1)_2$  and  $(p_1)_2 \leq d$  and  $a \leq (p_2)_1$  and  $(p_2)_1 < (p_3)_1$  and  $(p_3)_1 < (p_4)_1$  and  $(p_4)_1 \leq b$ . Then  $p_1, p_2, p_3, p_4$  are in this order on Rectangle(a, b, c, d).
- (25) Let  $p_1, p_2, p_3, p_4$  be points of  $\mathcal{E}_{\mathrm{T}}^2$  and a, b, c, d be real numbers. Suppose a < b and c < d and  $(p_1)_{\mathbf{1}} = a$  and  $(p_2)_{\mathbf{2}} = d$  and  $(p_3)_{\mathbf{2}} = d$  and  $(p_4)_{\mathbf{1}} = b$  and  $c \leq (p_1)_{\mathbf{2}}$  and  $(p_1)_{\mathbf{2}} \leq d$  and  $a \leq (p_2)_{\mathbf{1}}$  and  $(p_2)_{\mathbf{1}} < (p_3)_{\mathbf{1}}$  and

 $(p_3)_1 \leq b$  and  $c \leq (p_4)_2$  and  $(p_4)_2 \leq d$ . Then  $p_1, p_2, p_3, p_4$  are in this order on Rectangle(a, b, c, d).

- (26) Let  $p_1, p_2, p_3, p_4$  be points of  $\mathcal{E}_T^2$  and a, b, c, d be real numbers. Suppose a < b and c < d and  $(p_1)_1 = a$  and  $(p_2)_2 = d$  and  $(p_3)_2 = d$  and  $(p_4)_2 = c$  and  $c \leq (p_1)_2$  and  $(p_1)_2 \leq d$  and  $a \leq (p_2)_1$  and  $(p_2)_1 < (p_3)_1$  and  $(p_3)_1 \leq b$  and  $a < (p_4)_1$  and  $(p_4)_1 \leq b$ . Then  $p_1, p_2, p_3, p_4$  are in this order on Rectangle(a, b, c, d).
- (27) Let  $p_1, p_2, p_3, p_4$  be points of  $\mathcal{E}_T^2$  and a, b, c, d be real numbers. Suppose a < b and c < d and  $(p_1)_1 = a$  and  $(p_2)_2 = d$  and  $(p_3)_1 = b$  and  $(p_4)_1 = b$  and  $c \leq (p_1)_2$  and  $(p_1)_2 \leq d$  and  $a \leq (p_2)_1$  and  $(p_2)_1 \leq b$  and  $c \leq (p_4)_2$  and  $(p_4)_2 < (p_3)_2$  and  $(p_3)_2 \leq d$ . Then  $p_1, p_2, p_3, p_4$  are in this order on Rectangle(a, b, c, d).
- (28) Let  $p_1$ ,  $p_2$ ,  $p_3$ ,  $p_4$  be points of  $\mathcal{E}_{\mathrm{T}}^2$ . Suppose  $(p_1)_1 \neq (p_3)_1$  and  $(p_4)_2 \neq (p_2)_2$  and  $(p_4)_2 \leq (p_1)_2$  and  $(p_1)_2 \leq (p_2)_2$  and  $(p_1)_1 \leq (p_2)_1$  and  $(p_2)_1 \leq (p_3)_1$  and  $(p_4)_2 \leq (p_3)_2$  and  $(p_3)_2 \leq (p_2)_2$  and  $(p_1)_1 < (p_4)_1$  and  $(p_4)_1 \leq (p_3)_1$ . Then  $p_1$ ,  $p_2$ ,  $p_3$ ,  $p_4$  are in this order on Rectangle $((p_1)_1, (p_3)_1, (p_4)_2, (p_2)_2)$ .
- (29) Let  $p_1, p_2, p_3, p_4$  be points of  $\mathcal{E}_T^2$  and a, b, c, d be real numbers. Suppose a < b and c < d and  $(p_1)_1 = a$  and  $(p_2)_2 = d$  and  $(p_3)_2 = c$  and  $(p_4)_2 = c$  and  $c \leq (p_1)_2$  and  $(p_1)_2 \leq d$  and  $a \leq (p_2)_1$  and  $(p_2)_1 \leq b$  and  $a < (p_4)_1$  and  $(p_4)_1 < (p_3)_1$  and  $(p_3)_1 \leq b$ . Then  $p_1, p_2, p_3, p_4$  are in this order on Rectangle(a, b, c, d).
- (30) Let  $p_1, p_2, p_3, p_4$  be points of  $\mathcal{E}_{\mathrm{T}}^2$  and a, b, c, d be real numbers. Suppose a < b and c < d and  $(p_1)_1 = a$  and  $(p_2)_1 = b$  and  $(p_3)_1 = b$  and  $(p_4)_1 = b$  and  $c \leq (p_1)_2$  and  $(p_1)_2 \leq d$  and  $d \geq (p_2)_2$  and  $(p_2)_2 > (p_3)_2$  and  $(p_3)_2 > (p_4)_2$  and  $(p_4)_2 \geq c$ . Then  $p_1, p_2, p_3, p_4$  are in this order on Rectangle(a, b, c, d).
- (31) Let  $p_1$ ,  $p_2$ ,  $p_3$ ,  $p_4$  be points of  $\mathcal{E}_{\mathrm{T}}^2$  and a, b, c, d be real numbers. Suppose a < b and c < d and  $(p_1)_{\mathbf{1}} = a$  and  $(p_2)_{\mathbf{1}} = b$  and  $(p_3)_{\mathbf{1}} = b$  and  $(p_4)_{\mathbf{2}} = c$  and  $c \leq (p_1)_{\mathbf{2}}$  and  $(p_1)_{\mathbf{2}} \leq d$  and  $d \geq (p_2)_{\mathbf{2}}$  and  $(p_2)_{\mathbf{2}} > (p_3)_{\mathbf{2}}$  and  $(p_3)_{\mathbf{2}} \geq c$  and  $a < (p_4)_{\mathbf{1}}$  and  $(p_4)_{\mathbf{1}} \leq b$ . Then  $p_1$ ,  $p_2$ ,  $p_3$ ,  $p_4$  are in this order on Rectangle(a, b, c, d).
- (32) Let  $p_1$ ,  $p_2$ ,  $p_3$ ,  $p_4$  be points of  $\mathcal{E}_T^2$  and a, b, c, d be real numbers. Suppose a < b and c < d and  $(p_1)_1 = a$  and  $(p_2)_1 = b$  and  $(p_3)_2 = c$  and  $(p_4)_2 = c$  and  $c \leq (p_1)_2$  and  $(p_1)_2 \leq d$  and  $c \leq (p_2)_2$  and  $(p_2)_2 \leq d$  and  $a < (p_4)_1$  and  $(p_4)_1 < (p_3)_1$  and  $(p_3)_1 \leq b$ . Then  $p_1$ ,  $p_2$ ,  $p_3$ ,  $p_4$  are in this order on Rectangle(a, b, c, d).
- (33) Let  $p_1, p_2, p_3, p_4$  be points of  $\mathcal{E}_T^2$  and a, b, c, d be real numbers. Suppose a < b and c < d and  $(p_1)_1 = a$  and  $(p_2)_2 = c$  and  $(p_3)_2 = c$  and  $(p_4)_2 = c$  and  $c \leq (p_1)_2$  and  $(p_1)_2 \leq d$  and  $a < (p_4)_1$  and  $(p_4)_1 < (p_3)_1$  and  $(p_3)_1 < (p_2)_1$  and  $(p_2)_1 \leq b$ . Then  $p_1, p_2, p_3, p_4$  are in this order on

 $\operatorname{Rectangle}(a, b, c, d).$ 

- (34) Let  $p_1$ ,  $p_2$ ,  $p_3$ ,  $p_4$  be points of  $\mathcal{E}_T^2$  and a, b, c, d be real numbers. Suppose a < b and c < d and  $(p_1)_2 = d$  and  $(p_2)_2 = d$  and  $(p_3)_2 = d$  and  $(p_4)_2 = d$  and  $a \le (p_1)_1$  and  $(p_1)_1 < (p_2)_1$  and  $(p_2)_1 < (p_3)_1$  and  $(p_3)_1 < (p_4)_1$  and  $(p_4)_1 \le b$ . Then  $p_1$ ,  $p_2$ ,  $p_3$ ,  $p_4$  are in this order on Rectangle(a, b, c, d).
- (35) Let  $p_1, p_2, p_3, p_4$  be points of  $\mathcal{E}_{\mathrm{T}}^2$  and a, b, c, d be real numbers. Suppose a < b and c < d and  $(p_1)_2 = d$  and  $(p_2)_2 = d$  and  $(p_3)_2 = d$  and  $(p_4)_1 = b$  and  $a \leq (p_1)_1$  and  $(p_1)_1 < (p_2)_1$  and  $(p_2)_1 < (p_3)_1$  and  $(p_3)_1 \leq b$  and  $c \leq (p_4)_2$  and  $(p_4)_2 \leq d$ . Then  $p_1, p_2, p_3, p_4$  are in this order on Rectangle(a, b, c, d).
- (36) Let  $p_1, p_2, p_3, p_4$  be points of  $\mathcal{E}_{\mathrm{T}}^2$  and a, b, c, d be real numbers. Suppose a < b and c < d and  $(p_1)_2 = d$  and  $(p_2)_2 = d$  and  $(p_3)_2 = d$  and  $(p_4)_2 = c$  and  $a \leq (p_1)_1$  and  $(p_1)_1 < (p_2)_1$  and  $(p_2)_1 < (p_3)_1$  and  $(p_3)_1 \leq b$  and  $a < (p_4)_1$  and  $(p_4)_1 \leq b$ . Then  $p_1, p_2, p_3, p_4$  are in this order on Rectangle(a, b, c, d).
- (37) Let  $p_1, p_2, p_3, p_4$  be points of  $\mathcal{E}_{\mathrm{T}}^2$  and a, b, c, d be real numbers. Suppose a < b and c < d and  $(p_1)_2 = d$  and  $(p_2)_2 = d$  and  $(p_3)_1 = b$  and  $(p_4)_1 = b$  and  $a \leq (p_1)_1$  and  $(p_1)_1 < (p_2)_1$  and  $(p_2)_1 \leq b$  and  $c \leq (p_4)_2$  and  $(p_4)_2 < (p_3)_2$  and  $(p_3)_2 \leq d$ . Then  $p_1, p_2, p_3, p_4$  are in this order on Rectangle(a, b, c, d).
- (38) Let  $p_1$ ,  $p_2$ ,  $p_3$ ,  $p_4$  be points of  $\mathcal{E}_T^2$  and a, b, c, d be real numbers. Suppose a < b and c < d and  $(p_1)_2 = d$  and  $(p_2)_2 = d$  and  $(p_3)_1 = b$  and  $(p_4)_2 = c$  and  $a \leq (p_1)_1$  and  $(p_1)_1 < (p_2)_1$  and  $(p_2)_1 \leq b$  and  $c \leq (p_3)_2$  and  $(p_3)_2 \leq d$  and  $a < (p_4)_1$  and  $(p_4)_1 \leq b$ . Then  $p_1$ ,  $p_2$ ,  $p_3$ ,  $p_4$  are in this order on Rectangle(a, b, c, d).
- (39) Let  $p_1, p_2, p_3, p_4$  be points of  $\mathcal{E}_T^2$  and a, b, c, d be real numbers. Suppose a < b and c < d and  $(p_1)_2 = d$  and  $(p_2)_2 = d$  and  $(p_3)_2 = c$  and  $(p_4)_2 = c$  and  $a \leq (p_1)_1$  and  $(p_1)_1 < (p_2)_1$  and  $(p_2)_1 \leq b$  and  $a < (p_4)_1$  and  $(p_4)_1 < (p_3)_1$  and  $(p_3)_1 \leq b$ . Then  $p_1, p_2, p_3, p_4$  are in this order on Rectangle(a, b, c, d).
- (40) Let  $p_1$ ,  $p_2$ ,  $p_3$ ,  $p_4$  be points of  $\mathcal{E}_T^2$  and a, b, c, d be real numbers. Suppose a < b and c < d and  $(p_1)_2 = d$  and  $(p_2)_1 = b$  and  $(p_3)_1 = b$  and  $(p_4)_1 = b$  and  $a \leq (p_1)_1$  and  $(p_1)_1 \leq b$  and  $d \geq (p_2)_2$  and  $(p_2)_2 > (p_3)_2$  and  $(p_3)_2 > (p_4)_2$  and  $(p_4)_2 \geq c$ . Then  $p_1$ ,  $p_2$ ,  $p_3$ ,  $p_4$  are in this order on Rectangle(a, b, c, d).
- (41) Let  $p_1, p_2, p_3, p_4$  be points of  $\mathcal{E}_{\mathrm{T}}^2$  and a, b, c, d be real numbers. Suppose a < b and c < d and  $(p_1)_2 = d$  and  $(p_2)_1 = b$  and  $(p_3)_1 = b$  and  $(p_4)_2 = c$  and  $a \leq (p_1)_1$  and  $(p_1)_1 \leq b$  and  $d \geq (p_2)_2$  and  $(p_2)_2 > (p_3)_2$  and  $(p_3)_2 \geq c$  and  $a < (p_4)_1$  and  $(p_4)_1 \leq b$ . Then  $p_1, p_2, p_3, p_4$  are in this order on Rectangle(a, b, c, d).

- (42) Let  $p_1$ ,  $p_2$ ,  $p_3$ ,  $p_4$  be points of  $\mathcal{E}_T^2$  and a, b, c, d be real numbers. Suppose a < b and c < d and  $(p_1)_2 = d$  and  $(p_2)_1 = b$  and  $(p_3)_2 = c$  and  $(p_4)_2 = c$  and  $a \le (p_1)_1$  and  $(p_1)_1 \le b$  and  $c \le (p_2)_2$  and  $(p_2)_2 \le d$  and  $a < (p_4)_1$  and  $(p_4)_1 < (p_3)_1$  and  $(p_3)_1 \le b$ . Then  $p_1$ ,  $p_2$ ,  $p_3$ ,  $p_4$  are in this order on Rectangle(a, b, c, d).
- (43) Let  $p_1, p_2, p_3, p_4$  be points of  $\mathcal{E}_T^2$  and a, b, c, d be real numbers. Suppose a < b and c < d and  $(p_1)_2 = d$  and  $(p_2)_2 = c$  and  $(p_3)_2 = c$  and  $(p_4)_2 = c$  and  $a \leq (p_1)_1$  and  $(p_1)_1 \leq b$  and  $a < (p_4)_1$  and  $(p_4)_1 < (p_3)_1$  and  $(p_3)_1 < (p_2)_1$  and  $(p_2)_1 \leq b$ . Then  $p_1, p_2, p_3, p_4$  are in this order on Rectangle(a, b, c, d).
- (44) Let  $p_1$ ,  $p_2$ ,  $p_3$ ,  $p_4$  be points of  $\mathcal{E}_T^2$  and a, b, c, d be real numbers. Suppose a < b and c < d and  $(p_1)_1 = b$  and  $(p_2)_1 = b$  and  $(p_3)_1 = b$  and  $(p_4)_1 = b$  and  $d \ge (p_1)_2$  and  $(p_1)_2 > (p_2)_2$  and  $(p_2)_2 > (p_3)_2$  and  $(p_3)_2 > (p_4)_2$  and  $(p_4)_2 \ge c$ . Then  $p_1$ ,  $p_2$ ,  $p_3$ ,  $p_4$  are in this order on Rectangle(a, b, c, d).
- (45) Let  $p_1, p_2, p_3, p_4$  be points of  $\mathcal{E}_T^2$  and a, b, c, d be real numbers. Suppose a < b and c < d and  $(p_1)_1 = b$  and  $(p_2)_1 = b$  and  $(p_3)_1 = b$  and  $(p_4)_2 = c$  and  $d \ge (p_1)_2$  and  $(p_1)_2 > (p_2)_2$  and  $(p_2)_2 > (p_3)_2$  and  $(p_3)_2 \ge c$  and  $a < (p_4)_1$  and  $(p_4)_1 \le b$ . Then  $p_1, p_2, p_3, p_4$  are in this order on Rectangle(a, b, c, d).
- (46) Let  $p_1, p_2, p_3, p_4$  be points of  $\mathcal{E}_T^2$  and a, b, c, d be real numbers. Suppose a < b and c < d and  $(p_1)_1 = b$  and  $(p_2)_1 = b$  and  $(p_3)_2 = c$  and  $(p_4)_2 = c$  and  $d \ge (p_1)_2$  and  $(p_1)_2 > (p_2)_2$  and  $(p_2)_2 \ge c$  and  $b \ge (p_3)_1$  and  $(p_3)_1 > (p_4)_1$  and  $(p_4)_1 > a$ . Then  $p_1, p_2, p_3, p_4$  are in this order on Rectangle(a, b, c, d).
- (47) Let  $p_1, p_2, p_3, p_4$  be points of  $\mathcal{E}_T^2$  and a, b, c, d be real numbers. Suppose a < b and c < d and  $(p_1)_1 = b$  and  $(p_2)_2 = c$  and  $(p_3)_2 = c$  and  $(p_4)_2 = c$  and  $c \leq (p_1)_2$  and  $(p_1)_2 \leq d$  and  $b \geq (p_2)_1$  and  $(p_2)_1 > (p_3)_1$  and  $(p_3)_1 > (p_4)_1$  and  $(p_4)_1 > a$ . Then  $p_1, p_2, p_3, p_4$  are in this order on Rectangle(a, b, c, d).
- (48) Let  $p_1, p_2, p_3, p_4$  be points of  $\mathcal{E}_T^2$  and a, b, c, d be real numbers. Suppose a < b and c < d and  $(p_1)_2 = c$  and  $(p_2)_2 = c$  and  $(p_3)_2 = c$  and  $(p_4)_2 = c$  and  $b \ge (p_1)_1$  and  $(p_1)_1 > (p_2)_1$  and  $(p_2)_1 > (p_3)_1$  and  $(p_3)_1 > (p_4)_1$  and  $(p_4)_1 > a$ . Then  $p_1, p_2, p_3, p_4$  are in this order on Rectangle(a, b, c, d).
- (49) Let A, B, C, D be real numbers and h, g be maps from  $\mathcal{E}_{T}^{2}$  into  $\mathcal{E}_{T}^{2}$ . Suppose A > 0 and C > 0 and  $h = \operatorname{AffineMap}(A, B, C, D)$  and  $g = \operatorname{AffineMap}(\frac{1}{A}, -\frac{B}{A}, \frac{1}{C}, -\frac{D}{C})$ . Then  $g = h^{-1}$  and  $h = g^{-1}$ .
- (50) Let A, B, C, D be real numbers and h be a map from  $\mathcal{E}_{\mathrm{T}}^2$  into  $\mathcal{E}_{\mathrm{T}}^2$ . Suppose A > 0 and C > 0 and  $h = \mathrm{AffineMap}(A, B, C, D)$ . Then h is a homeomorphism and for all points  $p_1, p_2$  of  $\mathcal{E}_{\mathrm{T}}^2$  such that  $(p_1)_1 < (p_2)_1$  holds  $h(p_1)_1 < h(p_2)_1$ .

- (51) Let A, B, C, D be real numbers and h be a map from  $\mathcal{E}_{\mathrm{T}}^2$  into  $\mathcal{E}_{\mathrm{T}}^2$ . Suppose A > 0 and C > 0 and  $h = \operatorname{AffineMap}(A, B, C, D)$ . Then h is a homeomorphism and for all points  $p_1$ ,  $p_2$  of  $\mathcal{E}_{\mathrm{T}}^2$  such that  $(p_1)_2 < (p_2)_2$  holds  $h(p_1)_2 < h(p_2)_2$ .
- (52) Let a, b, c, d be real numbers, h be a map from  $\mathcal{E}_{\mathrm{T}}^2$ into  $\mathcal{E}_{\mathrm{T}}^2$ , and f be a map from  $\mathbb{I}$  into  $\mathcal{E}_{\mathrm{T}}^2$ . Suppose a < b and c < d and  $h = \operatorname{AffineMap}(\frac{2}{b-a}, -\frac{b+a}{b-a}, \frac{2}{d-c}, -\frac{d+c}{d-c})$  and rng  $f \subseteq \operatorname{ClosedInsideOfRectangle}(a, b, c, d)$ . Then  $\operatorname{rng}(h \cdot f) \subseteq \operatorname{ClosedInsideOfRectangle}(-1, 1, -1, 1)$ .
- (53) Let a, b, c, d be real numbers, h be a map from  $\mathcal{E}_{\mathrm{T}}^2$  into  $\mathcal{E}_{\mathrm{T}}^2$ , and f be a map from  $\mathbb{I}$  into  $\mathcal{E}_{\mathrm{T}}^2$ . Suppose a < b and c < d and  $h = \mathrm{AffineMap}(\frac{2}{b-a}, -\frac{b+a}{b-a}, \frac{2}{d-c}, -\frac{d+c}{d-c})$  and f is continuous and one-to-one. Then  $h \cdot f$  is continuous and one-to-one.
- (54) Let a, b, c, d be real numbers, h be a map from  $\mathcal{E}_{\mathrm{T}}^2$  into  $\mathcal{E}_{\mathrm{T}}^2$ , f be a map from  $\mathbb{I}$  into  $\mathcal{E}_{\mathrm{T}}^2$ , and O be a point of  $\mathbb{I}$ . Suppose a < b and c < d and  $h = \operatorname{AffineMap}(\frac{2}{b-a}, -\frac{b+a}{b-a}, \frac{2}{d-c}, -\frac{d+c}{d-c})$  and  $f(O)_{\mathbf{1}} = a$ . Then  $(h \cdot f)(O)_{\mathbf{1}} = -1$ .
- (55) Let a, b, c, d be real numbers, h be a map from  $\mathcal{E}_{\mathrm{T}}^2$  into  $\mathcal{E}_{\mathrm{T}}^2$ , f be a map from  $\mathbb{I}$  into  $\mathcal{E}_{\mathrm{T}}^2$ , and I be a point of  $\mathbb{I}$ . Suppose a < b and c < d and  $h = \mathrm{AffineMap}(\frac{2}{b-a}, -\frac{b+a}{b-a}, \frac{2}{d-c}, -\frac{d+c}{d-c})$  and  $f(I)_2 = d$ . Then  $(h \cdot f)(I)_2 = 1$ .
- (56) Let a, b, c, d be real numbers, h be a map from  $\mathcal{E}_{\mathrm{T}}^2$  into  $\mathcal{E}_{\mathrm{T}}^2$ , f be a map from  $\mathbb{I}$  into  $\mathcal{E}_{\mathrm{T}}^2$ , and I be a point of  $\mathbb{I}$ . Suppose a < b and c < d and  $h = \mathrm{AffineMap}(\frac{2}{b-a}, -\frac{b+a}{b-a}, \frac{2}{d-c}, -\frac{d+c}{d-c})$  and  $f(I)_{\mathbf{1}} = b$ . Then  $(h \cdot f)(I)_{\mathbf{1}} = 1$ .
- (57) Let a, b, c, d be real numbers, h be a map from  $\mathcal{E}_{\mathrm{T}}^2$  into  $\mathcal{E}_{\mathrm{T}}^2$ , f be a map from  $\mathbb{I}$  into  $\mathcal{E}_{\mathrm{T}}^2$ , and I be a point of  $\mathbb{I}$ . Suppose a < b and c < d and  $h = \mathrm{AffineMap}(\frac{2}{b-a}, -\frac{b+a}{b-a}, \frac{2}{d-c}, -\frac{d+c}{d-c})$  and  $f(I)_2 = c$ . Then  $(h \cdot f)(I)_2 = -1$ .
- (58) Let a, b, c, d be real numbers, h be a map from  $\mathcal{E}_{\mathrm{T}}^2$  into  $\mathcal{E}_{\mathrm{T}}^2$ , f be a map from  $\mathbb{I}$  into  $\mathcal{E}_{\mathrm{T}}^2$ , and O, I be points of  $\mathbb{I}$ . Suppose a < b and c < d and  $h = \mathrm{AffineMap}(\frac{2}{b-a}, -\frac{b+a}{b-a}, \frac{2}{d-c}, -\frac{d+c}{d-c})$  and  $c \leq f(O)_2$  and  $f(O)_2 < f(I)_2$  and  $f(I)_2 \leq d$ . Then  $-1 \leq (h \cdot f)(O)_2$  and  $(h \cdot f)(O)_2 < (h \cdot f)(I)_2$  and  $(h \cdot f)(I)_2 \leq 1$ .
- (59) Let a, b, c, d be real numbers, h be a map from  $\mathcal{E}_{\mathrm{T}}^2$  into  $\mathcal{E}_{\mathrm{T}}^2$ , f be a map from  $\mathbb{I}$  into  $\mathcal{E}_{\mathrm{T}}^2$ , and O, I be points of  $\mathbb{I}$ . Suppose a < b and c < d and  $h = \mathrm{AffineMap}(\frac{2}{b-a}, -\frac{b+a}{b-a}, \frac{2}{d-c}, -\frac{d+c}{d-c})$  and  $c \leq f(O)_2$  and  $f(O)_2 \leq d$  and  $a \leq f(I)_1$  and  $f(I)_1 \leq b$ . Then  $-1 \leq (h \cdot f)(O)_2$  and  $(h \cdot f)(O)_2 \leq 1$  and  $-1 \leq (h \cdot f)(I)_1$  and  $(h \cdot f)(I)_1 \leq 1$ .
- (60) Let a, b, c, d be real numbers, h be a map from  $\mathcal{E}_{\mathrm{T}}^2$  into  $\mathcal{E}_{\mathrm{T}}^2$ , f be a map from  $\mathbb{I}$  into  $\mathcal{E}_{\mathrm{T}}^2$ , and O, I be points of  $\mathbb{I}$ . Suppose a < b and c < d and  $h = \operatorname{AffineMap}(\frac{2}{b-a}, -\frac{b+a}{b-a}, \frac{2}{d-c}, -\frac{d+c}{d-c})$  and  $c \leq f(O)_2$  and  $f(O)_2 \leq d$  and  $c \leq f(I)_2$  and  $f(I)_2 \leq d$ . Then  $-1 \leq (h \cdot f)(O)_2$  and  $(h \cdot f)(O)_2 \leq 1$  and  $-1 \leq (h \cdot f)(I)_2$  and  $(h \cdot f)(I)_2 \leq 1$ .

- (61) Let a, b, c, d be real numbers, h be a map from  $\mathcal{E}_{\mathrm{T}}^2$  into  $\mathcal{E}_{\mathrm{T}}^2$ , f be a map from  $\mathbb{I}$  into  $\mathcal{E}_{\mathrm{T}}^2$ , and O, I be points of  $\mathbb{I}$ . Suppose a < b and c < d and  $h = \mathrm{AffineMap}(\frac{2}{b-a}, -\frac{b+a}{b-a}, \frac{2}{d-c}, -\frac{d+c}{d-c})$  and  $c \leq f(O)_2$  and  $f(O)_2 \leq d$  and  $a < f(I)_1$  and  $f(I)_1 \leq b$ . Then  $-1 \leq (h \cdot f)(O)_2$  and  $(h \cdot f)(O)_2 \leq 1$  and  $-1 < (h \cdot f)(I)_1$  and  $(h \cdot f)(I)_1 \leq 1$ .
- (62) Let a, b, c, d be real numbers, h be a map from  $\mathcal{E}_{\mathrm{T}}^2$  into  $\mathcal{E}_{\mathrm{T}}^2$ , f be a map from  $\mathbb{I}$  into  $\mathcal{E}_{\mathrm{T}}^2$ , and O, I be points of  $\mathbb{I}$ . Suppose a < b and c < d and  $h = \operatorname{AffineMap}(\frac{2}{b-a}, -\frac{b+a}{b-a}, \frac{2}{d-c}, -\frac{d+c}{d-c})$  and  $a \leq f(O)_{\mathbf{1}}$  and  $f(O)_{\mathbf{1}} < f(I)_{\mathbf{1}}$  and  $f(I)_{\mathbf{1}} \leq b$ . Then  $-1 \leq (h \cdot f)(O)_{\mathbf{1}}$  and  $(h \cdot f)(O)_{\mathbf{1}} < (h \cdot f)(I)_{\mathbf{1}}$  and  $(h \cdot f)(I)_{\mathbf{1}} \leq 1$ .
- (63) Let a, b, c, d be real numbers, h be a map from  $\mathcal{E}_{\mathrm{T}}^2$  into  $\mathcal{E}_{\mathrm{T}}^2$ , f be a map from  $\mathbb{I}$  into  $\mathcal{E}_{\mathrm{T}}^2$ , and O, I be points of  $\mathbb{I}$ . Suppose a < b and c < d and  $h = \mathrm{AffineMap}(\frac{2}{b-a}, -\frac{b+a}{b-a}, \frac{2}{d-c}, -\frac{d+c}{d-c})$  and  $a \leq f(O)_{\mathbf{1}}$  and  $f(O)_{\mathbf{1}} \leq b$  and  $c \leq f(I)_{\mathbf{2}}$  and  $f(I)_{\mathbf{2}} \leq d$ . Then  $-1 \leq (h \cdot f)(O)_{\mathbf{1}}$  and  $(h \cdot f)(O)_{\mathbf{1}} \leq 1$  and  $-1 \leq (h \cdot f)(I)_{\mathbf{2}}$  and  $(h \cdot f)(I)_{\mathbf{2}} \leq 1$ .
- (64) Let a, b, c, d be real numbers, h be a map from  $\mathcal{E}_{\mathrm{T}}^2$  into  $\mathcal{E}_{\mathrm{T}}^2$ , f be a map from  $\mathbb{I}$  into  $\mathcal{E}_{\mathrm{T}}^2$ , and O, I be points of  $\mathbb{I}$ . Suppose a < b and c < d and  $h = \operatorname{AffineMap}(\frac{2}{b-a}, -\frac{b+a}{b-a}, \frac{2}{d-c}, -\frac{d+c}{d-c})$  and  $a \leq f(O)_{\mathbf{1}}$  and  $f(O)_{\mathbf{1}} \leq b$  and  $a < f(I)_{\mathbf{1}}$  and  $f(I)_{\mathbf{1}} \leq b$ . Then  $-1 \leq (h \cdot f)(O)_{\mathbf{1}}$  and  $(h \cdot f)(O)_{\mathbf{1}} \leq 1$  and  $-1 < (h \cdot f)(I)_{\mathbf{1}}$  and  $(h \cdot f)(I)_{\mathbf{1}} \leq 1$ .
- (65) Let a, b, c, d be real numbers, h be a map from  $\mathcal{E}_{\mathrm{T}}^2$  into  $\mathcal{E}_{\mathrm{T}}^2$ , f be a map from  $\mathbb{I}$  into  $\mathcal{E}_{\mathrm{T}}^2$ , and O, I be points of  $\mathbb{I}$ . Suppose a < b and c < d and  $h = \mathrm{AffineMap}(\frac{2}{b-a}, -\frac{b+a}{b-a}, \frac{2}{d-c}, -\frac{d+c}{d-c})$  and  $d \ge f(O)_2$  and  $f(O)_2 > f(I)_2$  and  $f(I)_2 \ge c$ . Then  $1 \ge (h \cdot f)(O)_2$  and  $(h \cdot f)(O)_2 > (h \cdot f)(I)_2$  and  $(h \cdot f)(I)_2 \ge -1$ .
- (66) Let a, b, c, d be real numbers, h be a map from  $\mathcal{E}_{\mathrm{T}}^2$  into  $\mathcal{E}_{\mathrm{T}}^2$ , f be a map from  $\mathbb{I}$  into  $\mathcal{E}_{\mathrm{T}}^2$ , and O, I be points of  $\mathbb{I}$ . Suppose a < b and c < d and  $h = \operatorname{AffineMap}(\frac{2}{b-a}, -\frac{b+a}{b-a}, \frac{2}{d-c}, -\frac{d+c}{d-c})$  and  $c \leq f(O)_2$  and  $f(O)_2 \leq d$  and  $a < f(I)_1$  and  $f(I)_1 \leq b$ . Then  $-1 \leq (h \cdot f)(O)_2$  and  $(h \cdot f)(O)_2 \leq 1$  and  $-1 < (h \cdot f)(I)_1$  and  $(h \cdot f)(I)_1 \leq 1$ .
- (67) Let a, b, c, d be real numbers, h be a map from  $\mathcal{E}_{\mathrm{T}}^2$  into  $\mathcal{E}_{\mathrm{T}}^2$ , f be a map from  $\mathbb{I}$  into  $\mathcal{E}_{\mathrm{T}}^2$ , and O, I be points of  $\mathbb{I}$ . Suppose a < b and c < d and  $h = \operatorname{AffineMap}(\frac{2}{b-a}, -\frac{b+a}{b-a}, \frac{2}{d-c}, -\frac{d+c}{d-c})$  and  $a < f(I)_1$  and  $f(I)_1 < f(O)_1$  and  $f(O)_1 \leq b$ . Then  $-1 < (h \cdot f)(I)_1$  and  $(h \cdot f)(I)_1 < (h \cdot f)(O)_1$  and  $(h \cdot f)(O)_1 \leq 1$ .

One can prove the following propositions:

(68) Let  $p_1$ ,  $p_2$ ,  $p_3$ ,  $p_4$  be points of  $\mathcal{E}_{\mathrm{T}}^2$ , a, b, c, d be real numbers, and f, g be maps from  $\mathbb{I}$  into  $\mathcal{E}_{\mathrm{T}}^2$ . Suppose that a < b and c < d and  $(p_1)_1 = a$  and  $(p_2)_1 = a$  and  $(p_3)_1 = a$  and  $(p_4)_1 = a$  and  $c \leq (p_1)_2$  and  $(p_1)_2 < (p_2)_2$  and  $(p_2)_2 < (p_3)_2$  and  $(p_3)_2 < (p_4)_2$  and  $(p_4)_2 \leq d$  and  $f(0) = p_1$  and  $f(1) = p_3$  and  $g(0) = p_2$  and  $g(1) = p_4$  and f is continuous and one-to-one and g is

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continuous and one-to-one and rng  $f \subseteq$  ClosedInsideOfRectangle(a, b, c, d) and rng  $g \subseteq$  ClosedInsideOfRectangle(a, b, c, d). Then rng f meets rng g.

- (69) Let  $p_1$ ,  $p_2$ ,  $p_3$ ,  $p_4$  be points of  $\mathcal{E}^2_{\mathrm{T}}$ , a, b, c, d be real numbers, and P, Q be subsets of  $\mathcal{E}^2_{\mathrm{T}}$ . Suppose that a < b and c < d and  $(p_1)_1 = a$  and  $(p_2)_1 = a$  and  $(p_3)_1 = a$  and  $(p_4)_1 = a$  and  $c \leq (p_1)_2$  and  $(p_1)_2 < (p_2)_2$  and  $(p_2)_2 < (p_3)_2$  and  $(p_3)_2 < (p_4)_2$  and  $(p_4)_2 \leq d$  and P is an arc from  $p_1$  to  $p_3$ and Q is an arc from  $p_2$  to  $p_4$  and  $P \subseteq$  ClosedInsideOfRectangle(a, b, c, d)and  $Q \subseteq$  ClosedInsideOfRectangle(a, b, c, d). Then P meets Q.
- (70) Let  $p_1$ ,  $p_2$ ,  $p_3$ ,  $p_4$  be points of  $\mathcal{E}_{\mathrm{T}}^2$ , a, b, c, d be real numbers, and f, g be maps from  $\mathbb{I}$  into  $\mathcal{E}_{\mathrm{T}}^2$ . Suppose that a < b and c < d and  $(p_1)_1 = a$  and  $(p_2)_1 = a$  and  $(p_3)_1 = a$  and  $(p_4)_2 = d$  and  $c \leq (p_1)_2$  and  $(p_1)_2 < (p_2)_2$  and  $(p_2)_2 < (p_3)_2$  and  $(p_3)_2 \leq d$  and  $a \leq (p_4)_1$  and  $(p_4)_1 \leq b$  and  $f(0) = p_1$  and  $f(1) = p_3$  and  $g(0) = p_2$  and  $g(1) = p_4$  and f is continuous and one-to-one and g is continuous and one-to-one and  $\operatorname{rng} f \subseteq \operatorname{ClosedInsideOfRectangle}(a, b, c, d)$ . Then  $\operatorname{rng} f$  meets  $\operatorname{rng} g$ .
- (71) Let  $p_1$ ,  $p_2$ ,  $p_3$ ,  $p_4$  be points of  $\mathcal{E}_{\mathrm{T}}^2$ , a, b, c, d be real numbers, and P, Q be subsets of  $\mathcal{E}_{\mathrm{T}}^2$ . Suppose that a < b and c < d and  $(p_1)_1 = a$ and  $(p_2)_1 = a$  and  $(p_3)_1 = a$  and  $(p_4)_2 = d$  and  $c \leq (p_1)_2$  and  $(p_1)_2 < (p_2)_2$  and  $(p_2)_2 < (p_3)_2$  and  $(p_3)_2 \leq d$  and  $a \leq (p_4)_1$ and  $(p_4)_1 \leq b$  and P is an arc from  $p_1$  to  $p_3$  and Q is an arc from  $p_2$  to  $p_4$  and  $P \subseteq$  ClosedInsideOfRectangle(a, b, c, d) and  $Q \subseteq$ ClosedInsideOfRectangle(a, b, c, d). Then P meets Q.
- (72) Let  $p_1$ ,  $p_2$ ,  $p_3$ ,  $p_4$  be points of  $\mathcal{E}_{\mathrm{T}}^2$ , a, b, c, d be real numbers, and f, g be maps from I into  $\mathcal{E}_{\mathrm{T}}^2$ . Suppose that a < b and c < d and  $(p_1)_1 = a$  and  $(p_2)_1 = a$  and  $(p_3)_1 = a$  and  $(p_4)_1 = b$  and  $c \leq (p_1)_2$  and  $(p_1)_2 < (p_2)_2$  and  $(p_2)_2 < (p_3)_2$  and  $(p_3)_2 \leq d$  and  $c \leq (p_4)_2$  and  $(p_4)_2 \leq d$  and  $f(0) = p_1$  and  $f(1) = p_3$  and  $g(0) = p_2$  and  $g(1) = p_4$  and f is continuous and one-to-one and g is continuous and one-to-one and  $\operatorname{rng} f \subseteq \operatorname{ClosedInsideOfRectangle}(a, b, c, d)$ . Then  $\operatorname{rng} f$  meets  $\operatorname{rng} g$ .
- (73) Let  $p_1, p_2, p_3, p_4$  be points of  $\mathcal{E}^2_{\mathrm{T}}$ , a, b, c, d be real numbers, and P, Q be subsets of  $\mathcal{E}^2_{\mathrm{T}}$ . Suppose that a < b and c < d and  $(p_1)_1 = a$ and  $(p_2)_1 = a$  and  $(p_3)_1 = a$  and  $(p_4)_1 = b$  and  $c \leq (p_1)_2$  and  $(p_1)_2 < (p_2)_2$  and  $(p_2)_2 < (p_3)_2$  and  $(p_3)_2 \leq d$  and  $c \leq (p_4)_2$ and  $(p_4)_2 \leq d$  and P is an arc from  $p_1$  to  $p_3$  and Q is an arc from  $p_2$  to  $p_4$  and  $P \subseteq$  ClosedInsideOfRectangle(a, b, c, d) and  $Q \subseteq$ ClosedInsideOfRectangle(a, b, c, d). Then P meets Q.
- (74) Let  $p_1$ ,  $p_2$ ,  $p_3$ ,  $p_4$  be points of  $\mathcal{E}_{\mathrm{T}}^2$ , a, b, c, d be real numbers, and f, g be maps from  $\mathbb{I}$  into  $\mathcal{E}_{\mathrm{T}}^2$ . Suppose that a < b and c < d and  $(p_1)_1 = a$  and  $(p_2)_1 = a$  and  $(p_3)_1 = a$  and  $(p_4)_2 = c$  and  $c \leq (p_1)_2$

and  $(p_1)_2 < (p_2)_2$  and  $(p_2)_2 < (p_3)_2$  and  $(p_3)_2 \leq d$  and  $a < (p_4)_1$ and  $(p_4)_1 \leq b$  and  $f(0) = p_1$  and  $f(1) = p_3$  and  $g(0) = p_2$  and  $g(1) = p_4$  and f is continuous and one-to-one and g is continuous and one-to-one and rng  $f \subseteq$  ClosedInsideOfRectangle(a, b, c, d) and rng  $g \subseteq$ ClosedInsideOfRectangle(a, b, c, d). Then rng f meets rng g.

- (75) Let  $p_1, p_2, p_3, p_4$  be points of  $\mathcal{E}^2_{\mathrm{T}}$ , a, b, c, d be real numbers, and P, Q be subsets of  $\mathcal{E}^2_{\mathrm{T}}$ . Suppose that a < b and c < d and  $(p_1)_1 = a$ and  $(p_2)_1 = a$  and  $(p_3)_1 = a$  and  $(p_4)_2 = c$  and  $c \leq (p_1)_2$  and  $(p_1)_2 < (p_2)_2$  and  $(p_2)_2 < (p_3)_2$  and  $(p_3)_2 \leq d$  and  $a < (p_4)_1$ and  $(p_4)_1 \leq b$  and P is an arc from  $p_1$  to  $p_3$  and Q is an arc from  $p_2$  to  $p_4$  and  $P \subseteq$  ClosedInsideOfRectangle(a, b, c, d) and  $Q \subseteq$ ClosedInsideOfRectangle(a, b, c, d). Then P meets Q.
- (76) Let  $p_1, p_2, p_3, p_4$  be points of  $\mathcal{E}_{\mathrm{T}}^2$ , a, b, c, d be real numbers, and f, g be maps from  $\mathbb{I}$  into  $\mathcal{E}_{\mathrm{T}}^2$ . Suppose that a < b and c < d and  $(p_1)_1 = a$  and  $(p_2)_1 = a$  and  $(p_3)_2 = d$  and  $(p_4)_2 = d$  and  $c \leq (p_1)_2$  and  $(p_1)_2 < (p_2)_2$  and  $(p_2)_2 \leq d$  and  $a \leq (p_3)_1$  and  $(p_3)_1 < (p_4)_1$  and  $(p_4)_1 \leq b$  and  $f(0) = p_1$  and  $f(1) = p_3$  and  $g(0) = p_2$  and  $g(1) = p_4$  and f is continuous and one-to-one and g is continuous and one-to-one and  $\operatorname{rng} f \subseteq \operatorname{ClosedInsideOfRectangle}(a, b, c, d)$ . Then  $\operatorname{rng} f$  meets  $\operatorname{rng} g$ .
- (77) Let  $p_1, p_2, p_3, p_4$  be points of  $\mathcal{E}_{\mathrm{T}}^2$ , a, b, c, d be real numbers, and P, Q be subsets of  $\mathcal{E}_{\mathrm{T}}^2$ . Suppose that a < b and c < d and  $(p_1)_1 = a$  and  $(p_2)_1 = a$ and  $(p_3)_2 = d$  and  $(p_4)_2 = d$  and  $c \leq (p_1)_2$  and  $(p_1)_2 < (p_2)_2$  and  $(p_2)_2 \leq d$ and  $a \leq (p_3)_1$  and  $(p_3)_1 < (p_4)_1$  and  $(p_4)_1 \leq b$  and P is an arc from  $p_1$  to  $p_3$ and Q is an arc from  $p_2$  to  $p_4$  and  $P \subseteq$  ClosedInsideOfRectangle(a, b, c, d)and  $Q \subseteq$  ClosedInsideOfRectangle(a, b, c, d). Then P meets Q.
- (78) Let  $p_1, p_2, p_3, p_4$  be points of  $\mathcal{E}_{\mathrm{T}}^2$ , a, b, c, d be real numbers, and f, g be maps from  $\mathbb{I}$  into  $\mathcal{E}_{\mathrm{T}}^2$ . Suppose that a < b and c < d and  $(p_1)_1 = a$  and  $(p_2)_1 = a$  and  $(p_3)_2 = d$  and  $(p_4)_1 = b$  and  $c \leq (p_1)_2$  and  $(p_1)_2 < (p_2)_2$  and  $(p_2)_2 \leq d$  and  $a \leq (p_3)_1$  and  $(p_3)_1 \leq b$  and  $c \leq (p_4)_2$  and  $(p_4)_2 \leq d$  and  $f(0) = p_1$  and  $f(1) = p_3$  and  $g(0) = p_2$  and  $g(1) = p_4$  and f is continuous and one-to-one and g is continuous and one-to-one and  $\operatorname{rng} f \subseteq \operatorname{ClosedInsideOfRectangle}(a, b, c, d)$  and  $\operatorname{rng} g \subseteq \operatorname{ClosedInsideOfRectangle}(a, b, c, d)$ . Then  $\operatorname{rng} f$  meets  $\operatorname{rng} g$ .
- (79) Let  $p_1$ ,  $p_2$ ,  $p_3$ ,  $p_4$  be points of  $\mathcal{E}_{\mathrm{T}}^2$ , a, b, c, d be real numbers, and P, Q be subsets of  $\mathcal{E}_{\mathrm{T}}^2$ . Suppose that a < b and c < d and  $(p_1)_1 = a$ and  $(p_2)_1 = a$  and  $(p_3)_2 = d$  and  $(p_4)_1 = b$  and  $c \leq (p_1)_2$  and  $(p_1)_2 < (p_2)_2$  and  $(p_2)_2 \leq d$  and  $a \leq (p_3)_1$  and  $(p_3)_1 \leq b$  and  $c \leq (p_4)_2$  and  $(p_4)_2 \leq d$  and P is an arc from  $p_1$  to  $p_3$  and Q is an arc from  $p_2$  to  $p_4$  and  $P \subseteq$  ClosedInsideOfRectangle(a, b, c, d) and  $Q \subseteq$  ClosedInsideOfRectangle(a, b, c, d). Then P meets Q.

- (80) Let  $p_1$ ,  $p_2$ ,  $p_3$ ,  $p_4$  be points of  $\mathcal{E}_{\mathrm{T}}^2$ , a, b, c, d be real numbers, and f, g be maps from  $\mathbb{I}$  into  $\mathcal{E}_{\mathrm{T}}^2$ . Suppose that a < b and c < d and  $(p_1)_1 = a$  and  $(p_2)_1 = a$  and  $(p_3)_2 = d$  and  $(p_4)_2 = c$  and  $c \leq (p_1)_2$  and  $(p_1)_2 < (p_2)_2$  and  $(p_2)_2 \leq d$  and  $a \leq (p_3)_1$  and  $(p_3)_1 \leq b$  and  $a < (p_4)_1$  and  $(p_4)_1 \leq b$  and  $f(0) = p_1$  and  $f(1) = p_3$  and  $g(0) = p_2$  and  $g(1) = p_4$  and f is continuous and one-to-one and g is continuous and one-to-one and  $\operatorname{rng} f \subseteq \operatorname{ClosedInsideOfRectangle}(a, b, c, d)$ . Then  $\operatorname{rng} f$  meets  $\operatorname{rng} g$ .
- (81) Let  $p_1$ ,  $p_2$ ,  $p_3$ ,  $p_4$  be points of  $\mathcal{E}_{\mathrm{T}}^2$ , a, b, c, d be real numbers, and P, Q be subsets of  $\mathcal{E}_{\mathrm{T}}^2$ . Suppose that a < b and c < d and  $(p_1)_1 = a$ and  $(p_2)_1 = a$  and  $(p_3)_2 = d$  and  $(p_4)_2 = c$  and  $c \leq (p_1)_2$  and  $(p_1)_2 < (p_2)_2$  and  $(p_2)_2 \leq d$  and  $a \leq (p_3)_1$  and  $(p_3)_1 \leq b$  and  $a < (p_4)_1$  and  $(p_4)_1 \leq b$  and P is an arc from  $p_1$  to  $p_3$  and Q is an arc from  $p_2$  to  $p_4$  and  $P \subseteq$  ClosedInsideOfRectangle(a, b, c, d) and  $Q \subseteq$  ClosedInsideOfRectangle(a, b, c, d). Then P meets Q.
- (82) Let  $p_1$ ,  $p_2$ ,  $p_3$ ,  $p_4$  be points of  $\mathcal{E}_{\mathrm{T}}^2$ , a, b, c, d be real numbers, and f, g be maps from I into  $\mathcal{E}_{\mathrm{T}}^2$ . Suppose that a < b and c < d and  $(p_1)_1 = a$  and  $(p_2)_1 = a$  and  $(p_3)_1 = b$  and  $(p_4)_1 = b$  and  $c \leq (p_1)_2$  and  $(p_1)_2 < (p_2)_2$  and  $(p_2)_2 \leq d$  and  $c \leq (p_4)_2$  and  $(p_4)_2 < (p_3)_2$  and  $(p_3)_2 \leq d$  and  $f(0) = p_1$  and  $f(1) = p_3$  and  $g(0) = p_2$  and  $g(1) = p_4$  and f is continuous and one-to-one and g is continuous and one-to-one and  $\operatorname{rng} f \subseteq \operatorname{ClosedInsideOfRectangle}(a, b, c, d)$ . Then  $\operatorname{rng} f$  meets  $\operatorname{rng} g$ .
- (83) Let  $p_1, p_2, p_3, p_4$  be points of  $\mathcal{E}_{\mathrm{T}}^2$ , a, b, c, d be real numbers, and P, Q be subsets of  $\mathcal{E}_{\mathrm{T}}^2$ . Suppose that a < b and c < d and  $(p_1)_1 = a$  and  $(p_2)_1 = a$ and  $(p_3)_1 = b$  and  $(p_4)_1 = b$  and  $c \leq (p_1)_2$  and  $(p_1)_2 < (p_2)_2$  and  $(p_2)_2 \leq d$ and  $c \leq (p_4)_2$  and  $(p_4)_2 < (p_3)_2$  and  $(p_3)_2 \leq d$  and P is an arc from  $p_1$  to  $p_3$ and Q is an arc from  $p_2$  to  $p_4$  and  $P \subseteq$  ClosedInsideOfRectangle(a, b, c, d)and  $Q \subseteq$  ClosedInsideOfRectangle(a, b, c, d). Then P meets Q.
- (84) Let  $p_1$ ,  $p_2$ ,  $p_3$ ,  $p_4$  be points of  $\mathcal{E}_{\mathrm{T}}^2$ , a, b, c, d be real numbers, and f, g be maps from  $\mathbb{I}$  into  $\mathcal{E}_{\mathrm{T}}^2$ . Suppose that a < b and c < d and  $(p_1)_1 = a$  and  $(p_2)_1 = a$  and  $(p_3)_1 = b$  and  $(p_4)_2 = c$  and  $c \leq (p_1)_2$  and  $(p_1)_2 < (p_2)_2$  and  $(p_2)_2 \leq d$  and  $c \leq (p_3)_2$  and  $(p_3)_2 \leq d$  and  $a < (p_4)_1$  and  $(p_4)_1 \leq b$  and  $f(0) = p_1$  and  $f(1) = p_3$  and  $g(0) = p_2$  and  $g(1) = p_4$  and f is continuous and one-to-one and  $\operatorname{rg} f \subseteq \operatorname{ClosedInsideOfRectangle}(a, b, c, d)$  and  $\operatorname{rng} g \subseteq \operatorname{ClosedInsideOfRectangle}(a, b, c, d)$ . Then  $\operatorname{rng} f$  meets  $\operatorname{rng} g$ .
- (85) Let  $p_1$ ,  $p_2$ ,  $p_3$ ,  $p_4$  be points of  $\mathcal{E}^2_{\mathrm{T}}$ , a, b, c, d be real numbers, and P, Q be subsets of  $\mathcal{E}^2_{\mathrm{T}}$ . Suppose that a < b and c < d and  $(p_1)_1 = a$ and  $(p_2)_1 = a$  and  $(p_3)_1 = b$  and  $(p_4)_2 = c$  and  $c \leq (p_1)_2$  and  $(p_1)_2 < (p_2)_2$  and  $(p_2)_2 \leq d$  and  $c \leq (p_3)_2$  and  $(p_3)_2 \leq d$  and

 $a < (p_4)_1$  and  $(p_4)_1 \le b$  and P is an arc from  $p_1$  to  $p_3$  and Q is an arc from  $p_2$  to  $p_4$  and  $P \subseteq$  ClosedInsideOfRectangle(a, b, c, d) and  $Q \subseteq$  ClosedInsideOfRectangle(a, b, c, d). Then P meets Q.

- (86) Let  $p_1$ ,  $p_2$ ,  $p_3$ ,  $p_4$  be points of  $\mathcal{E}_{\mathrm{T}}^2$ , a, b, c, d be real numbers, and f, g be maps from  $\mathbb{I}$  into  $\mathcal{E}_{\mathrm{T}}^2$ . Suppose that a < b and c < d and  $(p_1)_1 = a$  and  $(p_2)_1 = a$  and  $(p_3)_2 = c$  and  $(p_4)_2 = c$  and  $c \leq (p_1)_2$  and  $(p_1)_2 < (p_2)_2$  and  $(p_2)_2 \leq d$  and  $a < (p_4)_1$  and  $(p_4)_1 < (p_3)_1$  and  $(p_3)_1 \leq b$  and  $f(0) = p_1$  and  $f(1) = p_3$  and  $g(0) = p_2$  and  $g(1) = p_4$  and f is continuous and one-to-one and g is continuous and one-to-one and  $\operatorname{rng} f \subseteq \operatorname{ClosedInsideOfRectangle}(a, b, c, d)$ . Then  $\operatorname{rng} f$  meets  $\operatorname{rng} g$ .
- (87) Let  $p_1, p_2, p_3, p_4$  be points of  $\mathcal{E}_{\mathrm{T}}^2$ , a, b, c, d be real numbers, and P, Q be subsets of  $\mathcal{E}_{\mathrm{T}}^2$ . Suppose that a < b and c < d and  $(p_1)_1 = a$  and  $(p_2)_1 = a$ and  $(p_3)_2 = c$  and  $(p_4)_2 = c$  and  $c \leq (p_1)_2$  and  $(p_1)_2 < (p_2)_2$  and  $(p_2)_2 \leq d$ and  $a < (p_4)_1$  and  $(p_4)_1 < (p_3)_1$  and  $(p_3)_1 \leq b$  and P is an arc from  $p_1$  to  $p_3$ and Q is an arc from  $p_2$  to  $p_4$  and  $P \subseteq$  ClosedInsideOfRectangle(a, b, c, d)and  $Q \subseteq$  ClosedInsideOfRectangle(a, b, c, d). Then P meets Q.
- (88) Let  $p_1$ ,  $p_2$ ,  $p_3$ ,  $p_4$  be points of  $\mathcal{E}_{\mathrm{T}}^2$ , a, b, c, d be real numbers, and f, g be maps from  $\mathbb{I}$  into  $\mathcal{E}_{\mathrm{T}}^2$ . Suppose that a < b and c < d and  $(p_1)_1 = a$  and  $(p_2)_2 = d$  and  $(p_3)_2 = d$  and  $(p_4)_2 = d$  and  $c \leq (p_1)_2$  and  $(p_1)_2 \leq d$  and  $a \leq (p_2)_1$  and  $(p_2)_1 < (p_3)_1$  and  $(p_3)_1 < (p_4)_1$  and  $(p_4)_1 \leq b$  and  $f(0) = p_1$  and  $f(1) = p_3$  and  $g(0) = p_2$  and  $g(1) = p_4$  and f is continuous and one-to-one and g is continuous and one-to-one and  $\operatorname{rng} f \subseteq \operatorname{ClosedInsideOfRectangle}(a, b, c, d)$ . Then  $\operatorname{rng} f$  meets  $\operatorname{rng} g$ .
- (89) Let  $p_1, p_2, p_3, p_4$  be points of  $\mathcal{E}_{\mathrm{T}}^2$ , a, b, c, d be real numbers, and P, Q be subsets of  $\mathcal{E}_{\mathrm{T}}^2$ . Suppose that a < b and c < d and  $(p_1)_1 = a$ and  $(p_2)_2 = d$  and  $(p_3)_2 = d$  and  $(p_4)_2 = d$  and  $c \leq (p_1)_2$  and  $(p_1)_2 \leq d$  and  $a \leq (p_2)_1$  and  $(p_2)_1 < (p_3)_1$  and  $(p_3)_1 < (p_4)_1$ and  $(p_4)_1 \leq b$  and P is an arc from  $p_1$  to  $p_3$  and Q is an arc from  $p_2$  to  $p_4$  and  $P \subseteq$  ClosedInsideOfRectangle(a, b, c, d) and  $Q \subseteq$ ClosedInsideOfRectangle(a, b, c, d). Then P meets Q.
- (90) Let  $p_1$ ,  $p_2$ ,  $p_3$ ,  $p_4$  be points of  $\mathcal{E}_{\mathrm{T}}^2$ , a, b, c, d be real numbers, and f, g be maps from  $\mathbb{I}$  into  $\mathcal{E}_{\mathrm{T}}^2$ . Suppose that a < b and c < d and  $(p_1)_1 = a$  and  $(p_2)_2 = d$  and  $(p_3)_2 = d$  and  $(p_4)_1 = b$  and  $c \leq (p_1)_2$  and  $(p_1)_2 \leq d$  and  $a \leq (p_2)_1$  and  $(p_2)_1 < (p_3)_1$  and  $(p_3)_1 \leq b$  and  $c \leq (p_4)_2$  and  $(p_4)_2 \leq d$  and  $f(0) = p_1$  and  $f(1) = p_3$  and  $g(0) = p_2$  and  $g(1) = p_4$  and f is continuous and one-to-one and g is continuous and one-to-one and  $\operatorname{rng} f \subseteq \operatorname{ClosedInsideOfRectangle}(a, b, c, d)$ . Then  $\operatorname{rng} f$  meets  $\operatorname{rng} g$ .
- (91) Let  $p_1$ ,  $p_2$ ,  $p_3$ ,  $p_4$  be points of  $\mathcal{E}_{T}^2$ , a, b, c, d be real numbers, and P,

Q be subsets of  $\mathcal{E}_{\mathrm{T}}^2$ . Suppose that a < b and c < d and  $(p_1)_1 = a$ and  $(p_2)_2 = d$  and  $(p_3)_2 = d$  and  $(p_4)_1 = b$  and  $c \leq (p_1)_2$  and  $(p_1)_2 \leq d$  and  $a \leq (p_2)_1$  and  $(p_2)_1 < (p_3)_1$  and  $(p_3)_1 \leq b$  and  $c \leq (p_4)_2$  and  $(p_4)_2 \leq d$  and P is an arc from  $p_1$  to  $p_3$  and Q is an arc from  $p_2$  to  $p_4$  and  $P \subseteq$  ClosedInsideOfRectangle(a, b, c, d) and  $Q \subseteq$  ClosedInsideOfRectangle(a, b, c, d). Then P meets Q.

- (92) Let  $p_1$ ,  $p_2$ ,  $p_3$ ,  $p_4$  be points of  $\mathcal{E}_{\mathrm{T}}^2$ , a, b, c, d be real numbers, and f, g be maps from I into  $\mathcal{E}_{\mathrm{T}}^2$ . Suppose that a < b and c < d and  $(p_1)_1 = a$  and  $(p_2)_2 = d$  and  $(p_3)_2 = d$  and  $(p_4)_2 = c$  and  $c \leq (p_1)_2$  and  $(p_1)_2 \leq d$  and  $a \leq (p_2)_1$  and  $(p_2)_1 < (p_3)_1$  and  $(p_3)_1 \leq b$  and  $a < (p_4)_1$  and  $(p_4)_1 \leq b$  and  $f(0) = p_1$  and  $f(1) = p_3$  and  $g(0) = p_2$  and  $g(1) = p_4$  and f is continuous and one-to-one and  $rng f \subseteq$  ClosedInsideOfRectangle(a, b, c, d) and  $rng g \subseteq$  ClosedInsideOfRectangle(a, b, c, d). Then rng f meets rng g.
- (93) Let  $p_1, p_2, p_3, p_4$  be points of  $\mathcal{E}^2_{\mathrm{T}}$ , a, b, c, d be real numbers, and P, Q be subsets of  $\mathcal{E}^2_{\mathrm{T}}$ . Suppose that a < b and c < d and  $(p_1)_1 = a$ and  $(p_2)_2 = d$  and  $(p_3)_2 = d$  and  $(p_4)_2 = c$  and  $c \leq (p_1)_2$  and  $(p_1)_2 \leq d$  and  $a \leq (p_2)_1$  and  $(p_2)_1 < (p_3)_1$  and  $(p_3)_1 \leq b$  and  $a < (p_4)_1$  and  $(p_4)_1 \leq b$  and P is an arc from  $p_1$  to  $p_3$  and Q is an arc from  $p_2$  to  $p_4$  and  $P \subseteq$  ClosedInsideOfRectangle(a, b, c, d) and  $Q \subseteq$  ClosedInsideOfRectangle(a, b, c, d). Then P meets Q.
- (94) Let  $p_1$ ,  $p_2$ ,  $p_3$ ,  $p_4$  be points of  $\mathcal{E}_{\mathrm{T}}^2$ , a, b, c, d be real numbers, and f, g be maps from  $\mathbb{I}$  into  $\mathcal{E}_{\mathrm{T}}^2$ . Suppose that a < b and c < d and  $(p_1)_1 = a$  and  $(p_2)_2 = d$  and  $(p_3)_1 = b$  and  $(p_4)_1 = b$  and  $c \leq (p_1)_2$  and  $(p_1)_2 \leq d$  and  $a \leq (p_2)_1$  and  $(p_2)_1 \leq b$  and  $c \leq (p_4)_2$  and  $(p_4)_2 < (p_3)_2$  and  $(p_3)_2 \leq d$  and  $f(0) = p_1$  and  $f(1) = p_3$  and  $g(0) = p_2$  and  $g(1) = p_4$  and f is continuous and one-to-one and  $\operatorname{rng} f \subseteq \operatorname{ClosedInsideOfRectangle}(a, b, c, d)$  and  $\operatorname{rng} g \subseteq \operatorname{ClosedInsideOfRectangle}(a, b, c, d)$ . Then  $\operatorname{rng} f$  meets  $\operatorname{rng} g$ .
- (95) Let  $p_1$ ,  $p_2$ ,  $p_3$ ,  $p_4$  be points of  $\mathcal{E}_{\mathrm{T}}^2$ , a, b, c, d be real numbers, and P, Q be subsets of  $\mathcal{E}_{\mathrm{T}}^2$ . Suppose that a < b and c < d and  $(p_1)_1 = a$  and  $(p_2)_2 = d$  and  $(p_3)_1 = b$  and  $(p_4)_1 = b$  and  $c \leq (p_1)_2$ and  $(p_1)_2 \leq d$  and  $a \leq (p_2)_1$  and  $(p_2)_1 \leq b$  and  $c \leq (p_4)_2$  and  $(p_4)_2 < (p_3)_2$  and  $(p_3)_2 \leq d$  and P is an arc from  $p_1$  to  $p_3$  and Q is an arc from  $p_2$  to  $p_4$  and  $P \subseteq$  ClosedInsideOfRectangle(a, b, c, d) and  $Q \subseteq$  ClosedInsideOfRectangle(a, b, c, d). Then P meets Q.
- (96) Let  $p_1$ ,  $p_2$ ,  $p_3$ ,  $p_4$  be points of  $\mathcal{E}_{\mathrm{T}}^2$ , a, b, c, d be real numbers, and f, g be maps from  $\mathbb{I}$  into  $\mathcal{E}_{\mathrm{T}}^2$ . Suppose that a < b and c < d and  $(p_1)_1 = a$  and  $(p_2)_2 = d$  and  $(p_3)_1 = b$  and  $(p_4)_2 = c$  and  $c \leq (p_1)_2$  and  $(p_1)_2 \leq d$  and  $a \leq (p_2)_1$  and  $(p_2)_1 \leq b$  and  $c \leq (p_3)_2$  and  $(p_3)_2 \leq d$  and  $a < (p_4)_1$  and  $(p_4)_1 \leq b$  and  $f(0) = p_1$  and  $f(1) = p_3$  and  $g(0) = p_2$  and

 $g(1) = p_4$  and f is continuous and one-to-one and g is continuous and one-to-one and rng  $f \subseteq$  ClosedInsideOfRectangle(a, b, c, d) and rng  $g \subseteq$  ClosedInsideOfRectangle(a, b, c, d). Then rng f meets rng g.

- (97) Let  $p_1$ ,  $p_2$ ,  $p_3$ ,  $p_4$  be points of  $\mathcal{E}_{\mathrm{T}}^2$ , a, b, c, d be real numbers, and P, Q be subsets of  $\mathcal{E}_{\mathrm{T}}^2$ . Suppose that a < b and c < d and  $(p_1)_1 = a$ and  $(p_2)_2 = d$  and  $(p_3)_1 = b$  and  $(p_4)_2 = c$  and  $c \leq (p_1)_2$  and  $(p_1)_2 \leq d$  and  $a \leq (p_2)_1$  and  $(p_2)_1 \leq b$  and  $c \leq (p_3)_2$  and  $(p_3)_2 \leq d$ and  $a < (p_4)_1$  and  $(p_4)_1 \leq b$  and P is an arc from  $p_1$  to  $p_3$  and Qis an arc from  $p_2$  to  $p_4$  and  $P \subseteq$  ClosedInsideOfRectangle(a, b, c, d) and  $Q \subseteq$  ClosedInsideOfRectangle(a, b, c, d). Then P meets Q.
- (98) Let  $p_1$ ,  $p_2$ ,  $p_3$ ,  $p_4$  be points of  $\mathcal{E}_{\mathrm{T}}^2$ , a, b, c, d be real numbers, and f, g be maps from  $\mathbb{I}$  into  $\mathcal{E}_{\mathrm{T}}^2$ . Suppose that a < b and c < d and  $(p_1)_1 = a$  and  $(p_2)_2 = d$  and  $(p_3)_2 = c$  and  $(p_4)_2 = c$  and  $c \leq (p_1)_2$  and  $(p_1)_2 \leq d$  and  $a \leq (p_2)_1$  and  $(p_2)_1 \leq b$  and  $a < (p_4)_1$  and  $(p_4)_1 < (p_3)_1$  and  $(p_3)_1 \leq b$  and  $f(0) = p_1$  and  $f(1) = p_3$  and  $g(0) = p_2$  and  $g(1) = p_4$  and f is continuous and one-to-one and  $\operatorname{rng} f \subseteq \operatorname{ClosedInsideOfRectangle}(a, b, c, d)$ . Then  $\operatorname{rng} f$  meets  $\operatorname{rng} g$ .
- (99) Let  $p_1$ ,  $p_2$ ,  $p_3$ ,  $p_4$  be points of  $\mathcal{E}_{\mathrm{T}}^2$ , a, b, c, d be real numbers, and P, Q be subsets of  $\mathcal{E}_{\mathrm{T}}^2$ . Suppose that a < b and c < d and  $(p_1)_1 = a$  and  $(p_2)_2 = d$  and  $(p_3)_2 = c$  and  $(p_4)_2 = c$  and  $c \leq (p_1)_2$ and  $(p_1)_2 \leq d$  and  $a \leq (p_2)_1$  and  $(p_2)_1 \leq b$  and  $a < (p_4)_1$  and  $(p_4)_1 < (p_3)_1$  and  $(p_3)_1 \leq b$  and P is an arc from  $p_1$  to  $p_3$  and Q is an arc from  $p_2$  to  $p_4$  and  $P \subseteq$  ClosedInsideOfRectangle(a, b, c, d) and  $Q \subseteq$  ClosedInsideOfRectangle(a, b, c, d). Then P meets Q.
- (100) Let  $p_1$ ,  $p_2$ ,  $p_3$ ,  $p_4$  be points of  $\mathcal{E}_{\mathrm{T}}^2$ , a, b, c, d be real numbers, and f, g be maps from  $\mathbb{I}$  into  $\mathcal{E}_{\mathrm{T}}^2$ . Suppose that a < b and c < d and  $(p_1)_1 = a$  and  $(p_2)_1 = b$  and  $(p_3)_1 = b$  and  $(p_4)_1 = b$  and  $c \leq (p_1)_2$  and  $(p_1)_2 \leq d$  and  $c \leq (p_4)_2$  and  $(p_4)_2 < (p_3)_2$  and  $(p_3)_2 < (p_2)_2$  and  $(p_2)_2 \leq d$  and  $f(0) = p_1$  and  $f(1) = p_3$  and  $g(0) = p_2$  and  $g(1) = p_4$  and f is continuous and one-to-one and g is continuous and one-to-one and  $\operatorname{rng} g \subseteq$  ClosedInsideOfRectangle(a, b, c, d). Then  $\operatorname{rng} f$  meets  $\operatorname{rng} g$ .
- (101) Let  $p_1, p_2, p_3, p_4$  be points of  $\mathcal{E}_{\mathrm{T}}^2, a, b, c, d$  be real numbers, and P, Q be subsets of  $\mathcal{E}_{\mathrm{T}}^2$ . Suppose that a < b and c < d and  $(p_1)_1 = a$  and  $(p_2)_1 = b$ and  $(p_3)_1 = b$  and  $(p_4)_1 = b$  and  $c \leq (p_1)_2$  and  $(p_1)_2 \leq d$  and  $c \leq (p_4)_2$  and  $(p_4)_2 < (p_3)_2$  and  $(p_3)_2 < (p_2)_2$  and  $(p_2)_2 \leq d$  and P is an arc from  $p_1$  to  $p_3$ and Q is an arc from  $p_2$  to  $p_4$  and  $P \subseteq$  ClosedInsideOfRectangle(a, b, c, d)and  $Q \subseteq$  ClosedInsideOfRectangle(a, b, c, d). Then P meets Q.
- (102) Let  $p_1$ ,  $p_2$ ,  $p_3$ ,  $p_4$  be points of  $\mathcal{E}_T^2$ , a, b, c, d be real numbers, and f, g be maps from  $\mathbb{I}$  into  $\mathcal{E}_T^2$ . Suppose that a < b and c < d and

 $(p_1)_1 = a$  and  $(p_2)_1 = b$  and  $(p_3)_1 = b$  and  $(p_4)_2 = c$  and  $c \leq (p_1)_2$ and  $(p_1)_2 \leq d$  and  $c \leq (p_3)_2$  and  $(p_3)_2 < (p_2)_2$  and  $(p_2)_2 \leq d$  and  $a < (p_4)_1$  and  $(p_4)_1 \leq b$  and  $f(0) = p_1$  and  $f(1) = p_3$  and  $g(0) = p_2$ and  $g(1) = p_4$  and f is continuous and one-to-one and g is continuous and one-to-one and rng  $f \subseteq$  ClosedInsideOfRectangle(a, b, c, d) and rng  $g \subseteq$  ClosedInsideOfRectangle(a, b, c, d). Then rng f meets rng g.

- (103) Let  $p_1$ ,  $p_2$ ,  $p_3$ ,  $p_4$  be points of  $\mathcal{E}_{\mathrm{T}}^2$ , a, b, c, d be real numbers, and P, Q be subsets of  $\mathcal{E}_{\mathrm{T}}^2$ . Suppose that a < b and c < d and  $(p_1)_1 = a$ and  $(p_2)_1 = b$  and  $(p_3)_1 = b$  and  $(p_4)_2 = c$  and  $c \leq (p_1)_2$  and  $(p_1)_2 \leq d$  and  $c \leq (p_3)_2$  and  $(p_3)_2 < (p_2)_2$  and  $(p_2)_2 \leq d$  and  $a < (p_4)_1$  and  $(p_4)_1 \leq b$  and P is an arc from  $p_1$  to  $p_3$  and Q is an arc from  $p_2$  to  $p_4$  and  $P \subseteq$  ClosedInsideOfRectangle(a, b, c, d) and  $Q \subseteq$  ClosedInsideOfRectangle(a, b, c, d). Then P meets Q.
- (104) Let  $p_1$ ,  $p_2$ ,  $p_3$ ,  $p_4$  be points of  $\mathcal{E}_{\mathrm{T}}^2$ , a, b, c, d be real numbers, and f, g be maps from  $\mathbb{I}$  into  $\mathcal{E}_{\mathrm{T}}^2$ . Suppose that a < b and c < d and  $(p_1)_1 = a$  and  $(p_2)_1 = b$  and  $(p_3)_2 = c$  and  $(p_4)_2 = c$  and  $c \leq (p_1)_2$  and  $(p_1)_2 \leq d$  and  $c \leq (p_2)_2$  and  $(p_2)_2 \leq d$  and  $a < (p_4)_1$  and  $(p_4)_1 < (p_3)_1$  and  $(p_3)_1 \leq b$  and  $f(0) = p_1$  and  $f(1) = p_3$  and  $g(0) = p_2$  and  $g(1) = p_4$  and f is continuous and one-to-one and g is continuous and one-to-one and  $\operatorname{rng} f \subseteq \operatorname{ClosedInsideOfRectangle}(a, b, c, d)$ . Then  $\operatorname{rng} f$  meets  $\operatorname{rng} g$ .
- (105) Let  $p_1$ ,  $p_2$ ,  $p_3$ ,  $p_4$  be points of  $\mathcal{E}_{\mathrm{T}}^2$ , a, b, c, d be real numbers, and P, Q be subsets of  $\mathcal{E}_{\mathrm{T}}^2$ . Suppose that a < b and c < d and  $(p_1)_{\mathbf{1}} = a$  and  $(p_2)_{\mathbf{1}} = b$  and  $(p_3)_{\mathbf{2}} = c$  and  $(p_4)_{\mathbf{2}} = c$  and  $c \leq (p_1)_{\mathbf{2}}$ and  $(p_1)_{\mathbf{2}} \leq d$  and  $c \leq (p_2)_{\mathbf{2}}$  and  $(p_2)_{\mathbf{2}} \leq d$  and  $a < (p_4)_{\mathbf{1}}$  and  $(p_4)_{\mathbf{1}} < (p_3)_{\mathbf{1}}$  and  $(p_3)_{\mathbf{1}} \leq b$  and P is an arc from  $p_1$  to  $p_3$  and Q is an arc from  $p_2$  to  $p_4$  and  $P \subseteq$  ClosedInsideOfRectangle(a, b, c, d) and  $Q \subseteq$  ClosedInsideOfRectangle(a, b, c, d). Then P meets Q.
- (106) Let  $p_1$ ,  $p_2$ ,  $p_3$ ,  $p_4$  be points of  $\mathcal{E}_{\mathrm{T}}^2$ , a, b, c, d be real numbers, and f, g be maps from  $\mathbb{I}$  into  $\mathcal{E}_{\mathrm{T}}^2$ . Suppose that a < b and c < d and  $(p_1)_1 = a$  and  $(p_2)_2 = c$  and  $(p_3)_2 = c$  and  $(p_4)_2 = c$  and  $c \leq (p_1)_2$  and  $(p_1)_2 \leq d$  and  $a < (p_4)_1$  and  $(p_4)_1 < (p_3)_1$  and  $(p_3)_1 < (p_2)_1$  and  $(p_2)_1 \leq b$  and  $f(0) = p_1$  and  $f(1) = p_3$  and  $g(0) = p_2$  and  $g(1) = p_4$  and f is continuous and one-to-one and  $rng f \subseteq$  ClosedInsideOfRectangle(a, b, c, d) and  $rng g \subseteq$  ClosedInsideOfRectangle(a, b, c, d). Then rng f meets rng g.
- (107) Let  $p_1, p_2, p_3, p_4$  be points of  $\mathcal{E}_{\mathrm{T}}^2, a, b, c, d$  be real numbers, and P, Q be subsets of  $\mathcal{E}_{\mathrm{T}}^2$ . Suppose that a < b and c < d and  $(p_1)_1 = a$  and  $(p_2)_2 = c$ and  $(p_3)_2 = c$  and  $(p_4)_2 = c$  and  $c \leq (p_1)_2$  and  $(p_1)_2 \leq d$  and  $a < (p_4)_1$  and  $(p_4)_1 < (p_3)_1$  and  $(p_3)_1 < (p_2)_1$  and  $(p_2)_1 \leq b$  and P is an arc from  $p_1$  to  $p_3$ and Q is an arc from  $p_2$  to  $p_4$  and  $P \subseteq$  ClosedInsideOfRectangle(a, b, c, d)

and  $Q \subseteq$  ClosedInsideOfRectangle(a, b, c, d). Then P meets Q.

- (108) Let  $p_1$ ,  $p_2$ ,  $p_3$ ,  $p_4$  be points of  $\mathcal{E}_{\mathrm{T}}^2$ , a, b, c, d be real numbers, and f, g be maps from  $\mathbb{I}$  into  $\mathcal{E}_{\mathrm{T}}^2$ . Suppose that a < b and c < d and  $(p_1)_2 = d$  and  $(p_2)_2 = d$  and  $(p_3)_2 = d$  and  $(p_4)_2 = d$  and  $a \leq (p_1)_1$  and  $(p_1)_1 < (p_2)_1$  and  $(p_2)_1 < (p_3)_1$  and  $(p_3)_1 < (p_4)_1$  and  $(p_4)_1 \leq b$  and  $f(0) = p_1$  and  $f(1) = p_3$  and  $g(0) = p_2$  and  $g(1) = p_4$  and f is continuous and one-to-one and  $rng f \subseteq$  ClosedInsideOfRectangle(a, b, c, d) and  $rng g \subseteq$  ClosedInsideOfRectangle(a, b, c, d). Then rng f meets rng g.
- (109) Let  $p_1, p_2, p_3, p_4$  be points of  $\mathcal{E}_{\mathrm{T}}^2, a, b, c, d$  be real numbers, and P, Q be subsets of  $\mathcal{E}_{\mathrm{T}}^2$ . Suppose that a < b and c < d and  $(p_1)_2 = d$  and  $(p_2)_2 = d$ and  $(p_3)_2 = d$  and  $(p_4)_2 = d$  and  $a \leq (p_1)_1$  and  $(p_1)_1 < (p_2)_1$  and  $(p_2)_1 < (p_3)_1$  and  $(p_3)_1 < (p_4)_1$  and  $(p_4)_1 \leq b$  and P is an arc from  $p_1$  to  $p_3$ and Q is an arc from  $p_2$  to  $p_4$  and  $P \subseteq \text{ClosedInsideOfRectangle}(a, b, c, d)$ and  $Q \subseteq \text{ClosedInsideOfRectangle}(a, b, c, d)$ . Then P meets Q.
- (110) Let  $p_1$ ,  $p_2$ ,  $p_3$ ,  $p_4$  be points of  $\mathcal{E}_{\mathrm{T}}^2$ , a, b, c, d be real numbers, and f, g be maps from  $\mathbb{I}$  into  $\mathcal{E}_{\mathrm{T}}^2$ . Suppose that a < b and c < d and  $(p_1)_2 = d$  and  $(p_2)_2 = d$  and  $(p_3)_2 = d$  and  $(p_4)_1 = b$  and  $a \leq (p_1)_1$  and  $(p_1)_1 < (p_2)_1$  and  $(p_2)_1 < (p_3)_1$  and  $(p_3)_1 \leq b$  and  $c \leq (p_4)_2$  and  $(p_4)_2 \leq d$  and  $f(0) = p_1$  and  $f(1) = p_3$  and  $g(0) = p_2$  and  $g(1) = p_4$  and f is continuous and one-to-one and  $\operatorname{rng} f \subseteq \operatorname{ClosedInsideOfRectangle}(a, b, c, d)$  and  $\operatorname{rng} g \subseteq \operatorname{ClosedInsideOfRectangle}(a, b, c, d)$ . Then  $\operatorname{rng} f$  meets  $\operatorname{rng} g$ .
- (111) Let  $p_1$ ,  $p_2$ ,  $p_3$ ,  $p_4$  be points of  $\mathcal{E}_{\mathrm{T}}^2$ , a, b, c, d be real numbers, and P, Q be subsets of  $\mathcal{E}_{\mathrm{T}}^2$ . Suppose that a < b and c < d and  $(p_1)_2 = d$  and  $(p_2)_2 = d$  and  $(p_3)_2 = d$  and  $(p_4)_1 = b$  and  $a \leq (p_1)_1$ and  $(p_1)_1 < (p_2)_1$  and  $(p_2)_1 < (p_3)_1$  and  $(p_3)_1 \leq b$  and  $c \leq (p_4)_2$ and  $(p_4)_2 \leq d$  and P is an arc from  $p_1$  to  $p_3$  and Q is an arc from  $p_2$  to  $p_4$  and  $P \subseteq$  ClosedInsideOfRectangle(a, b, c, d) and  $Q \subseteq$ ClosedInsideOfRectangle(a, b, c, d). Then P meets Q.
- (112) Let  $p_1$ ,  $p_2$ ,  $p_3$ ,  $p_4$  be points of  $\mathcal{E}_{\mathrm{T}}^2$ , a, b, c, d be real numbers, and f, g be maps from  $\mathbb{I}$  into  $\mathcal{E}_{\mathrm{T}}^2$ . Suppose that a < b and c < d and  $(p_1)_2 = d$  and  $(p_2)_2 = d$  and  $(p_3)_2 = d$  and  $(p_4)_2 = c$  and  $a \leq (p_1)_1$  and  $(p_1)_1 < (p_2)_1$  and  $(p_2)_1 < (p_3)_1$  and  $(p_3)_1 \leq b$  and  $a < (p_4)_1$  and  $(p_4)_1 \leq b$  and  $f(0) = p_1$  and  $f(1) = p_3$  and  $g(0) = p_2$  and  $g(1) = p_4$  and f is continuous and one-to-one and  $\operatorname{rng} f \subseteq \operatorname{ClosedInsideOfRectangle}(a, b, c, d)$  and  $\operatorname{rng} g \subseteq \operatorname{ClosedInsideOfRectangle}(a, b, c, d)$ . Then  $\operatorname{rng} f$  meets  $\operatorname{rng} g$ .
- (113) Let  $p_1$ ,  $p_2$ ,  $p_3$ ,  $p_4$  be points of  $\mathcal{E}_{\mathrm{T}}^2$ , a, b, c, d be real numbers, and P, Q be subsets of  $\mathcal{E}_{\mathrm{T}}^2$ . Suppose that a < b and c < d and  $(p_1)_2 = d$ and  $(p_2)_2 = d$  and  $(p_3)_2 = d$  and  $(p_4)_2 = c$  and  $a \leq (p_1)_1$  and

 $(p_1)_1 < (p_2)_1$  and  $(p_2)_1 < (p_3)_1$  and  $(p_3)_1 \leq b$  and  $a < (p_4)_1$ and  $(p_4)_1 \leq b$  and P is an arc from  $p_1$  to  $p_3$  and Q is an arc from  $p_2$  to  $p_4$  and  $P \subseteq$  ClosedInsideOfRectangle(a, b, c, d) and  $Q \subseteq$ ClosedInsideOfRectangle(a, b, c, d). Then P meets Q.

- (114) Let  $p_1$ ,  $p_2$ ,  $p_3$ ,  $p_4$  be points of  $\mathcal{E}_{\mathrm{T}}^2$ , a, b, c, d be real numbers, and f, g be maps from  $\mathbb{I}$  into  $\mathcal{E}_{\mathrm{T}}^2$ . Suppose that a < b and c < d and  $(p_1)_2 = d$  and  $(p_2)_2 = d$  and  $(p_3)_1 = b$  and  $(p_4)_1 = b$  and  $a \leq (p_1)_1$  and  $(p_1)_1 < (p_2)_1$  and  $(p_2)_1 \leq b$  and  $c \leq (p_4)_2$  and  $(p_4)_2 < (p_3)_2$  and  $(p_3)_2 \leq d$  and  $f(0) = p_1$  and  $f(1) = p_3$  and  $g(0) = p_2$  and  $g(1) = p_4$  and f is continuous and one-to-one and  $\operatorname{rng} f \subseteq \operatorname{ClosedInsideOfRectangle}(a, b, c, d)$  and  $\operatorname{rng} g \subseteq \operatorname{ClosedInsideOfRectangle}(a, b, c, d)$ . Then  $\operatorname{rng} f$  meets  $\operatorname{rng} g$ .
- (115) Let  $p_1, p_2, p_3, p_4$  be points of  $\mathcal{E}_{\mathrm{T}}^2, a, b, c, d$  be real numbers, and P, Q be subsets of  $\mathcal{E}_{\mathrm{T}}^2$ . Suppose that a < b and c < d and  $(p_1)_2 = d$  and  $(p_2)_2 = d$ and  $(p_3)_1 = b$  and  $(p_4)_1 = b$  and  $a \leq (p_1)_1$  and  $(p_1)_1 < (p_2)_1$  and  $(p_2)_1 \leq b$ and  $c \leq (p_4)_2$  and  $(p_4)_2 < (p_3)_2$  and  $(p_3)_2 \leq d$  and P is an arc from  $p_1$  to  $p_3$ and Q is an arc from  $p_2$  to  $p_4$  and  $P \subseteq$  ClosedInsideOfRectangle(a, b, c, d)and  $Q \subseteq$  ClosedInsideOfRectangle(a, b, c, d). Then P meets Q.
- (116) Let  $p_1$ ,  $p_2$ ,  $p_3$ ,  $p_4$  be points of  $\mathcal{E}_{\mathrm{T}}^2$ , a, b, c, d be real numbers, and f, g be maps from  $\mathbb{I}$  into  $\mathcal{E}_{\mathrm{T}}^2$ . Suppose that a < b and c < d and  $(p_1)_2 = d$  and  $(p_2)_2 = d$  and  $(p_3)_1 = b$  and  $(p_4)_2 = c$  and  $a \leq (p_1)_1$  and  $(p_1)_1 < (p_2)_1$  and  $(p_2)_1 \leq b$  and  $c \leq (p_3)_2$  and  $(p_3)_2 \leq d$  and  $a < (p_4)_1$  and  $(p_4)_1 \leq b$  and  $f(0) = p_1$  and  $f(1) = p_3$  and  $g(0) = p_2$  and  $g(1) = p_4$  and f is continuous and one-to-one and  $\operatorname{rng} f \subseteq \operatorname{ClosedInsideOfRectangle}(a, b, c, d)$ . Then  $\operatorname{rng} f$  meets  $\operatorname{rng} g$ .
- (117) Let  $p_1$ ,  $p_2$ ,  $p_3$ ,  $p_4$  be points of  $\mathcal{E}_{\mathrm{T}}^2$ , a, b, c, d be real numbers, and P, Q be subsets of  $\mathcal{E}_{\mathrm{T}}^2$ . Suppose that a < b and c < d and  $(p_1)_2 = d$ and  $(p_2)_2 = d$  and  $(p_3)_1 = b$  and  $(p_4)_2 = c$  and  $a \leq (p_1)_1$  and  $(p_1)_1 < (p_2)_1$  and  $(p_2)_1 \leq b$  and  $c \leq (p_3)_2$  and  $(p_3)_2 \leq d$  and  $a < (p_4)_1$  and  $(p_4)_1 \leq b$  and P is an arc from  $p_1$  to  $p_3$  and Q is an arc from  $p_2$  to  $p_4$  and  $P \subseteq$  ClosedInsideOfRectangle(a, b, c, d) and  $Q \subseteq$  ClosedInsideOfRectangle(a, b, c, d). Then P meets Q.
- (118) Let  $p_1$ ,  $p_2$ ,  $p_3$ ,  $p_4$  be points of  $\mathcal{E}_{\mathrm{T}}^2$ , a, b, c, d be real numbers, and f, g be maps from  $\mathbb{I}$  into  $\mathcal{E}_{\mathrm{T}}^2$ . Suppose that a < b and c < d and  $(p_1)_2 = d$  and  $(p_2)_2 = d$  and  $(p_3)_2 = c$  and  $(p_4)_2 = c$  and  $a \leq (p_1)_1$  and  $(p_1)_1 < (p_2)_1$  and  $(p_2)_1 \leq b$  and  $a < (p_4)_1$  and  $(p_4)_1 < (p_3)_1$  and  $(p_3)_1 \leq b$  and  $f(0) = p_1$  and  $f(1) = p_3$  and  $g(0) = p_2$  and  $g(1) = p_4$  and f is continuous and one-to-one and  $\operatorname{rng} f \subseteq \operatorname{ClosedInsideOfRectangle}(a, b, c, d)$  and  $\operatorname{rng} g \subseteq \operatorname{ClosedInsideOfRectangle}(a, b, c, d)$ . Then  $\operatorname{rng} f$  meets  $\operatorname{rng} g$ .

- (119) Let  $p_1, p_2, p_3, p_4$  be points of  $\mathcal{E}_{\mathrm{T}}^2$ , a, b, c, d be real numbers, and P, Q be subsets of  $\mathcal{E}_{\mathrm{T}}^2$ . Suppose that a < b and c < d and  $(p_1)_2 = d$  and  $(p_2)_2 = d$ and  $(p_3)_2 = c$  and  $(p_4)_2 = c$  and  $a \leq (p_1)_1$  and  $(p_1)_1 < (p_2)_1$  and  $(p_2)_1 \leq b$ and  $a < (p_4)_1$  and  $(p_4)_1 < (p_3)_1$  and  $(p_3)_1 \leq b$  and P is an arc from  $p_1$  to  $p_3$ and Q is an arc from  $p_2$  to  $p_4$  and  $P \subseteq \text{ClosedInsideOfRectangle}(a, b, c, d)$ and  $Q \subseteq \text{ClosedInsideOfRectangle}(a, b, c, d)$ . Then P meets Q.
- (120) Let  $p_1$ ,  $p_2$ ,  $p_3$ ,  $p_4$  be points of  $\mathcal{E}_{\mathrm{T}}^2$ , a, b, c, d be real numbers, and f, g be maps from  $\mathbb{I}$  into  $\mathcal{E}_{\mathrm{T}}^2$ . Suppose that a < b and c < d and  $(p_1)_2 = d$  and  $(p_2)_1 = b$  and  $(p_3)_1 = b$  and  $(p_4)_1 = b$  and  $a \leq (p_1)_1$  and  $(p_1)_1 \leq b$  and  $d \geq (p_2)_2$  and  $(p_2)_2 > (p_3)_2$  and  $(p_3)_2 > (p_4)_2$  and  $(p_4)_2 \geq c$  and  $f(0) = p_1$  and  $f(1) = p_3$  and  $g(0) = p_2$  and  $g(1) = p_4$  and f is continuous and one-to-one and g is continuous and one-to-one and  $\operatorname{rng} f \subseteq \operatorname{ClosedInsideOfRectangle}(a, b, c, d)$ . Then  $\operatorname{rng} f$  meets  $\operatorname{rng} g$ .
- (121) Let  $p_1, p_2, p_3, p_4$  be points of  $\mathcal{E}_{\mathrm{T}}^2$ , a, b, c, d be real numbers, and P, Q be subsets of  $\mathcal{E}_{\mathrm{T}}^2$ . Suppose that a < b and c < d and  $(p_1)_2 = d$  and  $(p_2)_1 = b$ and  $(p_3)_1 = b$  and  $(p_4)_1 = b$  and  $a \leq (p_1)_1$  and  $(p_1)_1 \leq b$  and  $d \geq (p_2)_2$  and  $(p_2)_2 > (p_3)_2$  and  $(p_3)_2 > (p_4)_2$  and  $(p_4)_2 \geq c$  and P is an arc from  $p_1$  to  $p_3$ and Q is an arc from  $p_2$  to  $p_4$  and  $P \subseteq \text{ClosedInsideOfRectangle}(a, b, c, d)$ and  $Q \subseteq \text{ClosedInsideOfRectangle}(a, b, c, d)$ . Then P meets Q.
- (122) Let  $p_1$ ,  $p_2$ ,  $p_3$ ,  $p_4$  be points of  $\mathcal{E}_{\mathrm{T}}^2$ , a, b, c, d be real numbers, and f, g be maps from  $\mathbb{I}$  into  $\mathcal{E}_{\mathrm{T}}^2$ . Suppose that a < b and c < d and  $(p_1)_2 = d$  and  $(p_2)_1 = b$  and  $(p_3)_1 = b$  and  $(p_4)_2 = c$  and  $a \leq (p_1)_1$  and  $(p_1)_1 \leq b$  and  $d \geq (p_2)_2$  and  $(p_2)_2 > (p_3)_2$  and  $(p_3)_2 \geq c$  and  $a < (p_4)_1$  and  $(p_4)_1 \leq b$  and  $f(0) = p_1$  and  $f(1) = p_3$  and  $g(0) = p_2$  and  $g(1) = p_4$  and f is continuous and one-to-one and  $\operatorname{rng} f \subseteq \operatorname{ClosedInsideOfRectangle}(a, b, c, d)$  and  $\operatorname{rng} g \subseteq \operatorname{ClosedInsideOfRectangle}(a, b, c, d)$ . Then  $\operatorname{rng} f$  meets  $\operatorname{rng} g$ .
- (123) Let  $p_1$ ,  $p_2$ ,  $p_3$ ,  $p_4$  be points of  $\mathcal{E}_{\mathrm{T}}^2$ , a, b, c, d be real numbers, and P, Q be subsets of  $\mathcal{E}_{\mathrm{T}}^2$ . Suppose that a < b and c < d and  $(p_1)_2 = d$ and  $(p_2)_1 = b$  and  $(p_3)_1 = b$  and  $(p_4)_2 = c$  and  $a \leq (p_1)_1$  and  $(p_1)_1 \leq b$  and  $d \geq (p_2)_2$  and  $(p_2)_2 > (p_3)_2$  and  $(p_3)_2 \geq c$  and  $a < (p_4)_1$  and  $(p_4)_1 \leq b$  and P is an arc from  $p_1$  to  $p_3$  and Q is an arc from  $p_2$  to  $p_4$  and  $P \subseteq$  ClosedInsideOfRectangle(a, b, c, d) and  $Q \subseteq$  ClosedInsideOfRectangle(a, b, c, d). Then P meets Q.
- (124) Let  $p_1$ ,  $p_2$ ,  $p_3$ ,  $p_4$  be points of  $\mathcal{E}_T^2$ , a, b, c, d be real numbers, and f, g be maps from  $\mathbb{I}$  into  $\mathcal{E}_T^2$ . Suppose that a < b and c < d and  $(p_1)_2 = d$  and  $(p_2)_1 = b$  and  $(p_3)_2 = c$  and  $(p_4)_2 = c$  and  $a \leq (p_1)_1$  and  $(p_1)_1 \leq b$  and  $c \leq (p_2)_2$  and  $(p_2)_2 \leq d$  and  $a < (p_4)_1$  and  $(p_4)_1 < (p_3)_1$  and  $(p_3)_1 \leq b$  and  $f(0) = p_1$  and  $f(1) = p_3$  and  $g(0) = p_2$  and  $g(1) = p_4$  and f is continuous and one-to-one and g is continu-

ous and one-to-one and  $\operatorname{rng} f \subseteq \operatorname{ClosedInsideOfRectangle}(a, b, c, d)$  and  $\operatorname{rng} g \subseteq \operatorname{ClosedInsideOfRectangle}(a, b, c, d)$ . Then  $\operatorname{rng} f$  meets  $\operatorname{rng} g$ .

- (125) Let  $p_1$ ,  $p_2$ ,  $p_3$ ,  $p_4$  be points of  $\mathcal{E}_{\mathrm{T}}^2$ , a, b, c, d be real numbers, and P, Q be subsets of  $\mathcal{E}_{\mathrm{T}}^2$ . Suppose that a < b and c < d and  $(p_1)_2 = d$  and  $(p_2)_1 = b$  and  $(p_3)_2 = c$  and  $(p_4)_2 = c$  and  $a \leq (p_1)_1$ and  $(p_1)_1 \leq b$  and  $c \leq (p_2)_2$  and  $(p_2)_2 \leq d$  and  $a < (p_4)_1$  and  $(p_4)_1 < (p_3)_1$  and  $(p_3)_1 \leq b$  and P is an arc from  $p_1$  to  $p_3$  and Q is an arc from  $p_2$  to  $p_4$  and  $P \subseteq$  ClosedInsideOfRectangle(a, b, c, d) and  $Q \subseteq$  ClosedInsideOfRectangle(a, b, c, d). Then P meets Q.
- (126) Let  $p_1$ ,  $p_2$ ,  $p_3$ ,  $p_4$  be points of  $\mathcal{E}_{\mathrm{T}}^2$ , a, b, c, d be real numbers, and f, g be maps from  $\mathbb{I}$  into  $\mathcal{E}_{\mathrm{T}}^2$ . Suppose that a < b and c < d and  $(p_1)_2 = d$  and  $(p_2)_2 = c$  and  $(p_3)_2 = c$  and  $(p_4)_2 = c$  and  $a \leq (p_1)_1$  and  $(p_1)_1 \leq b$  and  $a < (p_4)_1$  and  $(p_4)_1 < (p_3)_1$  and  $(p_3)_1 < (p_2)_1$  and  $(p_2)_1 \leq b$  and  $f(0) = p_1$  and  $f(1) = p_3$  and  $g(0) = p_2$  and  $g(1) = p_4$  and f is continuous and one-to-one and  $\operatorname{rng} f \subseteq \operatorname{ClosedInsideOfRectangle}(a, b, c, d)$  and  $\operatorname{rng} g \subseteq \operatorname{ClosedInsideOfRectangle}(a, b, c, d)$ . Then  $\operatorname{rng} f$  meets  $\operatorname{rng} g$ .
- (127) Let  $p_1, p_2, p_3, p_4$  be points of  $\mathcal{E}_{\mathrm{T}}^2, a, b, c, d$  be real numbers, and P, Q be subsets of  $\mathcal{E}_{\mathrm{T}}^2$ . Suppose that a < b and c < d and  $(p_1)_2 = d$  and  $(p_2)_2 = c$ and  $(p_3)_2 = c$  and  $(p_4)_2 = c$  and  $a \leq (p_1)_1$  and  $(p_1)_1 \leq b$  and  $a < (p_4)_1$  and  $(p_4)_1 < (p_3)_1$  and  $(p_3)_1 < (p_2)_1$  and  $(p_2)_1 \leq b$  and P is an arc from  $p_1$  to  $p_3$ and Q is an arc from  $p_2$  to  $p_4$  and  $P \subseteq$  ClosedInsideOfRectangle(a, b, c, d)and  $Q \subseteq$  ClosedInsideOfRectangle(a, b, c, d). Then P meets Q.
- (128) Let  $p_1$ ,  $p_2$ ,  $p_3$ ,  $p_4$  be points of  $\mathcal{E}_{\mathrm{T}}^2$ , a, b, c, d be real numbers, and f, g be maps from  $\mathbb{I}$  into  $\mathcal{E}_{\mathrm{T}}^2$ . Suppose that a < b and c < d and  $(p_1)_1 = b$  and  $(p_2)_1 = b$  and  $(p_3)_1 = b$  and  $(p_4)_1 = b$  and  $d \ge (p_1)_2$  and  $(p_1)_2 > (p_2)_2$  and  $(p_2)_2 > (p_3)_2$  and  $(p_3)_2 > (p_4)_2$  and  $(p_4)_2 \ge c$  and  $f(0) = p_1$  and  $f(1) = p_3$  and  $g(0) = p_2$  and  $g(1) = p_4$  and f is continuous and one-to-one and g is continuous and one-to-one and rng  $f \subseteq$  ClosedInsideOfRectangle(a, b, c, d) and rng  $g \subseteq$  ClosedInsideOfRectangle(a, b, c, d). Then rng f meets rng g.
- (129) Let  $p_1$ ,  $p_2$ ,  $p_3$ ,  $p_4$  be points of  $\mathcal{E}^2_{\mathrm{T}}$ , a, b, c, d be real numbers, and P, Q be subsets of  $\mathcal{E}^2_{\mathrm{T}}$ . Suppose that a < b and c < d and  $(p_1)_1 = b$  and  $(p_2)_1 = b$  and  $(p_3)_1 = b$  and  $(p_4)_1 = b$  and  $d \ge (p_1)_2$  and  $(p_1)_2 > (p_2)_2$  and  $(p_2)_2 > (p_3)_2$  and  $(p_3)_2 > (p_4)_2$  and  $(p_4)_2 \ge c$  and P is an arc from  $p_1$  to  $p_3$ and Q is an arc from  $p_2$  to  $p_4$  and  $P \subseteq \text{ClosedInsideOfRectangle}(a, b, c, d)$ and  $Q \subseteq \text{ClosedInsideOfRectangle}(a, b, c, d)$ . Then P meets Q.
- (130) Let  $p_1$ ,  $p_2$ ,  $p_3$ ,  $p_4$  be points of  $\mathcal{E}_{\mathrm{T}}^2$ , a, b, c, d be real numbers, and f, g be maps from  $\mathbb{I}$  into  $\mathcal{E}_{\mathrm{T}}^2$ . Suppose that a < b and c < d and  $(p_1)_1 = b$  and  $(p_2)_1 = b$  and  $(p_3)_1 = b$  and  $(p_4)_2 = c$  and  $d \ge (p_1)_2$  and  $(p_1)_2 > (p_2)_2$  and  $(p_2)_2 > (p_3)_2$  and  $(p_3)_2 \ge c$  and  $a < (p_4)_1$  and  $(p_4)_1 \le b$  and  $f(0) = p_1$  and  $f(1) = p_3$  and  $g(0) = p_2$  and

 $g(1) = p_4$  and f is continuous and one-to-one and g is continuous and one-to-one and rng  $f \subseteq$  ClosedInsideOfRectangle(a, b, c, d) and rng  $g \subseteq$  ClosedInsideOfRectangle(a, b, c, d). Then rng f meets rng g.

- (131) Let  $p_1$ ,  $p_2$ ,  $p_3$ ,  $p_4$  be points of  $\mathcal{E}^2_{\mathrm{T}}$ , a, b, c, d be real numbers, and P, Q be subsets of  $\mathcal{E}^2_{\mathrm{T}}$ . Suppose that a < b and c < d and  $(p_1)_1 = b$ and  $(p_2)_1 = b$  and  $(p_3)_1 = b$  and  $(p_4)_2 = c$  and  $d \ge (p_1)_2$  and  $(p_1)_2 > (p_2)_2$  and  $(p_2)_2 > (p_3)_2$  and  $(p_3)_2 \ge c$  and  $a < (p_4)_1$ and  $(p_4)_1 \le b$  and P is an arc from  $p_1$  to  $p_3$  and Q is an arc from  $p_2$  to  $p_4$  and  $P \subseteq$  ClosedInsideOfRectangle(a, b, c, d) and  $Q \subseteq$ ClosedInsideOfRectangle(a, b, c, d). Then P meets Q.
- (132) Let  $p_1$ ,  $p_2$ ,  $p_3$ ,  $p_4$  be points of  $\mathcal{E}_{\mathrm{T}}^2$ , a, b, c, d be real numbers, and f, g be maps from  $\mathbb{I}$  into  $\mathcal{E}_{\mathrm{T}}^2$ . Suppose that a < b and c < d and  $(p_1)_1 = b$  and  $(p_2)_1 = b$  and  $(p_3)_2 = c$  and  $(p_4)_2 = c$  and  $d \ge (p_1)_2$  and  $(p_1)_2 > (p_2)_2$  and  $(p_2)_2 \ge c$  and  $b \ge (p_3)_1$  and  $(p_3)_1 > (p_4)_1$  and  $(p_4)_1 > a$  and  $f(0) = p_1$  and  $f(1) = p_3$  and  $g(0) = p_2$  and  $g(1) = p_4$  and f is continuous and one-to-one and  $\operatorname{rng} f \subseteq \operatorname{ClosedInsideOfRectangle}(a, b, c, d)$  and  $\operatorname{rng} g \subseteq \operatorname{ClosedInsideOfRectangle}(a, b, c, d)$ . Then  $\operatorname{rng} f$  meets  $\operatorname{rng} g$ .
- (133) Let  $p_1, p_2, p_3, p_4$  be points of  $\mathcal{E}_{\mathrm{T}}^2, a, b, c, d$  be real numbers, and P, Q be subsets of  $\mathcal{E}_{\mathrm{T}}^2$ . Suppose that a < b and c < d and  $(p_1)_1 = b$  and  $(p_2)_1 = b$ and  $(p_3)_2 = c$  and  $(p_4)_2 = c$  and  $d \ge (p_1)_2$  and  $(p_1)_2 > (p_2)_2$  and  $(p_2)_2 \ge c$ and  $b \ge (p_3)_1$  and  $(p_3)_1 > (p_4)_1$  and  $(p_4)_1 > a$  and P is an arc from  $p_1$  to  $p_3$ and Q is an arc from  $p_2$  to  $p_4$  and  $P \subseteq$  ClosedInsideOfRectangle(a, b, c, d)and  $Q \subseteq$  ClosedInsideOfRectangle(a, b, c, d). Then P meets Q.
- (134) Let  $p_1$ ,  $p_2$ ,  $p_3$ ,  $p_4$  be points of  $\mathcal{E}_{\mathrm{T}}^2$ , a, b, c, d be real numbers, and f, g be maps from  $\mathbb{I}$  into  $\mathcal{E}_{\mathrm{T}}^2$ . Suppose that a < b and c < d and  $(p_1)_1 = b$  and  $(p_2)_2 = c$  and  $(p_3)_2 = c$  and  $(p_4)_2 = c$  and  $c \leq (p_1)_2$  and  $(p_1)_2 \leq d$  and  $b \geq (p_2)_1$  and  $(p_2)_1 > (p_3)_1$  and  $(p_3)_1 > (p_4)_1$  and  $(p_4)_1 > a$  and  $f(0) = p_1$  and  $f(1) = p_3$  and  $g(0) = p_2$  and  $g(1) = p_4$  and f is continuous and one-to-one and g is continuous and one-to-one and  $\operatorname{rng} f \subseteq \operatorname{ClosedInsideOfRectangle}(a, b, c, d)$  and  $\operatorname{rng} g \subseteq \operatorname{ClosedInsideOfRectangle}(a, b, c, d)$ . Then  $\operatorname{rng} f$  meets  $\operatorname{rng} g$ .
- (135) Let  $p_1, p_2, p_3, p_4$  be points of  $\mathcal{E}_{\mathrm{T}}^2$ , a, b, c, d be real numbers, and P, Q be subsets of  $\mathcal{E}_{\mathrm{T}}^2$ . Suppose that a < b and c < d and  $(p_1)_1 = b$  and  $(p_2)_2 = c$ and  $(p_3)_2 = c$  and  $(p_4)_2 = c$  and  $c \leq (p_1)_2$  and  $(p_1)_2 \leq d$  and  $b \geq (p_2)_1$  and  $(p_2)_1 > (p_3)_1$  and  $(p_3)_1 > (p_4)_1$  and  $(p_4)_1 > a$  and P is an arc from  $p_1$  to  $p_3$ and Q is an arc from  $p_2$  to  $p_4$  and  $P \subseteq \text{ClosedInsideOfRectangle}(a, b, c, d)$ and  $Q \subseteq \text{ClosedInsideOfRectangle}(a, b, c, d)$ . Then P meets Q.
- (136) Let  $p_1$ ,  $p_2$ ,  $p_3$ ,  $p_4$  be points of  $\mathcal{E}_T^2$ , a, b, c, d be real numbers, and f, g be maps from  $\mathbb{I}$  into  $\mathcal{E}_T^2$ . Suppose that a < b and c < d and  $(p_1)_2 = c$  and  $(p_2)_2 = c$  and  $(p_3)_2 = c$  and  $(p_4)_2 = c$  and  $b \ge (p_1)_1$  and  $(p_1)_1 > (p_2)_1$  and

 $(p_2)_1 > (p_3)_1$  and  $(p_3)_1 > (p_4)_1$  and  $(p_4)_1 > a$  and  $f(0) = p_1$  and  $f(1) = p_3$  and  $g(0) = p_2$  and  $g(1) = p_4$  and f is continuous and one-to-one and g is continuous and one-to-one and rng  $f \subseteq$  ClosedInsideOfRectangle(a, b, c, d) and rng  $g \subseteq$  ClosedInsideOfRectangle(a, b, c, d). Then rng f meets rng g.

(137) Let  $p_1$ ,  $p_2$ ,  $p_3$ ,  $p_4$  be points of  $\mathcal{E}_{\mathrm{T}}^2$ , a, b, c, d be real numbers, and P, Q be subsets of  $\mathcal{E}_{\mathrm{T}}^2$ . Suppose that a < b and c < d and  $(p_1)_2 = c$  and  $(p_2)_2 = c$  and  $(p_3)_2 = c$  and  $(p_4)_2 = c$  and  $b \ge (p_1)_1$  and  $(p_1)_1 > (p_2)_1$  and  $(p_2)_1 > (p_3)_1$  and  $(p_3)_1 > (p_4)_1$  and  $(p_4)_1 > a$  and P is an arc from  $p_1$  to  $p_3$ and Q is an arc from  $p_2$  to  $p_4$  and  $P \subseteq \text{ClosedInsideOfRectangle}(a, b, c, d)$ and  $Q \subseteq \text{ClosedInsideOfRectangle}(a, b, c, d)$ . Then P meets Q.

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# Preliminaries to Mathematical Morphology and Its Properties

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**Summary.** The article is a translation of chapter 2 of the book *Mathematical Morphological Method and Application* by Changqing Tang, Hongbo Lu, Zheng Huang, Fang Zhang, Science Press, China, 1990. In this article, the basic mathematical morphological operators such as Erosion, Dilation, Adjunction Opening, Adjunction Closing and their properties are given. And these operators are usually used in processing and analysing the images.

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The terminology and notation used here are introduced in the following articles: [5], [1], [2], [6], [4], and [3].

## 1. The Definition of Erosion and Dilation and Their Algebraic Properties

In this paper n denotes a natural number and q, y, b denote points of  $\mathcal{E}_{\mathrm{T}}^{n}$ . Let us consider n, let p be a point of  $\mathcal{E}_{\mathrm{T}}^{n}$ , and let X be a subset of  $\mathcal{E}_{\mathrm{T}}^{n}$ . The functor X + p yielding a subset of  $\mathcal{E}_{\mathrm{T}}^{n}$  is defined by:

(Def. 1)  $X + p = \{q + p : q \in X\}.$ 

Let us consider n and let X be a subset of  $\mathcal{E}_{\mathrm{T}}^{n}$ . The functor X! yielding a subset of  $\mathcal{E}_{\mathrm{T}}^{n}$  is defined as follows:

(Def. 2)  $X! = \{-q : q \in X\}.$ 

Let us consider n and let X, B be subsets of  $\mathcal{E}_{T}^{n}$ . The functor  $X \ominus B$  yields a subset of  $\mathcal{E}_{T}^{n}$  and is defined as follows:

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(Def. 3)  $X \ominus B = \{y : B + y \subseteq X\}.$ 

Let us consider n and let X, B be subsets of  $\mathcal{E}^n_T$ . The functor  $X \oplus B$  yields a subset of  $\mathcal{E}^n_T$  and is defined as follows:

(Def. 4)  $X \oplus B = \{y + b : y \in X \land b \in B\}.$ 

We follow the rules: n is a natural number,  $X, Y, Z, B, C, B_1, B_2$  are subsets of  $\mathcal{E}^n_{\mathrm{T}}$ , and x, y, p are points of  $\mathcal{E}^n_{\mathrm{T}}$ .

One can prove the following propositions:

- $(1) \quad B!! = B.$
- (2)  $\{0_{\mathcal{E}^n_T}\} + x = \{x\}.$
- (3) If  $B_1 \subseteq B_2$ , then  $B_1 + p \subseteq B_2 + p$ .
- (4) For every X such that  $X = \emptyset$  holds  $X + x = \emptyset$ .
- (5)  $X \ominus \{0_{\mathcal{E}^n_{\mathcal{T}}}\} = X.$
- $(6) \quad X \oplus \{0_{\mathcal{E}^n_{\mathcal{T}}}\} = X.$
- $(7) \quad X \oplus \{x\} = X + x.$
- (8) For all X, Y such that  $Y = \emptyset$  holds  $X \ominus Y = \mathcal{R}^n$ .
- (9) If  $X \subseteq Y$ , then  $X \ominus B \subseteq Y \ominus B$  and  $X \oplus B \subseteq Y \oplus B$ .
- (10) If  $B_1 \subseteq B_2$ , then  $X \ominus B_2 \subseteq X \ominus B_1$  and  $X \oplus B_1 \subseteq X \oplus B_2$ .
- (11) If  $0_{\mathcal{E}^n_{\mathcal{T}}} \in B$ , then  $X \ominus B \subseteq X$  and  $X \subseteq X \oplus B$ .
- (12)  $X \oplus Y = Y \oplus X.$
- (13)  $Y + y \subseteq X + x$  iff  $Y + (y x) \subseteq X$ .
- (14)  $(X+p) \ominus Y = X \ominus Y + p.$
- (15)  $(X+p) \oplus Y = X \oplus Y + p.$
- (16) (X + x) + y = X + (x + y).
- (17)  $X \ominus (Y+p) = X \ominus Y + -p.$
- (18)  $X \oplus (Y+p) = X \oplus Y + p.$
- (19) If  $x \in X$ , then  $B + x \subseteq B \oplus X$ .
- (20)  $X \subseteq (X \oplus B) \ominus B$ .
- $(21) \quad X + 0_{\mathcal{E}^n_{\mathrm{T}}} = X.$
- $(22) \quad X \ominus \{x\} = X + -x.$
- (23)  $X \ominus (Y \oplus Z) = X \ominus Y \ominus Z.$
- (24)  $X \ominus (Y \oplus Z) = X \ominus Z \ominus Y.$
- (25)  $X \oplus (Y \ominus Z) \subseteq (X \oplus Y) \ominus Z.$
- (26)  $X \oplus (Y \oplus Z) = (X \oplus Y) \oplus Z.$
- (27)  $(B \cup C) + y = (B + y) \cup (C + y).$
- (28)  $B \cap C + y = (B + y) \cap (C + y).$
- (29)  $X \ominus (B \cup C) = (X \ominus B) \cap (X \ominus C).$
- $(30) \quad X \oplus (B \cup C) = X \oplus B \cup X \oplus C.$

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- (31)  $X \ominus B \cup Y \ominus B \subseteq (X \cup Y) \ominus B$ .
- (32)  $(X \cup Y) \oplus B = X \oplus B \cup Y \oplus B.$
- $(33) \quad X \cap Y \ominus B = (X \ominus B) \cap (Y \ominus B).$
- $(34) \quad X \cap Y \oplus B \subseteq (X \oplus B) \cap (Y \oplus B).$
- $(35) \quad B \oplus X \cap Y \subseteq (B \oplus X) \cap (B \oplus Y).$
- $(36) \quad B \ominus X \cup B \ominus Y \subseteq B \ominus X \cap Y.$
- $(37) \quad (X^{c} \ominus B)^{c} = X \oplus B!.$
- $(38) \quad (X \ominus B)^{c} = X^{c} \oplus B!.$

# 2. The Definition of Adjunction Opening and Closing and Their Algebraic Properties

Let *n* be a natural number and let *X*, *B* be subsets of  $\mathcal{E}_{\mathrm{T}}^n$ . The functor  $X \bigcirc B$  yielding a subset of  $\mathcal{E}_{\mathrm{T}}^n$  is defined by:

(Def. 5)  $X \bigcirc B = (X \ominus B) \oplus B$ .

Let *n* be a natural number and let *X*, *B* be subsets of  $\mathcal{E}^n_{\mathrm{T}}$ . The functor  $X \odot B$  yielding a subset of  $\mathcal{E}^n_{\mathrm{T}}$  is defined as follows:

(Def. 6)  $X \odot B = (X \oplus B) \ominus B$ .

We now state a number of propositions:

- $(39) \quad (X^{c} \bigcirc B!)^{c} = X \odot B.$
- $(40) \quad (X^{c} \odot B!)^{c} = X \bigcirc B.$
- (41)  $X \bigcirc B \subseteq X$  and  $X \subseteq X \odot B$ .
- $(42) \quad X \bigcirc X = X.$
- (43)  $X \bigcirc B \ominus B \subseteq X \ominus B$  and  $X \bigcirc B \oplus B \subseteq X \oplus B$ .
- (44)  $X \ominus B \subseteq X \odot B \ominus B$  and  $X \oplus B \subseteq X \odot B \oplus B$ .
- (45) If  $X \subseteq Y$ , then  $X \bigcirc B \subseteq Y \bigcirc B$  and  $X \odot B \subseteq Y \odot B$ .
- (46)  $(X+p) \bigcirc Y = X \bigcirc Y + p.$
- $(47) \quad (X+p) \odot Y = X \odot Y + p.$
- (48) If  $C \subseteq B$ , then  $X \bigcirc B \subseteq (X \ominus C) \oplus B$ .
- (49) If  $B \subseteq C$ , then  $X \odot B \subseteq (X \oplus C) \ominus B$ .
- (50)  $X \oplus Y = X \odot Y \oplus Y$  and  $X \oplus Y = X \bigcirc Y \oplus Y$ .
- (51)  $X \oplus Y = (X \oplus Y) \bigcirc Y$  and  $X \oplus Y = (X \oplus Y) \odot Y$ .
- (52)  $X \bigcirc B \bigcirc B = X \bigcirc B$ .
- (53)  $X \odot B \odot B = X \odot B$ .
- (54)  $X \bigcirc B \subseteq (X \cup Y) \bigcirc B$ .
- (55) If  $B = B \bigcirc B_1$ , then  $X \bigcirc B \subseteq X \bigcirc B_1$ .

## 3. The Definition of Scaling Transformation and Its Algebraic Properties

In the sequel a is a point of  $\mathcal{E}^n_{\mathrm{T}}$ .

Let t be a real number, let us consider n, and let A be a subset of  $\mathcal{E}_{\mathrm{T}}^{n}$ . The functor  $t \odot A$  yields a subset of  $\mathcal{E}_{\mathrm{T}}^{n}$  and is defined as follows:

(Def. 7)  $t \odot A = \{t \cdot a : a \in A\}.$ 

In the sequel t, s denote real numbers.

One can prove the following propositions:

- (56) For every subset X of  $\mathcal{E}^n_{\mathrm{T}}$  such that  $X = \emptyset$  holds  $0 \odot X = \emptyset$ .
- (57) For every non empty subset X of  $\mathcal{E}^n_{\mathrm{T}}$  holds  $0 \odot X = \{0_{\mathcal{E}^n_{\mathrm{T}}}\}$ .
- $(58) \quad 1 \odot X = X.$
- (59)  $2 \odot X \subseteq X \oplus X$ .
- (60)  $(t \cdot s) \odot X = t \odot (s \odot X).$
- (61) If  $X \subseteq Y$ , then  $t \odot X \subseteq t \odot Y$ .
- (62)  $t \odot (X+x) = t \odot X + t \cdot x.$
- (63)  $t \odot (X \oplus Y) = t \odot X \oplus t \odot Y.$
- (64) If  $t \neq 0$ , then  $t \odot (X \ominus Y) = t \odot X \ominus t \odot Y$ .
- (65) If  $t \neq 0$ , then  $t \odot (X \bigcirc Y) = (t \odot X) \bigcirc (t \odot Y)$ .
- (66) If  $t \neq 0$ , then  $t \odot (X \odot Y) = (t \odot X) \odot (t \odot Y)$ .

# 4. The Definition of Thinning and Thickening and Their Algebraic Properties

Let *n* be a natural number and let *X*, *B*<sub>1</sub>, *B*<sub>2</sub> be subsets of  $\mathcal{E}^n_{\mathrm{T}}$ . The functor  $X \circledast (B_1, B_2)$  yielding a subset of  $\mathcal{E}^n_{\mathrm{T}}$  is defined as follows:

(Def. 8)  $X \circledast (B_1, B_2) = (X \ominus B_1) \cap (X^c \ominus B_2).$ 

Let *n* be a natural number and let *X*, *B*<sub>1</sub>, *B*<sub>2</sub> be subsets of  $\mathcal{E}^n_{\mathrm{T}}$ . The functor  $X \otimes (B_1, B_2)$  yields a subset of  $\mathcal{E}^n_{\mathrm{T}}$  and is defined as follows:

(Def. 9)  $X \otimes (B_1, B_2) = X \cup (X \circledast (B_1, B_2)).$ 

Let *n* be a natural number and let *X*, *B*<sub>1</sub>, *B*<sub>2</sub> be subsets of  $\mathcal{E}^n_{\mathrm{T}}$ . The functor  $X \circledast (B_1, B_2)$  yielding a subset of  $\mathcal{E}^n_{\mathrm{T}}$  is defined by:

(Def. 10)  $X \circledast (B_1, B_2) = X \setminus (X \circledast (B_1, B_2)).$ 

The following propositions are true:

- (67) If  $B_1 = \emptyset$ , then  $X \circledast (B_1, B_2) = X^c \ominus B_2$ .
- (68) If  $B_2 = \emptyset$ , then  $X \circledast (B_1, B_2) = X \ominus B_1$ .
- (69) If  $0_{\mathcal{E}^n_{\mathcal{T}}} \in B_1$ , then  $X \circledast (B_1, B_2) \subseteq X$ .
- (70) If  $0_{\mathcal{E}^n_{\mathcal{T}}} \in B_2$ , then  $(X \circledast (B_1, B_2)) \cap X = \emptyset$ .

- (71) If  $0_{\mathcal{E}^n_{\mathcal{T}}} \in B_1$ , then  $X \otimes (B_1, B_2) = X$ .
- (72) If  $0_{\mathcal{E}^n_{\mathcal{T}}} \in B_2$ , then  $X \circledast (B_1, B_2) = X$ .
- (73)  $X \otimes (B_2, B_1) = (X^c \circledast (B_1, B_2))^c.$
- (74)  $X \circledast (B_2, B_1) = (X^c \otimes (B_1, B_2))^c$ .

### 5. PROPERTIES OF EROSION, DILATION, ADJUNCTION OPENING, ADJUNCTION CLOSING ON CONVEX SETS

One can prove the following proposition

(75) Let *n* be a natural number and *B* be a subset of  $\mathcal{E}_{T}^{n}$ . Then *B* is convex if and only if for all points x, y of  $\mathcal{E}_{T}^{n}$  and for every real number *r* such that  $0 \leq r$  and  $r \leq 1$  and  $x \in B$  and  $y \in B$  holds  $r \cdot x + (1 - r) \cdot y \in B$ .

Let n be a natural number and let B be a subset of  $\mathcal{E}_{\mathrm{T}}^{n}$ . Let us observe that B is convex if and only if:

(Def. 11) For all points x, y of  $\mathcal{E}_{\mathrm{T}}^{n}$  and for every real number r such that  $0 \leq r$ and  $r \leq 1$  and  $x \in B$  and  $y \in B$  holds  $r \cdot x + (1-r) \cdot y \in B$ .

One can prove the following propositions:

- (76) If X is convex, then X! is convex.
- (77) If X is convex and B is convex, then  $X \oplus B$  is convex and  $X \oplus B$  is convex.
- (78) If X is convex and B is convex, then  $X \bigcirc B$  is convex and  $X \odot B$  is convex.
- (79) If B is convex and 0 < t and 0 < s, then  $(s + t) \odot B = s \odot B \oplus t \odot B$ .

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# Subsequences of Almost, Weakly and Poorly One-to-one Finite Sequences<sup>1</sup>

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The articles [21], [24], [1], [3], [2], [23], [4], [11], [9], [22], [16], [20], [19], [6], [7], [12], [8], [13], [17], [14], [15], [5], [18], and [10] provide the terminology and notation for this paper.

In this paper n is a natural number.

The following three propositions are true:

- (1) For every finite sequence f of elements of  $\mathcal{E}_{\mathrm{T}}^2$  and for every point p of  $\mathcal{E}_{\mathrm{T}}^2$  such that  $p \in \widetilde{\mathcal{L}}(f)$  holds len  $\downarrow p, f \ge 1$ .
- (2) For every non empty finite sequence f of elements of  $\mathcal{E}_{\mathrm{T}}^2$  and for every point p of  $\mathcal{E}_{\mathrm{T}}^2$  holds len  $| f, p \geq 1$ .
- (3) For every finite sequence f of elements of  $\mathcal{E}_{\mathrm{T}}^2$  and for all points p, q of  $\mathcal{E}_{\mathrm{T}}^2$  holds  $|\downarrow| p, f, q \neq \emptyset$ .

Let x be a set. One can check that  $\langle x \rangle$  is one-to-one.

Let f be a finite sequence. We say that f is almost one-to-one if and only if:

(Def. 1) For all natural numbers i, j such that  $i \in \text{dom } f$  and  $j \in \text{dom } f$  and  $i \neq 1$  or  $j \neq \text{len } f$  and  $i \neq \text{len } f$  or  $j \neq 1$  and f(i) = f(j) holds i = j.

Let f be a finite sequence. We say that f is weakly one-to-one if and only if:

(Def. 2) For every natural number *i* such that  $1 \le i$  and i < len f holds  $f(i) \ne f(i+1)$ .

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#### ROBERT MILEWSKI

Let f be a finite sequence. We say that f is poorly one-to-one if and only if: (Def. 3)(i) For every natural number i such that  $1 \le i$  and i < len f holds  $f(i) \ne f(i+1)$  if  $\text{len } f \ne 2$ ,

(ii) TRUE, otherwise.

The following three propositions are true:

- (4) Let D be a set and f be a finite sequence of elements of D. Then f is almost one-to-one if and only if for all natural numbers i, j such that  $i \in \text{dom } f$  and  $j \in \text{dom } f$  and  $i \neq 1$  or  $j \neq \text{len } f$  and  $i \neq \text{len } f$  or  $j \neq 1$  and  $f_i = f_j$  holds i = j.
- (5) Let D be a set and f be a finite sequence of elements of D. Then f is weakly one-to-one if and only if for every natural number i such that  $1 \le i$ and i < len f holds  $f_i \neq f_{i+1}$ .
- (6) Let D be a set and f be a finite sequence of elements of D. Then f is poorly one-to-one if and only if if len f ≠ 2, then for every natural number i such that 1 ≤ i and i < len f holds f<sub>i</sub> ≠ f<sub>i+1</sub>.

Let us note that every finite sequence which is one-to-one is also almost one-to-one.

One can check that every finite sequence which is almost one-to-one is also poorly one-to-one.

The following proposition is true

(7) For every finite sequence f such that len  $f \neq 2$  holds f is weakly one-to-one iff f is poorly one-to-one.

Let us note that  $\emptyset$  is weakly one-to-one.

Let x be a set. One can verify that  $\langle x \rangle$  is weakly one-to-one.

Let x, y be sets. Observe that  $\langle x, y \rangle$  is poorly one-to-one.

Let us mention that there exists a finite sequence which is weakly one-to-one and non empty.

Let D be a non empty set. Observe that there exists a finite sequence of elements of D which is weakly one-to-one, circular, and non empty.

We now state three propositions:

- (8) For every finite sequence f such that f is almost one-to-one holds  $\operatorname{Rev}(f)$  is almost one-to-one.
- (9) For every finite sequence f such that f is weakly one-to-one holds  $\operatorname{Rev}(f)$  is weakly one-to-one.
- (10) For every finite sequence f such that f is poorly one-to-one holds  $\operatorname{Rev}(f)$  is poorly one-to-one.

Let us observe that there exists a finite sequence which is one-to-one and non empty.

Let f be an almost one-to-one finite sequence. Observe that  $\operatorname{Rev}(f)$  is almost one-to-one.

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Let f be a weakly one-to-one finite sequence. Observe that  $\operatorname{Rev}(f)$  is weakly one-to-one.

Let f be a poorly one-to-one finite sequence. Observe that  $\operatorname{Rev}(f)$  is poorly one-to-one.

One can prove the following three propositions:

- (11) Let D be a non empty set and f be a finite sequence of elements of D. Suppose f is almost one-to-one. Let p be an element of D. Then  $f \circ p$  is almost one-to-one.
- (12) Let D be a non empty set and f be a finite sequence of elements of D. Suppose f is weakly one-to-one and circular. Let p be an element of D. Then  $f \bigcirc p$  is weakly one-to-one.
- (13) Let D be a non empty set and f be a finite sequence of elements of D. Suppose f is poorly one-to-one and circular. Let p be an element of D. Then  $f \circlearrowleft p$  is poorly one-to-one.

Let D be a non empty set. One can check that there exists a finite sequence of elements of D which is one-to-one, circular, and non empty.

Let D be a non empty set, let f be an almost one-to-one finite sequence of elements of D, and let p be an element of D. Note that  $f \circ p$  is almost one-to-one.

Let D be a non empty set, let f be a circular weakly one-to-one finite sequence of elements of D, and let p be an element of D. Note that  $f \circlearrowleft p$  is weakly one-to-one.

Let *D* be a non empty set, let *f* be a circular poorly one-to-one finite sequence of elements of *D*, and let *p* be an element of *D*. One can verify that  $f \circ p$  is poorly one-to-one.

The following proposition is true

(14) Let D be a non empty set and f be a finite sequence of elements of D. Then f is almost one-to-one if and only if  $f_{|1}$  is one-to-one and  $f \upharpoonright (\text{len } f - 1)$  is one-to-one.

Let C be a compact non vertical non horizontal subset of  $\mathcal{E}_{\mathrm{T}}^2$  and let n be a natural number. Observe that  $\mathrm{Cage}(C, n)$  is almost one-to-one.

Let C be a compact non vertical non horizontal subset of  $\mathcal{E}_{\mathrm{T}}^2$  and let n be a natural number. One can check that  $\mathrm{Cage}(C, n)$  is weakly one-to-one.

The following propositions are true:

- (15) Let f be a finite sequence of elements of  $\mathcal{E}_{\mathrm{T}}^2$  and p be a point of  $\mathcal{E}_{\mathrm{T}}^2$ . If  $p \in \widetilde{\mathcal{L}}(f)$  and f is weakly one-to-one, then  $|\downarrow p, f, p = \langle p \rangle$ .
- (16) For every finite sequence f such that f is one-to-one holds f is weakly one-to-one.

One can check that every finite sequence which is one-to-one is also weakly one-to-one.

The following propositions are true:

- (17) Let f be a finite sequence of elements of  $\mathcal{E}_{\mathrm{T}}^2$ . Suppose f is weakly one-toone. Let p, q be points of  $\mathcal{E}_{\mathrm{T}}^2$ . If  $p \in \widetilde{\mathcal{L}}(f)$  and  $q \in \widetilde{\mathcal{L}}(f)$ , then || p, f, q = $\mathrm{Rev}(|| q, f, p)$ .
- (18) Let f be a finite sequence of elements of  $\mathcal{E}_{\mathrm{T}}^2$ , p be a point of  $\mathcal{E}_{\mathrm{T}}^2$ , and  $i_1$  be a natural number. Suppose f is poorly one-to-one, unfolded, and s.n.c. and  $1 < i_1$  and  $i_1 \leq \text{len } f$  and  $p = f(i_1)$ . Then  $\text{Index}(p, f) + 1 = i_1$ .
- (19) Let f be a finite sequence of elements of  $\mathcal{E}_{\mathrm{T}}^2$ . Suppose f is weakly one-toone. Let p, q be points of  $\mathcal{E}_{\mathrm{T}}^2$ . If  $p \in \widetilde{\mathcal{L}}(f)$  and  $q \in \widetilde{\mathcal{L}}(f)$ , then  $(|\downarrow p, f, q)_1 = p$ .
- (20) Let f be a finite sequence of elements of  $\mathcal{E}_{\mathrm{T}}^2$ . Suppose f is weakly one-to-one. Let p, q be points of  $\mathcal{E}_{\mathrm{T}}^2$ . If  $p \in \mathcal{L}(f)$  and  $q \in \mathcal{\widetilde{L}}(f)$ , then  $(\downarrow \downarrow p, f, q)_{\mathrm{len} \downarrow \downarrow p, f, q} = q$ .
- (21) For every finite sequence f of elements of  $\mathcal{E}_{\mathrm{T}}^2$  and for every point p of  $\mathcal{E}_{\mathrm{T}}^2$  such that  $p \in \widetilde{\mathcal{L}}(f)$  holds  $\widetilde{\mathcal{L}}(\downarrow p, f) \subseteq \widetilde{\mathcal{L}}(f)$ .
- (22) Let f be a finite sequence of elements of  $\mathcal{E}_{\mathrm{T}}^2$  and p, q be points of  $\mathcal{E}_{\mathrm{T}}^2$ . If  $p \in \widetilde{\mathcal{L}}(f)$  and  $q \in \widetilde{\mathcal{L}}(f)$  and f is weakly one-to-one, then  $\widetilde{\mathcal{L}}(|\!\!|\, p, f, q) \subseteq \widetilde{\mathcal{L}}(f)$ .
- (23) For all finite sequences f, g holds dom  $f \subseteq \text{dom}(f \frown g)$ .
- (24) For every non empty finite sequence f and for every finite sequence g holds dom  $g \subseteq \text{dom}(f \frown g)$ .
- (25) For all finite sequences f, g such that  $f \frown g$  is constant holds f is constant.
- (26) For all finite sequences f, g such that  $f \frown g$  is constant and  $f(\operatorname{len} f) = g(1)$  and  $f \neq \emptyset$  holds g is constant.
- (27) For every special finite sequence f of elements of  $\mathcal{E}_{\mathrm{T}}^2$  and for all natural numbers i, j holds  $\operatorname{mid}(f, i, j)$  is special.
- (28) For every unfolded finite sequence f of elements of  $\mathcal{E}_{\mathrm{T}}^2$  and for all natural numbers i, j holds  $\operatorname{mid}(f, i, j)$  is unfolded.
- (29) Let f be a finite sequence of elements of  $\mathcal{E}_{\mathrm{T}}^2$ . Suppose f is special. Let p be a point of  $\mathcal{E}_{\mathrm{T}}^2$ . If  $p \in \widetilde{\mathcal{L}}(f)$ , then | p, f is special.
- (30) Let f be a finite sequence of elements of  $\mathcal{E}_{\mathrm{T}}^2$ . Suppose f is special. Let p be a point of  $\mathcal{E}_{\mathrm{T}}^2$ . If  $p \in \widetilde{\mathcal{L}}(f)$ , then | f, p is special.
- (31) Let f be a finite sequence of elements of  $\mathcal{E}^2_{\mathrm{T}}$ . Suppose f is special and weakly one-to-one. Let p, q be points of  $\mathcal{E}^2_{\mathrm{T}}$ . If  $p \in \widetilde{\mathcal{L}}(f)$  and  $q \in \widetilde{\mathcal{L}}(f)$ , then || p, f, q is special.
- (32) Let f be a finite sequence of elements of  $\mathcal{E}_{\mathrm{T}}^2$ . Suppose f is unfolded. Let p be a point of  $\mathcal{E}_{\mathrm{T}}^2$ . If  $p \in \widetilde{\mathcal{L}}(f)$ , then  $\downarrow p, f$  is unfolded.
- (33) Let f be a finite sequence of elements of  $\mathcal{E}_{T}^{2}$ . Suppose f is unfolded. Let

p be a point of  $\mathcal{E}^2_{\mathrm{T}}$ . If  $p \in \mathcal{L}(f)$ , then |f, p is unfolded.

- (34) Let f be a finite sequence of elements of  $\mathcal{E}^2_{\mathrm{T}}$ . Suppose f is unfolded and weakly one-to-one. Let p, q be points of  $\mathcal{E}^2_{\mathrm{T}}$ . If  $p \in \widetilde{\mathcal{L}}(f)$  and  $q \in \widetilde{\mathcal{L}}(f)$ , then  $\downarrow \downarrow p, f, q$  is unfolded.
- (35) Let f, g be finite sequences of elements of  $\mathcal{E}_{\mathrm{T}}^2$  and p be a point of  $\mathcal{E}_{\mathrm{T}}^2$ . Suppose f is almost one-to-one, special, unfolded, and s.n.c. and  $p \in \widetilde{\mathcal{L}}(f)$ and  $p \neq f(1)$  and  $g = (\mathrm{mid}(f, 1, \mathrm{Index}(p, f))) \cap \langle p \rangle$ . Then g is a special sequence joining  $f_1, p$ .
- (36) Let f be a finite sequence of elements of  $\mathcal{E}_{\mathrm{T}}^2$  and p be a point of  $\mathcal{E}_{\mathrm{T}}^2$ . Suppose f is poorly one-to-one, unfolded, and s.n.c. and  $p \in \widetilde{\mathcal{L}}(f)$  and  $p = f(\mathrm{Index}(p, f) + 1)$  and  $p \neq f(\mathrm{len} f)$ . Then  $\mathrm{Index}(p, \mathrm{Rev}(f)) + \mathrm{Index}(p, f) + 1 = \mathrm{len} f$ .
- (37) Let f be a non empty finite sequence of elements of  $\mathcal{E}_{\mathrm{T}}^2$  and p be a point of  $\mathcal{E}_{\mathrm{T}}^2$ . If f is weakly one-to-one and len  $f \geq 2$ , then  $\downarrow f_1, f = f$ .
- (38) Let f be a non empty finite sequence of elements of  $\mathcal{E}_{\mathrm{T}}^2$  and p be a point of  $\mathcal{E}_{\mathrm{T}}^2$ . Suppose f is poorly one-to-one, unfolded, and s.n.c. and  $p \in \widetilde{\mathcal{L}}(f)$  and  $p \neq f(\operatorname{len} f)$ . Then  $\downarrow p, \operatorname{Rev}(f) = \operatorname{Rev}(\downarrow f, p)$ .
- (39) Let f be a finite sequence of elements of  $\mathcal{E}_{\mathrm{T}}^2$  and p be a point of  $\mathcal{E}_{\mathrm{T}}^2$ . Suppose f is almost one-to-one, special, unfolded, and s.n.c. and  $p \in \widetilde{\mathcal{L}}(f)$  and  $p \neq f(1)$ . Then |f, p is a special sequence joining  $f_1, p$ .
- (40) Let f be a non empty finite sequence of elements of  $\mathcal{E}_{\mathrm{T}}^2$  and p be a point of  $\mathcal{E}_{\mathrm{T}}^2$ . Suppose f is almost one-to-one, special, unfolded, and s.n.c. and  $p \in \mathcal{L}(f)$  and  $p \neq f(\operatorname{len} f)$  and  $p \neq f(1)$ . Then | p, f is a special sequence joining p,  $f_{\operatorname{len} f}$ .
- (41) Let f be a finite sequence of elements of  $\mathcal{E}_{\mathrm{T}}^2$  and p be a point of  $\mathcal{E}_{\mathrm{T}}^2$ . Suppose f is almost one-to-one, special, unfolded, and s.n.c. and  $p \in \widetilde{\mathcal{L}}(f)$  and  $p \neq f(1)$ . Then |f, p is a special sequence.
- (42) Let f be a non empty finite sequence of elements of  $\mathcal{E}_{\mathrm{T}}^2$  and p be a point of  $\mathcal{E}_{\mathrm{T}}^2$ . Suppose f is almost one-to-one, special, unfolded, and s.n.c. and  $p \in \mathcal{L}(f)$  and  $p \neq f(\ln f)$  and  $p \neq f(1)$ . Then  $\downarrow p, f$  is a special sequence.
- (43) Let f be a non empty finite sequence of elements of  $\mathcal{E}_{\mathrm{T}}^2$  and p, q be points of  $\mathcal{E}_{\mathrm{T}}^2$ . Suppose that f is almost one-to-one, special, unfolded, and s.n.c. and len  $f \neq 2$  and  $p \in \widetilde{\mathcal{L}}(f)$  and  $q \in \widetilde{\mathcal{L}}(f)$  and  $p \neq q$  and  $p \neq f(1)$  and  $q \neq f(1)$ . Then || p, f, q is a special sequence joining p, q.
- (44) Let f be a non empty finite sequence of elements of  $\mathcal{E}_{\mathrm{T}}^2$  and p, q be points of  $\mathcal{E}_{\mathrm{T}}^2$ . Suppose that f is almost one-to-one, special, unfolded, and s.n.c. and len  $f \neq 2$  and  $p \in \widetilde{\mathcal{L}}(f)$  and  $q \in \widetilde{\mathcal{L}}(f)$  and  $p \neq q$  and  $p \neq f(1)$  and  $q \neq f(1)$ . Then || p, f, q is a special sequence.
- (45) Let C be a compact non vertical non horizontal subset of  $\mathcal{E}_{\mathrm{T}}^2$  and p, q be points of  $\mathcal{E}_{\mathrm{T}}^2$ . Suppose  $p \in \mathrm{BDD}\,\widetilde{\mathcal{L}}(\mathrm{Cage}(C,n))$ . Then there exists a

S-sequence B in  $\mathbb{R}^2$  such that

- (i)  $B = \bigcup \text{South-Bound}(p, \widetilde{\mathcal{L}}(\text{Cage}(C, n))),$   $(\text{Cage}(C, n) \circlearrowleft (\text{Cage}(C, n))_{\text{Index}(\text{South-Bound}(p, \widetilde{\mathcal{L}}(\text{Cage}(C, n))), \text{Cage}(C, n))}) \upharpoonright (\text{len}$   $(\text{Cage}(C, n) \circlearrowright (\text{Cage}(C, n))_{\text{Index}(\text{South-Bound}(p, \widetilde{\mathcal{L}}(\text{Cage}(C, n))), \text{Cage}(C, n))}) - 1),$ North-Bound $(p, \widetilde{\mathcal{L}}(\text{Cage}(C, n))), \text{ and}$
- (ii) there exists a S-sequence P in  $\mathbb{R}^2$  such that P is a sequence which elements belong to the Go-board of  $B \curvearrowright \langle \text{North-Bound}(p, \widetilde{\mathcal{L}}(\text{Cage}(C, n))), \text{South-Bound}(p, \widetilde{\mathcal{L}}(\text{Cage}(C, n))) \rangle$  and  $\widetilde{\mathcal{L}}(\langle \text{North-Bound}(p, \widetilde{\mathcal{L}}(\text{Cage}(C, n))), \text{South-Bound}(p, \widetilde{\mathcal{L}}(\text{Cage}(C, n))) \rangle = \widetilde{\mathcal{L}}(P)$  and

$$P_1 = \text{North-Bound}(p, \mathcal{L}(\text{Cage}(C, n)))$$
 and

 $P_{\text{len }P} = \text{South-Bound}(p, \widetilde{\mathcal{L}}(\text{Cage}(C, n))) \text{ and } \text{len }P \geq 2 \text{ and there exists a S-sequence } B_1 \text{ in } \mathbb{R}^2 \text{ such that } B_1 \text{ is a sequence which elements belong to the Go-board of } B \curvearrowleft \langle \text{North-Bound}(p, \widetilde{\mathcal{L}}(\text{Cage}(C, n))), \text{South-Bound}(p, \widetilde{\mathcal{L}}(\text{Cage}(C, n))) \rangle \text{ and } \widetilde{\mathcal{L}}(B) = \widetilde{\mathcal{L}}(B_1) \text{ and } B_1 = (B_1)_1 \text{ and } B_{\text{len }B} = (B_1)_{\text{len }B_1} \text{ and } \text{len } B \leq \text{len }B_1 \text{ and there exists a non constant standard special circular sequence } g \text{ such that } g = B_1 \frown P.$ 

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## Alternative Graph Structures<sup>1</sup>

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**Summary.** We define the notion of a graph anew without using the available Mizar structures. In our approach, we model graph structure as a finite function whose domain is a subset of natural numbers. The elements of the domain of the function play the role of selectors for accessing the components of the structure. As these selectors are first class objects, many future extensions of the new graph structure turned out to be easier to formalize in Mizar than with the traditional Mizar structures.

After introducing graph structure, we define its selectors and then conditions that the structure needs to satisfy to form a directed graph (in the spirit of [13]). For these graphs we define a collection of basic graph notions; the presentation of these notions is continued in articles [16, 15, 17].

We have tried to follow a number of graph theory books in choosing graph terminology but since the terminology is not commonly agreed upon, we had to make a number of compromises, see [14].

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The papers [20], [19], [22], [21], [24], [2], [1], [25], [7], [5], [12], [3], [8], [6], [23], [9], [4], [10], [11], and [18] provide the terminology and notation for this paper.

#### 1. Definitions

A finite function is called a graph structure if:

(Def. 1) dom it  $\subseteq \mathbb{N}$ .

The natural number VertexSelector is defined as follows:

(Def. 2) VertexSelector = 1.

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The natural number EdgeSelector is defined as follows:

(Def. 3) EdgeSelector = 2.

The natural number SourceSelector is defined by:

(Def. 4) SourceSelector = 3.

The natural number TargetSelector is defined by:

- (Def. 5) TargetSelector = 4.
  - The non empty subset the graph selectors of  $\mathbb{N}$  is defined by:
- (Def. 6) The graph selectors =

{VertexSelector, EdgeSelector, SourceSelector, TargetSelector}.

Let G be a graph structure. The vertices of G is defined by:

(Def. 7) The vertices of G = G(VertexSelector).

The edges of G is defined by:

- (Def. 8) The edges of G = G(EdgeSelector). The source of G is defined by:
- (Def. 9) The source of G = G(SourceSelector). The target of G is defined by:
- (Def. 10) The target of G = G(TargetSelector).

Let G be a graph structure. We say that G is graph-like if and only if the conditions (Def. 11) are satisfied.

(Def. 11) VertexSelector  $\in$  dom G and EdgeSelector  $\in$  dom G and SourceSelector  $\in$  dom G and TargetSelector  $\in$  dom G and the vertices of G is a non empty set and the source of G is a function from the edges of G into the vertices of G and the target of G is a function from the edges of G into the vertices of G.

Let us note that there exists a graph structure which is graph-like.

A graph is a graph-like graph structure.

Let G be a graph. Observe that the vertices of G is non empty.

Let G be a graph. Then the source of G is a function from the edges of G into the vertices of G. Then the target of G is a function from the edges of G into the vertices of G.

Let V be a non empty set, let E be a set, and let S, T be functions from E into V. The functor createGraph(V, E, S, T) yielding a graph is defined by:

(Def. 12) createGraph $(V, E, S, T) = \langle V, E, S, T \rangle$ .

Let x, y be sets. One can verify that  $x \mapsto y$  is finite.

Let G be a graph structure, let n be a natural number, and let x be a set. The functor G.set(n, x) yielding a graph structure is defined as follows:

(Def. 13)  $G.set(n, x) = G + (n \mapsto x).$ 

Let G be a graph structure and let X be a set. The functor G.strict(X) yielding a graph structure is defined by:

(Def. 14)  $G.strict(X) = G \upharpoonright X.$ 

Let G be a graph. Observe that G.strict(the graph selectors) is graph-like. Let G be a graph and let x, y, e be sets. We say that e joins x and y in G if and only if the conditions (Def. 15) are satisfied.

- (Def. 15)(i)  $e \in$  the edges of G, and
  - (ii) (the source of G)(e) = x and (the target of G)(e) = y or (the source of G)(e) = y and (the target of G)(e) = x.

Let G be a graph and let x, y, e be sets. We say that e joins x to y in G if and only if:

(Def. 16)  $e \in$  the edges of G and (the source of G)(e) = x and (the target of G)(e) = y.

Let G be a graph and let X, Y, e be sets. We say that e joins a vertex from X and a vertex from Y in G if and only if the conditions (Def. 17) are satisfied.

(Def. 17)(i)  $e \in$  the edges of G, and

(ii) (the source of G) $(e) \in X$  and (the target of G) $(e) \in Y$  or (the source of G) $(e) \in Y$  and (the target of G) $(e) \in X$ .

We say that e joins a vertex from X to a vertex from Y in G if and only if:

(Def. 18)  $e \in$  the edges of G and (the source of G) $(e) \in X$  and (the target of G) $(e) \in Y$ .

Let G be a graph. We say that G is finite if and only if:

- (Def. 19) The vertices of G is finite and the edges of G is finite. We say that G is loopless if and only if:
- (Def. 20) It is not true that there exists a set e such that  $e \in$  the edges of G and (the source of G)(e) = (the target of G)(e).

We say that G is trivial if and only if:

(Def. 21) The vertices of  $\overline{G} = \mathbf{1}$ .

We say that G is non-multi if and only if:

(Def. 22) For all sets  $e_1$ ,  $e_2$ ,  $v_1$ ,  $v_2$  such that  $e_1$  joins  $v_1$  and  $v_2$  in G and  $e_2$  joins  $v_1$  and  $v_2$  in G holds  $e_1 = e_2$ .

We say that G is non-directed-multi if and only if:

(Def. 23) For all sets  $e_1$ ,  $e_2$ ,  $v_1$ ,  $v_2$  such that  $e_1$  joins  $v_1$  to  $v_2$  in G and  $e_2$  joins  $v_1$  to  $v_2$  in G holds  $e_1 = e_2$ .

Let G be a graph. We say that G is simple if and only if:

(Def. 24) G is loopless and non-multi.

We say that G is directed-simple if and only if:

(Def. 25) G is loopless and non-directed-multi.

One can verify the following observations:

\* every graph which is non-multi is also non-directed-multi,

- \* every graph which is simple is also loopless and non-multi,
- \* every graph which is loopless and non-multi is also simple,
- \* every graph which is loopless and non-directed-multi is also directed-simple,
- \* every graph which is directed-simple is also loopless and non-directedmulti,
- \* every graph which is trivial and loopless is also finite, and
- \* every graph which is trivial and non-directed-multi is also finite.

Let us note that there exists a graph which is trivial and simple and there exists a graph which is finite, non trivial, and simple.

Let G be a finite graph. Observe that the vertices of G is finite and the edges of G is finite.

Let G be a trivial graph. One can verify that the vertices of G is finite.

Let V be a non empty finite set, let E be a finite set, and let S, T be functions from E into V. One can check that createGraph(V, E, S, T) is finite.

Let V be a non empty set, let E be an empty set, and let S, T be functions from E into V. One can check that createGraph(V, E, S, T) is simple.

Let v be a set, let E be a set, and let S, T be functions from E into  $\{v\}$ . Observe that createGraph $(\{v\}, E, S, T)$  is trivial.

Let G be a graph. The functor G.order() yielding a cardinal number is defined as follows:

(Def. 26)  $G.order() = \overline{the vertices of G}.$ 

Let G be a finite graph. Then G.order() is a non empty natural number. Let G be a graph. The functor G.size() yields a cardinal number and is defined by:

(Def. 27)  $G.size() = \overline{the edges of G}.$ 

Let G be a finite graph. Then G.size() is a natural number.

Let G be a graph and let X be a set. The functor G.edgesInto(X) yields a subset of the edges of G and is defined as follows:

(Def. 28) For every set e holds  $e \in G$ .edgesInto(X) iff  $e \in$  the edges of G and (the target of G) $(e) \in X$ .

The functor G.edgesOutOf(X) yields a subset of the edges of G and is defined by:

(Def. 29) For every set e holds  $e \in G$ .edgesOutOf(X) iff  $e \in$  the edges of G and (the source of G) $(e) \in X$ .

Let G be a graph and let X be a set. The functor G.edgesInOut(X) yields a subset of the edges of G and is defined by:

(Def. 30)  $G.edgesInOut(X) = G.edgesInto(X) \cup G.edgesOutOf(X).$ 

The functor G.edgesBetween(X) yielding a subset of the edges of G is defined as follows:

(Def. 31)  $G.edgesBetween(X) = G.edgesInto(X) \cap G.edgesOutOf(X).$ 

Let G be a graph and let X, Y be sets. The functor G.edgesBetween(X, Y) yielding a subset of the edges of G is defined by:

(Def. 32) For every set e holds  $e \in G$ .edgesBetween(X, Y) iff e joins a vertex from X and a vertex from Y in G.

The functor G.edgesDBetween(X, Y) yields a subset of the edges of G and is defined as follows:

(Def. 33) For every set e holds  $e \in G$ .edgesDBetween(X, Y) iff e joins a vertex from X to a vertex from Y in G.

In this article we present several logical schemes. The scheme FinGraphOrder-Ind concerns a unary predicate  $\mathcal{P}$ , and states that:

For every finite graph G holds  $\mathcal{P}[G]$ 

provided the following conditions are met:

- For every finite graph G such that G.order() = 1 holds  $\mathcal{P}[G]$ , and
- Let k be a non empty natural number. Suppose that for every finite graph  $G_1$  such that  $G_1$ .order() = k holds  $\mathcal{P}[G_1]$ . Let  $G_2$  be a finite graph. If  $G_2$ .order() = k + 1, then  $\mathcal{P}[G_2]$ .

The scheme FinGraphSizeInd concerns a unary predicate  $\mathcal{P}$ , and states that: For every finite graph G holds  $\mathcal{P}[G]$ 

provided the following requirements are met:

- For every finite graph G such that G.size() = 0 holds  $\mathcal{P}[G]$ , and
- Let k be a natural number. Suppose that for every finite graph  $G_1$  such that  $G_1$ .size() = k holds  $\mathcal{P}[G_1]$ . Let  $G_2$  be a finite graph. If  $G_2$ .size() = k + 1, then  $\mathcal{P}[G_2]$ .

Let G be a graph. A graph is called a subgraph of G if it satisfies the conditions (Def. 34).

(Def. 34)(i) The vertices of it  $\subseteq$  the vertices of G,

- (ii) the edges of it  $\subseteq$  the edges of G, and
- (iii) for every set e such that  $e \in$  the edges of it holds (the source of it)(e) = (the source of G)(e) and (the target of it)(e) = (the target of G)(e).

Let  $G_3$  be a graph and let  $G_4$  be a subgraph of  $G_3$ . Then the vertices of  $G_4$  is a non empty subset of the vertices of  $G_3$ . Then the edges of  $G_4$  is a subset of the edges of  $G_3$ .

Let G be a graph. Note that there exists a subgraph of G which is trivial and simple.

Let G be a finite graph. Note that every subgraph of G is finite.

Let G be a loopless graph. Observe that every subgraph of G is loopless.

Let G be a trivial graph. One can check that every subgraph of G is trivial.

Let G be a non-multi graph. Observe that every subgraph of G is non-multi.

Let  $G_3$  be a graph and let  $G_4$  be a subgraph of  $G_3$ . We say that  $G_4$  is spanning if and only if:

(Def. 35) The vertices of  $G_4$  = the vertices of  $G_3$ .

Let G be a graph. One can verify that there exists a subgraph of G which is spanning.

Let  $G_3$ ,  $G_4$  be graphs. The predicate  $G_3 =_G G_4$  is defined by the conditions (Def. 36).

(Def. 36)(i) The vertices of  $G_3$  = the vertices of  $G_4$ ,

(ii) the edges of  $G_3$  = the edges of  $G_4$ ,

(iii) the source of  $G_3$  = the source of  $G_4$ , and

(iv) the target of  $G_3$  = the target of  $G_4$ .

Let us notice that the predicate  $G_3 =_G G_4$  is reflexive and symmetric. Let  $G_3$ ,  $G_4$  be graphs. We introduce  $G_3 \neq_G G_4$  as an antonym of  $G_3 =_G G_4$ . Let  $G_3$ ,  $G_4$  be graphs. The predicate  $G_3 \subseteq G_4$  is defined as follows:

(Def. 37)  $G_3$  is a subgraph of  $G_4$ .

Let us note that the predicate  $G_3 \subseteq G_4$  is reflexive.

Let  $G_3$ ,  $G_4$  be graphs. The predicate  $G_3 \subset G_4$  is defined as follows:

(Def. 38) 
$$G_3 \subseteq G_4$$
 and  $G_3 \neq_G G_4$ .

Let us note that the predicate  $G_3 \subset G_4$  is irreflexive.

Let G be a graph and let V, E be sets. A subgraph of G is called a subgraph of G induced by V and E if:

(Def. 39)(i) The vertices of it = V and the edges of it = E if V is a non empty subset of the vertices of G and  $E \subseteq G.edgesBetween(V)$ ,

(ii) it  $=_G G$ , otherwise.

Let G be a graph and let V be a set. A subgraph of G induced by V is a subgraph of G induced by V and G.edgesBetween(V).

Let G be a graph, let V be a finite non empty subset of the vertices of G, and let E be a finite subset of G.edgesBetween(V). Observe that every subgraph of G induced by V and E is finite.

Let G be a graph, let v be an element of the vertices of G, and let E be a subset of G.edgesBetween( $\{v\}$ ). Note that every subgraph of G induced by  $\{v\}$  and E is trivial.

Let G be a graph and let v be an element of the vertices of G. Note that every subgraph of G induced by  $\{v\}$  and  $\emptyset$  is finite and trivial.

Let G be a graph and let V be a non empty subset of the vertices of G. Note that every subgraph of G induced by V and  $\emptyset$  is simple.

Let G be a graph and let E be a subset of the edges of G. Observe that every subgraph of G induced by the vertices of G and E is spanning.

Let G be a graph. One can check that every subgraph of G induced by the vertices of G and  $\emptyset$  is spanning.

Let G be a graph and let v be a set. A subgraph of G with vertex v removed is a subgraph of G induced by (the vertices of  $G \setminus \{v\}$ .

Let G be a graph and let V be a set. A subgraph of G with vertices V removed is a subgraph of G induced by (the vertices of  $G \setminus V$ .

Let G be a graph and let e be a set. A subgraph of G with edge e removed is a subgraph of G induced by the vertices of G and (the edges of G)  $\setminus \{e\}$ .

Let G be a graph and let E be a set. A subgraph of G with edges E removed is a subgraph of G induced by the vertices of G and (the edges of  $G \setminus E$ .

Let G be a graph and let e be a set. Observe that every subgraph of G with edge e removed is spanning.

Let G be a graph and let E be a set. Observe that every subgraph of G with edges E removed is spanning.

Let G be a graph. A vertex of G is an element of the vertices of G.

Let G be a graph and let v be a vertex of G. The functor v.edgesIn() yielding a subset of the edges of G is defined as follows:

(Def. 40)  $v.edgesIn() = G.edgesInto(\{v\}).$ 

The functor v.edgesOut() yields a subset of the edges of G and is defined as follows:

(Def. 41)  $v.edgesOut() = G.edgesOutOf(\{v\}).$ 

The functor v.edgesInOut() yields a subset of the edges of G and is defined by: (Def. 42) v.edgesInOut() = G.edgesInOut( $\{v\}$ ).

Let G be a graph, let v be a vertex of G, and let e be a set. The functor v.adj(e) yields a vertex of G and is defined by:

$$(\text{the source of } G)(e), \text{ if } e \in \text{the edges of } G \text{ and} \\ (\text{the target of } G)(e) = v, \\ (\text{the target of } G)(e) = if e \in \text{the edges of } G \text{ and} \\ (\text{the target of } G)(e) = if e \in \text{the edges of } G \text{ and} \\ (\text{the target of } G)(e) = if e \in \text{the edges of } G \text{ and} \\ (\text{the target of } G)(e) = if e \in \text{the edges of } G \text{ and} \\ (\text{the target of } G)(e) = if e \in \text{the edges of } G \text{ and} \\ (\text{the target of } G)(e) = if e \in \text{the edges of } G \text{ and} \\ (\text{the target of } G)(e) = if e \in \text{the edges of } G \text{ and} \\ (\text{the target of } G)(e) = if e \in \text{the edges of } G \text{ and} \\ (\text{the target of } G)(e) = if e \in \text{the edges of } G \text{ and} \\ (\text{the target of } G)(e) = if e \in \text{the edges of } G \text{ and} \\ (\text{the target of } G)(e) = if e \in \text{the edges of } G \text{ and} \\ (\text{the target of } G)(e) = if e \in \text{the edges of } G \text{ and} \\ (\text{the target of } G)(e) = if e \in \text{the edges of } G \text{ and} \\ (\text{the target of } G)(e) = if e \in \text{the edges of } G \text{ and} \\ (\text{the target of } G)(e) = if e \in \text{the edges of } G \text{ and} \\ (\text{the target of } G)(e) = if e \in \text{the edges of } G \text{ and} \\ (\text{the target of } G)(e) = if e \in \text{the edges of } G \text{ and} \\ (\text{the target of } G)(e) = if e \in \text{the edges of } G \text{ and} \\ (\text{the target of } G)(e) = if e \in \text{the edges of } G \text{ and} \\ (\text{the target of } G)(e) = if e \in \text{the edges of } G \text{ and} \\ (\text{the target of } G)(e) = if e \in \text{the edges of } G \text{ and} \\ (e) = if e \in \text{the edges of } G \text{ and} \\ (e) = if e \in \text{the edges of } G \text{ and} \\ (e) = if e \in \text{the edges of } G \text{ and} \\ (e) = if e \in \text{the edges of } G \text{ and} \\ (e) = if e \in \text{the edges of } G \text{ and} \\ (e) = if e \in \text{the edges of } G \text{ and} \\ (e) = if e \in \text{the edges of } G \text{ and} \\ (e) = if e \in \text{the edges of } G \text{ and} \\ (e) = if e \in \text{the edges of } G \text{ and} \\ (e) = if e \in \text{the edges of } G \text{ and} \\ (e) = if e \in \text{the edges of } G \text{ and} \\ (e) = if e \in \text{the edges of } G \text{ and} \\ (e) = if e \in \text{the edges of } G \text{ and} \\ (e) = if e \in \text{the edges of } G \text{ and} \\ (e) = if e \in \text{the edges of } G \text$$

(Def. 43)  $v.adj(e) = \begin{cases} \text{(the target of } G)(e), \text{ if } e \in \text{the edges of } G \text{ and} \\ (\text{the source of } G)(e) = v \text{ and (the target of } G)(e) \neq v, \\ v, \text{ otherwise.} \end{cases}$ 

Let G be a graph and let v be a vertex of G. The functor v.inDegree() yields a cardinal number and is defined as follows:

(Def. 44)  $v.inDegree() = \overline{v.edgesIn()}.$ 

The functor v.outDegree() yielding a cardinal number is defined as follows:

(Def. 45)  $v.outDegree() = \overline{v.edgesOut()}$ .

Let G be a finite graph and let v be a vertex of G. Then v.inDegree() is a natural number. Then v.outDegree() is a natural number.

Let G be a graph and let v be a vertex of G. The functor v.degree() yielding a cardinal number is defined as follows:

(Def. 46) v.degree() = v.inDegree() + v.outDegree().

Let G be a finite graph and let v be a vertex of G. Then v.degree() is a natural number and it can be characterized by the condition:

(Def. 47) v.degree() = v.inDegree() + v.outDegree().

Let G be a graph and let v be a vertex of G. The functor v.inNeighbors() yields a subset of the vertices of G and is defined as follows:

(Def. 48)  $v.inNeighbors() = (the source of G)^{\circ}v.edgesIn().$ 

The functor v.outNeighbors() yielding a subset of the vertices of G is defined by:

(Def. 49)  $v.outNeighbors() = (the target of G)^{\circ}v.edgesOut().$ 

Let G be a graph and let v be a vertex of G. The functor v.allNeighbors() yields a subset of the vertices of G and is defined by:

(Def. 50)  $v.allNeighbors() = v.inNeighbors() \cup v.outNeighbors().$ 

Let G be a graph and let v be a vertex of G. We say that v is isolated if and only if:

(Def. 51)  $v.edgesInOut() = \emptyset$ .

Let G be a finite graph and let v be a vertex of G. Let us observe that v is isolated if and only if:

(Def. 52) v.degree() = 0.

Let G be a graph and let v be a vertex of G. We say that v is endvertex if and only if:

(Def. 53) There exists a set e such that  $v.edgesInOut() = \{e\}$  and e does not join v and v in G.

Let G be a finite graph and let v be a vertex of G. Let us observe that v is endvertex if and only if:

 $(Def. 54) \quad v.degree() = 1.$ 

Let F be a many sorted set indexed by N. We say that F is graph-yielding if and only if:

(Def. 55) For every natural number n holds F(n) is a graph.

We say that F is halting if and only if:

(Def. 56) There exists a natural number n such that F(n) = F(n+1).

Let F be a many sorted set indexed by N. The functor F.Lifespan() yielding a natural number is defined by:

(Def. 57)(i) F(F.Lifespan()) = F(F.Lifespan()+1) and for every natural number n such that F(n) = F(n+1) holds  $F.Lifespan() \le n$  if F is halting,

(ii) F.Lifespan() = 0, otherwise.

Let F be a many sorted set indexed by N. The functor F.Result() yielding a set is defined by:

(Def. 58) F.Result() = F(F.Lifespan()).

Let us mention that there exists a many sorted set indexed by  $\mathbb N$  which is graph-yielding.

A graph sequence is a graph-yielding many sorted set indexed by  $\mathbb{N}$ .

Let  $G_5$  be a graph sequence and let x be a natural number. The functor  $G_5 \rightarrow x$  yields a graph and is defined by:

(Def. 59)  $G_5 \to x = G_5(x)$ .

Let  $G_5$  be a graph sequence. We say that  $G_5$  is finite if and only if:

(Def. 60) For every natural number x holds  $G_5 \rightarrow x$  is finite.

We say that  $G_5$  is loopless if and only if:

- (Def. 61) For every natural number x holds  $G_5 \rightarrow x$  is loopless. We say that  $G_5$  is trivial if and only if:
- (Def. 62) For every natural number x holds  $G_5 \rightarrow x$  is trivial. We say that  $G_5$  is non-trivial if and only if:
- (Def. 63) For every natural number x holds  $G_5 \rightarrow x$  is non trivial. We say that  $G_5$  is non-multi if and only if:
- (Def. 64) For every natural number x holds  $G_5 \rightarrow x$  is non-multi. We say that  $G_5$  is non-directed-multi if and only if:
- (Def. 65) For every natural number x holds  $G_5 \rightarrow x$  is non-directed-multi. We say that  $G_5$  is simple if and only if:
- (Def. 66) For every natural number x holds  $G_5 \rightarrow x$  is simple. We say that  $G_5$  is directed-simple if and only if:
- (Def. 67) For every natural number x holds  $G_5 \rightarrow x$  is directed-simple.
  - Let  $G_5$  be a graph sequence. Let us observe that  $G_5$  is halting if and only if:

(Def. 68) There exists a natural number n such that  $G_5 \rightarrow n = G_5 \rightarrow (n+1)$ .

One can verify that there exists a graph sequence which is halting, finite, loopless, trivial, non-multi, non-directed-multi, simple, and directed-simple and there exists a graph sequence which is halting, finite, loopless, non-trivial, nonmulti, non-directed-multi, simple, and directed-simple.

Let  $G_5$  be a finite graph sequence and let x be a natural number. One can check that  $G_5 \rightarrow x$  is finite.

Let  $G_5$  be a loopless graph sequence and let x be a natural number. Note that  $G_5 \rightarrow x$  is loopless.

Let  $G_5$  be a trivial graph sequence and let x be a natural number. Observe that  $G_5 \rightarrow x$  is trivial.

Let  $G_5$  be a non-trivial graph sequence and let x be a natural number. Observe that  $G_5 \rightarrow x$  is non trivial.

Let  $G_5$  be a non-multi graph sequence and let x be a natural number. Note that  $G_5 \rightarrow x$  is non-multi.

Let  $G_5$  be a non-directed-multi graph sequence and let x be a natural number. Observe that  $G_5 \rightarrow x$  is non-directed-multi.

Let  $G_5$  be a simple graph sequence and let x be a natural number. Note that  $G_5 \rightarrow x$  is simple.

Let  $G_5$  be a directed-simple graph sequence and let x be a natural number. Note that  $G_5 \rightarrow x$  is directed-simple.

One can check that every graph sequence which is non-multi is also nondirected-multi.

Let us observe that every graph sequence which is simple is also loopless and non-multi.

One can verify that every graph sequence which is loopless and non-multi is also simple.

Let us note that every graph sequence which is loopless and non-directedmulti is also directed-simple.

One can verify that every graph sequence which is directed-simple is also loopless and non-directed-multi.

Let us note that every graph sequence which is trivial and loopless is also finite.

Let us observe that every graph sequence which is trivial and non-directedmulti is also finite.

## 2. Theorems

For simplicity, we adopt the following convention:  $G_6$  denotes a graph structure, G,  $G_3$ ,  $G_4$ ,  $G_7$  denote graphs, e, x,  $x_1$ ,  $x_2$ , y,  $y_1$ ,  $y_2$ , E, V, X, Y denote sets, n,  $n_1$ ,  $n_2$  denote natural numbers, and v,  $v_1$ ,  $v_2$  denote vertices of G.

We now state a number of propositions:

- (1) VertexSelector = 1 and EdgeSelector = 2 and SourceSelector = 3 and TargetSelector = 4.
- (2)  $x \in$  the graph selectors iff x = VertexSelector or x = EdgeSelector or x = SourceSelector or x = TargetSelector.
- (3) The graph selectors  $\subseteq \operatorname{dom} G$ .
- (4) The vertices of  $G_6 = G_6$  (VertexSelector) and the edges of  $G_6 = G_6$  (EdgeSelector) and the source of  $G_6 = G_6$  (SourceSelector) and the target of  $G_6 = G_6$  (TargetSelector).

(5)(i) dom (the source of G) = the edges of G,

- (ii) dom (the target of G) = the edges of G,
- (iii) rng (the source of G)  $\subseteq$  the vertices of G, and
- (iv) rng (the target of G)  $\subseteq$  the vertices of G.
- $(7)^3$   $G_6$  is graph-like if and only if the following conditions are satisfied:
- (i) the graph selectors  $\subseteq \operatorname{dom} G_6$ ,

<sup>&</sup>lt;sup>3</sup>The proposition (6) has been removed.

- (ii) the vertices of  $G_6$  is non empty,
- (iii) the source of  $G_6$  is a function from the edges of  $G_6$  into the vertices of  $G_6$ , and
- (iv) the target of  $G_6$  is a function from the edges of  $G_6$  into the vertices of  $G_6$ .
- (8) Let V be a non empty set, E be a set, and S, T be functions from E into V. Then
- (i) the vertices of createGraph(V, E, S, T) = V,
- (ii) the edges of createGraph(V, E, S, T) = E,
- (iii) the source of createGraph(V, E, S, T) = S, and
- (iv) the target of createGraph(V, E, S, T) = T.
- (9)  $\operatorname{dom}(G_6.\operatorname{set}(n,x)) = \operatorname{dom} G_6 \cup \{n\}.$
- (10) dom  $G_6 \subseteq \text{dom}(G_6.\text{set}(n, x)).$
- (11)  $(G_6.set(n, x))(n) = x.$
- (12) If  $n_1 \neq n_2$ , then  $G_6(n_2) = (G_6.set(n_1, x))(n_2)$ .
- (13) Suppose  $n \notin$  the graph selectors. Then
- (i) the vertices of G = the vertices of G.set(n, x),
- (ii) the edges of G = the edges of G.set(n, x),
- (iii) the source of G = the source of G.set(n, x),
- (iv) the target of G = the target of G.set(n, x), and
- (v) G.set(n, x) is a graph.
- (14) The vertices of  $G_6$ .set(VertexSelector, x) = x and the edges of  $G_6$ .set(EdgeSelector, x) = x and the source of  $G_6$ .set(SourceSelector, x) = x and the target of  $G_6$ .set(TargetSelector, x) = x.
- (15) If  $n_1 \neq n_2$ , then  $n_1 \in \text{dom}(G_6.\text{set}(n_1, x).\text{set}(n_2, y))$  and  $n_2 \in \text{dom}(G_6.\text{set}(n_1, x).\text{set}(n_2, y))$  and  $(G_6.\text{set}(n_1, x).\text{set}(n_2, y))(n_1) = x$  and  $(G_6.\text{set}(n_1, x).\text{set}(n_2, y))(n_2) = y.$
- (16) If e joins x and y in G, then  $x \in$  the vertices of G and  $y \in$  the vertices of G.
- (17) If e joins x and y in G, then e joins y and x in G.
- (18) If e joins  $x_1$  and  $y_1$  in G and e joins  $x_2$  and  $y_2$  in G, then  $x_1 = x_2$  and  $y_1 = y_2$  or  $x_1 = y_2$  and  $y_1 = x_2$ .
- (19) e joins x and y in G iff e joins x to y in G or e joins y to x in G.
- (20) Suppose e joins x and y in G but  $x \in X$  and  $y \in Y$  or  $x \in Y$  and  $y \in X$ . Then e joins a vertex from X and a vertex from Y in G.
- (21) G is loopless iff for every set v it is not true that there exists a set e such that e joins v and v in G.
- (22) For every finite loopless graph G and for every vertex v of G holds v.degree() = card(v.edgesInOut()).

- (23) For every non trivial graph G and for every vertex v of G holds (the vertices of  $G \setminus \{v\}$  is non empty.
- (24) For every non trivial graph G there exist vertices  $v_1$ ,  $v_2$  of G such that  $v_1 \neq v_2$ .
- (25) For every trivial graph G there exists a vertex v of G such that the vertices of  $G = \{v\}$ .
- (26) For every trivial loopless graph G holds the edges of  $G = \emptyset$ .
- (27) If the edges of  $G = \emptyset$ , then G is simple.
- (28) For every finite graph G holds  $G.order() \ge 1$ .
- (29) For every finite graph G holds G.order() = 1 iff G is trivial.
- (30) For every finite graph G holds G.order() = 1 iff there exists a vertex v of G such that the vertices of  $G = \{v\}$ .
- (31)  $e \in$  the edges of G but (the source of G) $(e) \in X$  or (the target of G) $(e) \in X$  iff  $e \in G$ .edgesInOut(X).
- (32)  $G.edgesInto(X) \subseteq G.edgesInOut(X)$  and  $G.edgesOutOf(X) \subseteq G.edgesInOut(X)$ .
- (33) The edges of G = G.edgesInOut(the vertices of G).
- (34)  $e \in$  the edges of G and (the source of G) $(e) \in X$  and (the target of G) $(e) \in X$  iff  $e \in G$ .edgesBetween(X).
- (35) If  $x \in X$  and  $y \in X$  and e joins x and y in G, then  $e \in G.$ edgesBetween(X).
- (36)  $G.edgesBetween(X) \subseteq G.edgesInOut(X).$
- (37) The edges of G = G.edgesBetween(the vertices of G).
- (38) (The edges of G) \ G.edgesInOut(X) = G.edgesBetween((the vertices of G) \ X).
- (39) If  $X \subseteq Y$ , then G.edgesBetween $(X) \subseteq G.edgesBetween(Y)$ .
- (40) For every graph G and for all sets  $X_1, X_2, Y_1, Y_2$  such that  $X_1 \subseteq X_2$ and  $Y_1 \subseteq Y_2$  holds G.edgesBetween $(X_1, Y_1) \subseteq G.edgesBetween(X_2, Y_2)$ .
- (41) For every graph G and for all sets  $X_1, X_2, Y_1, Y_2$  such that  $X_1 \subseteq X_2$  and  $Y_1 \subseteq Y_2$  holds G.edgesDBetween $(X_1, Y_1) \subseteq G.edgesDBetween<math>(X_2, Y_2)$ .
- (42) For every graph G and for every vertex v of G holds v.edgesIn() =  $G.edgesDBetween(the vertices of G, \{v\})$  and v.edgesOut() =  $G.edgesDBetween(\{v\}, the vertices of G).$
- (43) G is a subgraph of G.
- (44)  $G_3$  is a subgraph of  $G_4$  and  $G_4$  is a subgraph of  $G_3$  if and only if the following conditions are satisfied:
  - (i) the vertices of  $G_3$  = the vertices of  $G_4$ ,
  - (ii) the edges of  $G_3$  = the edges of  $G_4$ ,
- (iii) the source of  $G_3$  = the source of  $G_4$ , and

- (iv) the target of  $G_3$  = the target of  $G_4$ .
- (45) Let  $G_3$  be a graph,  $G_4$  be a subgraph of  $G_3$ , and x be a set. Then
  - (i) if  $x \in$  the vertices of  $G_4$ , then  $x \in$  the vertices of  $G_3$ , and
- (ii) if  $x \in$  the edges of  $G_4$ , then  $x \in$  the edges of  $G_3$ .
- (46) For every graph  $G_3$  and for every subgraph  $G_4$  of  $G_3$  holds every subgraph of  $G_4$  is a subgraph of  $G_3$ .
- (47) Let G be a graph and  $G_3$ ,  $G_4$  be subgraphs of G. Suppose the vertices of  $G_3 \subseteq$  the vertices of  $G_4$  and the edges of  $G_3 \subseteq$  the edges of  $G_4$ . Then  $G_3$  is a subgraph of  $G_4$ .
- (48) Let  $G_3$  be a graph and  $G_4$  be a subgraph of  $G_3$ . Then
  - (i) the source of  $G_4 = ($ the source of  $G_3) \upharpoonright ($ the edges of  $G_4),$  and
  - (ii) the target of  $G_4 = (\text{the target of } G_3) \upharpoonright (\text{the edges of } G_4).$
- (49) Let G be a graph,  $V_1$ ,  $V_2$ ,  $E_1$ ,  $E_2$  be sets,  $G_3$  be a subgraph of G induced by  $V_1$  and  $E_1$ , and  $G_4$  be a subgraph of G induced by  $V_2$  and  $E_2$ . Suppose  $V_2 \subseteq V_1$  and  $E_2 \subseteq E_1$  and  $V_2$  is a non empty subset of the vertices of G and  $E_2 \subseteq G$ .edgesBetween $(V_2)$ . Then  $G_4$  is a subgraph of  $G_3$ .
- (50) Let  $G_3$  be a non trivial graph, v be a vertex of  $G_3$ , and  $G_4$  be a subgraph of  $G_3$  with vertex v removed. Then the vertices of  $G_4 =$  (the vertices of  $G_3) \setminus \{v\}$  and the edges of  $G_4 = G_3$ .edgesBetween((the vertices of  $G_3) \setminus \{v\}$ ).
- (51) Let  $G_3$  be a finite non trivial graph, v be a vertex of  $G_3$ , and  $G_4$  be a subgraph of  $G_3$  with vertex v removed. Then  $G_4$ .order() + 1 =  $G_3$ .order() and  $G_4$ .size() + card(v.edgesInOut()) =  $G_3$ .size().
- (52) Let  $G_3$  be a graph, V be a set, and  $G_4$  be a subgraph of  $G_3$  with vertices V removed. Suppose  $V \subset$  the vertices of  $G_3$ . Then the vertices of  $G_4 =$  (the vertices of  $G_3) \setminus V$  and the edges of  $G_4 = G_3$ .edgesBetween((the vertices of  $G_3) \setminus V$ ).
- (53) Let  $G_3$  be a finite graph, V be a subset of the vertices of  $G_3$ , and  $G_4$  be a subgraph of  $G_3$  with vertices V removed. If  $V \neq$  the vertices of  $G_3$ , then  $G_4$ .order() + card V = G\_3.order() and  $G_4.size() + card(G_3.edgesInOut(V)) = G_3.size().$
- (54) Let  $G_3$  be a graph, e be a set, and  $G_4$  be a subgraph of  $G_3$  with edge e removed. Then the vertices of  $G_4$  = the vertices of  $G_3$  and the edges of  $G_4$  = (the edges of  $G_3$ ) \ {e}.
- (55) Let  $G_3$  be a finite graph, e be a set, and  $G_4$  be a subgraph of  $G_3$  with edge e removed. Then  $G_3$ .order() =  $G_4$ .order() and if  $e \in$  the edges of  $G_3$ , then  $G_4$ .size() + 1 =  $G_3$ .size().
- (56) Let  $G_3$  be a graph, E be a set, and  $G_4$  be a subgraph of  $G_3$  with edges E removed. Then the vertices of  $G_4$  = the vertices of  $G_3$  and the edges of  $G_4$  = (the edges of  $G_3) \setminus E$ .

- (57) For every finite graph  $G_3$  and for every set E and for every subgraph  $G_4$  of  $G_3$  with edges E removed holds  $G_3$ .order() =  $G_4$ .order().
- (58) Let  $G_3$  be a finite graph, E be a subset of the edges of  $G_3$ , and  $G_4$  be a subgraph of  $G_3$  with edges E removed. Then  $G_4$ .size() + card  $E = G_3$ .size().
- (59)  $e \in v.edgesIn()$  iff  $e \in the edges of G and (the target of G)(e) = v.$
- (60)  $e \in v.edgesIn()$  iff there exists a set x such that e joins x to v in G.
- (61)  $e \in v.$ edgesOut() iff  $e \in$  the edges of G and (the source of G)(e) = v.
- (62)  $e \in v.edgesOut()$  iff there exists a set x such that e joins v to x in G.
- (63)  $v.edgesInOut() = v.edgesIn() \cup v.edgesOut().$
- (64)  $e \in v.edgesInOut()$  iff  $e \in the edges of G but (the source of G)(e) = v$ or (the target of G)(e) = v.
- (65) If e joins  $v_1$  and x in G, then  $e \in v_1$ .edgesInOut().
- (66) If e joins  $v_1$  and  $v_2$  in G, then  $e \in v_1$ .edgesIn() and  $e \in v_2$ .edgesOut() or  $e \in v_2$ .edgesIn() and  $e \in v_1$ .edgesOut().
- (67)  $e \in v_1.$ edgesInOut() iff there exists a vertex  $v_2$  of G such that e joins  $v_1$  and  $v_2$  in G.
- (68) If  $e \in v.edgesInOut()$  and e joins x and y in G, then v = x or v = y.
- (69) If e joins  $v_1$  and  $v_2$  in G, then  $v_1$ .adj $(e) = v_2$  and  $v_2$ .adj $(e) = v_1$ .
- (70)  $e \in v.edgesInOut()$  iff e joins v and v.adj(e) in G.
- (71) Let G be a finite graph, e be a set, and  $v_1, v_2$  be vertices of G. If e joins  $v_1$  and  $v_2$  in G, then  $1 \le v_1$ .degree() and  $1 \le v_2$ .degree().
- (72)  $x \in v.inNeighbors()$  iff there exists a set e such that e joins x to v in G.
- (73)  $x \in v.outNeighbors()$  iff there exists a set e such that e joins v to x in G.
- (74)  $x \in v.$ allNeighbors() iff there exists a set e such that e joins v and x in G.
- (75) Let  $G_3$  be a graph,  $G_4$  be a subgraph of  $G_3$ , and x, y, e be sets. Then
  - (i) if e joins x and y in  $G_4$ , then e joins x and y in  $G_3$ ,
  - (ii) if e joins x to y in  $G_4$ , then e joins x to y in  $G_3$ ,
- (iii) if e joins a vertex from x and a vertex from y in  $G_4$ , then e joins a vertex from x and a vertex from y in  $G_3$ , and
- (iv) if e joins a vertex from x to a vertex from y in  $G_4$ , then e joins a vertex from x to a vertex from y in  $G_3$ .
- (76) Let  $G_3$  be a graph,  $G_4$  be a subgraph of  $G_3$ , and x, y, e be sets such that  $e \in$  the edges of  $G_4$ . Then
  - (i) if e joins x and y in  $G_3$ , then e joins x and y in  $G_4$ ,
  - (ii) if e joins x to y in  $G_3$ , then e joins x to y in  $G_4$ ,

- (iii) if e joins a vertex from x and a vertex from y in  $G_3$ , then e joins a vertex from x and a vertex from y in  $G_4$ , and
- (iv) if e joins a vertex from x to a vertex from y in  $G_3$ , then e joins a vertex from x to a vertex from y in  $G_4$ .
- (77) For every graph  $G_3$  and for every spanning subgraph  $G_4$  of  $G_3$  holds every spanning subgraph of  $G_4$  is a spanning subgraph of  $G_3$ .
- (78) For every finite graph  $G_3$  and for every subgraph  $G_4$  of  $G_3$  holds  $G_4$ .order()  $\leq G_3$ .order() and  $G_4$ .size()  $\leq G_3$ .size().
- (79) Let  $G_3$  be a graph,  $G_4$  be a subgraph of  $G_3$ , and X be a set. Then  $G_4$ .edgesInto $(X) \subseteq G_3$ .edgesInto(X) and  $G_4$ .edgesOutOf $(X) \subseteq G_3$ .edgesOutOf(X) and  $G_4$ .edgesInOut $(X) \subseteq G_3$ .edgesInOut(X) and  $G_4$ .edgesBetween $(X) \subseteq G_3$ .edgesBetween(X).
- (80) For every graph  $G_3$  and for every subgraph  $G_4$  of  $G_3$  and for all sets X, Y holds  $G_4$ .edgesBetween $(X, Y) \subseteq G_3$ .edgesBetween(X, Y) and  $G_4$ .edgesDBetween $(X, Y) \subseteq G_3$ .edgesDBetween(X, Y).
- (81) Let  $G_3$  be a graph,  $G_4$  be a subgraph of  $G_3$ ,  $v_1$  be a vertex of  $G_3$ , and  $v_2$  be a vertex of  $G_4$ . If  $v_1 = v_2$ , then  $v_2$ .edgesIn()  $\subseteq v_1$ .edgesIn() and  $v_2$ .edgesOut()  $\subseteq v_1$ .edgesOut() and  $v_2$ .edgesInOut()  $\subseteq v_1$ .edgesInOut().
- (82) Let  $G_3$  be a graph,  $G_4$  be a subgraph of  $G_3$ ,  $v_1$  be a vertex of  $G_3$ , and  $v_2$  be a vertex of  $G_4$ . Suppose  $v_1 = v_2$ . Then  $v_2$ .edgesIn() =  $v_1$ .edgesIn()  $\cap$  the edges of  $G_4$  and  $v_2$ .edgesOut() =  $v_1$ .edgesOut()  $\cap$  the edges of  $G_4$  and  $v_2$ .edgesInOut() =  $v_1$ .edgesInOut()  $\cap$  the edges of  $G_4$ .
- (83) Let  $G_3$  be a graph,  $G_4$  be a subgraph of  $G_3$ ,  $v_1$  be a vertex of  $G_3$ ,  $v_2$  be a vertex of  $G_4$ , and e be a set. If  $v_1 = v_2$  and  $e \in$  the edges of  $G_4$ , then  $v_1$ .adj $(e) = v_2$ .adj(e).
- (84) Let  $G_3$  be a finite graph,  $G_4$  be a subgraph of  $G_3$ ,  $v_1$  be a vertex of  $G_3$ , and  $v_2$  be a vertex of  $G_4$ . If  $v_1 = v_2$ , then  $v_2$ .inDegree()  $\leq v_1$ .inDegree() and  $v_2$ .outDegree()  $\leq v_1$ .outDegree() and  $v_2$ .degree()  $\leq v_1$ .degree().
- (85) Let  $G_3$  be a graph,  $G_4$  be a subgraph of  $G_3$ ,  $v_1$  be a vertex of  $G_3$ , and  $v_2$  be a vertex of  $G_4$ . If  $v_1 = v_2$ , then  $v_2$ .inNeighbors()  $\subseteq v_1$ .inNeighbors() and  $v_2$ .outNeighbors()  $\subseteq v_1$ .outNeighbors() and  $v_2$ .allNeighbors()  $\subseteq v_1$ .allNeighbors().
- (86) Let  $G_3$  be a graph,  $G_4$  be a subgraph of  $G_3$ ,  $v_1$  be a vertex of  $G_3$ , and  $v_2$  be a vertex of  $G_4$ . If  $v_1 = v_2$  and  $v_1$  is isolated, then  $v_2$  is isolated.
- (87) Let  $G_3$  be a graph,  $G_4$  be a subgraph of  $G_3$ ,  $v_1$  be a vertex of  $G_3$ , and  $v_2$  be a vertex of  $G_4$ . If  $v_1 = v_2$  and  $v_1$  is endvertex, then  $v_2$  is endvertex or isolated.
- (88) If  $G_3 =_G G_4$  and  $G_4 =_G G_7$ , then  $G_3 =_G G_7$ .
- (89) Let G be a graph and  $G_3$ ,  $G_4$  be subgraphs of G. Suppose the vertices of  $G_3$  = the vertices of  $G_4$  and the edges of  $G_3$  = the edges of  $G_4$ . Then

 $G_3 =_G G_4.$ 

- (90)  $G_3 =_G G_4$  iff  $G_3$  is a subgraph of  $G_4$  and  $G_4$  is a subgraph of  $G_3$ .
- (91) Suppose  $G_3 =_G G_4$ . Then
  - (i) if e joins x and y in  $G_3$ , then e joins x and y in  $G_4$ ,
- (ii) if e joins x to y in  $G_3$ , then e joins x to y in  $G_4$ ,
- (iii) if e joins a vertex from X and a vertex from Y in  $G_3$ , then e joins a vertex from X and a vertex from Y in  $G_4$ , and
- (iv) if e joins a vertex from X to a vertex from Y in  $G_3$ , then e joins a vertex from X to a vertex from Y in  $G_4$ .
- (92) Suppose  $G_3 =_G G_4$ . Then
  - (i) if  $G_3$  is finite, then  $G_4$  is finite,
  - (ii) if  $G_3$  is loopless, then  $G_4$  is loopless,
- (iii) if  $G_3$  is trivial, then  $G_4$  is trivial,
- (iv) if  $G_3$  is non-multi, then  $G_4$  is non-multi,
- (v) if  $G_3$  is non-directed-multi, then  $G_4$  is non-directed-multi,
- (vi) if  $G_3$  is simple, then  $G_4$  is simple, and
- (vii) if  $G_3$  is directed-simple, then  $G_4$  is directed-simple.
- (93) If  $G_3 =_G G_4$ , then  $G_3$ .order() =  $G_4$ .order() and  $G_3$ .size() =  $G_4$ .size() and  $G_3$ .edgesInto(X) =  $G_4$ .edgesInto(X) and  $G_3$ .edgesOutOf(X) =  $G_4$ .edgesOutOf(X) and  $G_3$ .edgesInOut(X) =  $G_4$ .edgesInOut(X) and  $G_3$ .edgesBetween(X) =  $G_4$ .edgesBetween(X) and  $G_3$ .edgesDBetween(X, Y) =  $G_4$ .edgesDBetween(X, Y).
- (94) If  $G_3 =_G G_4$  and  $G_7$  is a subgraph of  $G_3$ , then  $G_7$  is a subgraph of  $G_4$ .
- (95) If  $G_3 =_G G_4$  and  $G_3$  is a subgraph of  $G_7$ , then  $G_4$  is a subgraph of  $G_7$ .
- (96) For all subgraphs  $G_3$ ,  $G_4$  of G induced by V and E holds  $G_3 =_G G_4$ .
- (97) For every graph  $G_3$  and for every subgraph  $G_4$  of  $G_3$  induced by the vertices of  $G_3$  holds  $G_3 =_G G_4$ .
- (98) Let  $G_3$ ,  $G_4$  be graphs, V, E be sets, and  $G_7$  be a subgraph of  $G_3$  induced by V and E. If  $G_3 =_G G_4$ , then  $G_7$  is a subgraph of  $G_4$  induced by V and E.
- (99) Let  $v_1$  be a vertex of  $G_3$  and  $v_2$  be a vertex of  $G_4$ . Suppose  $v_1 = v_2$  and  $G_3 =_G G_4$ . Then  $v_1.\text{edgesIn}() = v_2.\text{edgesIn}()$  and  $v_1.\text{edgesOut}() = v_2.\text{edgesOut}()$  and  $v_1.\text{edgesInOut}() = v_2.\text{edgesInOut}()$  and  $v_1.\text{adj}(e) = v_2.\text{adj}(e)$  and  $v_1.\text{inDegree}() = v_2.\text{inDegree}()$  and  $v_1.\text{outDegree}() = v_2.\text{outDegree}()$  and  $v_1.\text{degree}() = v_2.\text{degree}()$  and  $v_1.\text{inNeighbors}() = v_2.\text{inNeighbors}()$  and  $v_1.\text{outNeighbors}() = v_2.\text{outNeighbors}()$ .
- (100) Let  $v_1$  be a vertex of  $G_3$  and  $v_2$  be a vertex of  $G_4$  such that  $v_1 = v_2$  and  $G_3 =_G G_4$ . Then
  - (i) if  $v_1$  is isolated, then  $v_2$  is isolated, and

- (ii) if  $v_1$  is endvertex, then  $v_2$  is endvertex.
- (101) Let G be a graph and  $G_3$ ,  $G_4$  be subgraphs of G. Suppose  $G_3 \subset G_4$ . Then the vertices of  $G_3 \subset$  the vertices of  $G_4$  or the edges of  $G_3 \subset$  the edges of  $G_4$ .
- (102) Let G be a graph and  $G_3$ ,  $G_4$  be subgraphs of G. Suppose  $G_3 \subset G_4$ . Then
  - (i) there exists a set v such that  $v \in$  the vertices of  $G_4$  and  $v \notin$  the vertices of  $G_3$ , or
  - (ii) there exists a set e such that  $e \in$  the edges of  $G_4$  and  $e \notin$  the edges of  $G_3$ .

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# Walks in Graphs<sup>1</sup>

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**Summary.** We define walks for graphs introduced in [9], introduce walk attributes and functors for walk creation and modification of walks. Subwalks of a walk are also defined. In our rendition, walks are alternating finite sequences of vertices and edges.

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The notation and terminology used here are introduced in the following papers: [14], [12], [16], [13], [18], [6], [4], [5], [1], [10], [17], [7], [3], [19], [15], [8], [2], [9], and [11].

#### 1. Preliminaries

The following propositions are true:

- (1) For all odd natural numbers x, y holds x < y iff  $x + 2 \le y$ .
- (2) Let X be a set and k be a natural number. Suppose  $X \subseteq \text{Seg } k$ . Let m, n be natural numbers. If  $m \in \text{dom Sgm } X$  and n = (Sgm X)(m), then  $m \leq n$ .
- (3) For every set X and for every finite sequence  $f_2$  of elements of X and for every FinSubsequence  $f_1$  of  $f_2$  holds len Seq  $f_1 \leq \text{len } f_2$ .
- (4) Let X be a set,  $f_2$  be a finite sequence of elements of X,  $f_1$  be a Fin-Subsequence of  $f_2$ , and m be a natural number. Suppose  $m \in \text{dom Seq } f_1$ . Then there exists a natural number n such that  $n \in \text{dom } f_2$  and  $m \leq n$ and  $(\text{Seq } f_1)(m) = f_2(n)$ .

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- (5) For every set X and for every finite sequence  $f_2$  of elements of X and for every FinSubsequence  $f_1$  of  $f_2$  holds len Seq  $f_1 = \operatorname{card} f_1$ .
- (6) Let X be a set,  $f_2$  be a finite sequence of elements of X, and  $f_1$  be a FinSubsequence of  $f_2$ . Then dom Seq  $f_1 = \text{dom Sgm dom } f_1$ .

### 2. WALK DEFINITIONS

Let G be a graph. A finite sequence of elements of the vertices of G is said to be a vertex sequence of G if:

(Def. 1) For every natural number n such that  $1 \le n$  and n < len it there existsa set e such that e joins it(n) and it(n+1) in G.

Let G be a graph. A finite sequence of elements of the edges of G is said to be a edge sequence of G if it satisfies the condition (Def. 2).

(Def. 2) There exists a finite sequence  $v_1$  of elements of the vertices of G such that  $\operatorname{len} v_1 = \operatorname{len} \operatorname{it} + 1$  and for every natural number n such that  $1 \leq n$  and  $n \leq \operatorname{len} \operatorname{it}$  holds  $\operatorname{it}(n)$  joins  $v_1(n)$  and  $v_1(n+1)$  in G.

Let G be a graph. A finite sequence of elements of (the vertices of G)  $\cup$  (the edges of G) is said to be a walk of G if it satisfies the conditions (Def. 3).

(Def. 3)(i) len it is odd,

- (ii)  $it(1) \in the vertices of G, and$
- (iii) for every odd natural number n such that n < len it holds it(n+1) joins it(n) and it(n+2) in G.

Let G be a graph and let W be a walk of G. One can verify that  $\operatorname{len} W$  is odd and non empty.

Let G be a graph and let v be a vertex of G. The functor G.walkOf(v) yielding a walk of G is defined as follows:

(Def. 4)  $G.walkOf(v) = \langle v \rangle.$ 

Let G be a graph and let x, y, e be sets. The functor G.walkOf(x, e, y) yielding a walk of G is defined as follows:

(Def. 5) G.walkOf $(x, e, y) = \begin{cases} \langle x, e, y \rangle, \text{ if } e \text{ joins } x \text{ and } y \text{ in } G, \\ G.$ walkOf $(\text{choose}(\text{the vertices of } G)), \text{ otherwise.} \end{cases}$ 

Let G be a graph and let W be a walk of G. The functor W.first() yields a vertex of G and is defined as follows:

(Def. 6) W.first() = W(1).

The functor W.last() yields a vertex of G and is defined by:

(Def. 7) W.last() = W(len W).

Let G be a graph, let W be a walk of G, and let n be a natural number. The functor W.vertexAt(n) yielding a vertex of G is defined as follows: (Def. 8)  $W.vertexAt(n) = \begin{cases} W(n), \text{ if } n \text{ is odd and } n \leq \operatorname{len} W, \\ W.first(), \text{ otherwise.} \end{cases}$ 

Let G be a graph and let W be a walk of G. The functor W.reverse() yielding a walk of G is defined as follows:

(Def. 9) W.reverse() = Rev(W).

Let G be a graph and let  $W_1$ ,  $W_2$  be walks of G. The functor  $W_1$ .append $(W_2)$  yields a walk of G and is defined by:

(Def. 10) 
$$W_1$$
.append $(W_2) = \begin{cases} W_1 \frown W_2, \text{ if } W_1.\text{last}() = W_2.\text{first}(), \\ W_1, \text{ otherwise.} \end{cases}$ 

Let G be a graph, let W be a walk of G, and let m, n be natural numbers. The functor  $W.\operatorname{cut}(m, n)$  yields a walk of G and is defined by:

(Def. 11) 
$$W.\operatorname{cut}(m,n) = \begin{cases} \langle W(m), \dots, W(n) \rangle, & \text{if } m \text{ is odd and } n \text{ is odd and } \\ m \leq n \text{ and } n \leq \operatorname{len} W, \\ W, & \text{otherwise.} \end{cases}$$

Let G be a graph, let W be a walk of G, and let m, n be natural numbers. The functor W.remove(m, n) yielding a walk of G is defined by:

(Def. 12)  $W.\text{remove}(m,n) = \begin{cases} (W.\text{cut}(1,m)).\text{append}((W.\text{cut}(n, \text{len }W))), \\ \text{if } m \text{ is odd and } n \text{ is odd and } m \leq n \text{ and} \\ n \leq \text{len } W \text{ and } W(m) = W(n), \\ W, \text{ otherwise.} \end{cases}$ 

Let G be a graph, let W be a walk of G, and let e be a set. The functor W.addEdge(e) yields a walk of G and is defined as follows:

- (Def. 13) W.addEdge(e) = W.append((G.walkOf(W.last(), e, W.last().adj(e)))).Let G be a graph and let W be a walk of G. The functor W.vertexSeq() yielding a vertex sequence of G is defined by:
- (Def. 14)  $\operatorname{len} W + 1 = 2 \cdot \operatorname{len}(W.\operatorname{vertexSeq}())$  and for every natural number n such that  $1 \leq n$  and  $n \leq \operatorname{len}(W.\operatorname{vertexSeq}())$  holds  $W.\operatorname{vertexSeq}()(n) = W(2 \cdot n 1)$ .

Let G be a graph and let W be a walk of G. The functor W.edgeSeq() yields a edge sequence of G and is defined by:

(Def. 15)  $\operatorname{len} W = 2 \cdot \operatorname{len}(W.\operatorname{edgeSeq}()) + 1$  and for every natural number n such that  $1 \leq n$  and  $n \leq \operatorname{len}(W.\operatorname{edgeSeq}())$  holds  $W.\operatorname{edgeSeq}()(n) = W(2 \cdot n)$ .

Let G be a graph and let W be a walk of G. The functor W.vertices() yields a finite subset of the vertices of G and is defined as follows:

(Def. 16) W.vertices() = rng(W.vertexSeq()).

Let G be a graph and let W be a walk of G. The functor W.edges() yields a finite subset of the edges of G and is defined by:

(Def. 17) W.edges() = rng(W.edgeSeq()).

Let G be a graph and let W be a walk of G. The functor W.length() yielding a natural number is defined by: (Def. 18) W.length() = len(W.edgeSeq()).

Let G be a graph, let W be a walk of G, and let v be a set. The functor W.find(v) yields an odd natural number and is defined by:

- (Def. 19)(i)  $W.\operatorname{find}(v) \leq \operatorname{len} W$  and  $W(W.\operatorname{find}(v)) = v$  and for every odd natural number n such that  $n \leq \operatorname{len} W$  and W(n) = v holds  $W.\operatorname{find}(v) \leq n$  if  $v \in W.\operatorname{vertices}()$ ,
  - (ii)  $W.\operatorname{find}(v) = \operatorname{len} W$ , otherwise.

Let G be a graph, let W be a walk of G, and let n be a natural number. The functor  $W.\operatorname{find}(n)$  yielding an odd natural number is defined by:

- (Def. 20)(i)  $W.\operatorname{find}(n) \leq \operatorname{len} W$  and  $W(W.\operatorname{find}(n)) = W(n)$  and for every odd natural number k such that  $k \leq \operatorname{len} W$  and W(k) = W(n) holds  $W.\operatorname{find}(n) \leq k$  if n is odd and  $n \leq \operatorname{len} W$ ,
  - (ii)  $W.\operatorname{find}(n) = \operatorname{len} W$ , otherwise.

Let G be a graph, let W be a walk of G, and let v be a set. The functor W.rfind(v) yields an odd natural number and is defined as follows:

- (Def. 21)(i)  $W.rfind(v) \leq \operatorname{len} W$  and W(W.rfind(v)) = v and for every odd natural number n such that  $n \leq \operatorname{len} W$  and W(n) = v holds  $n \leq W.rfind(v)$  if  $v \in W.vertices()$ ,
  - (ii) W.rfind(v) = len W, otherwise.

Let G be a graph, let W be a walk of G, and let n be a natural number. The functor W.rfind(n) yields an odd natural number and is defined by:

- (Def. 22)(i)  $W.rfind(n) \leq \operatorname{len} W$  and W(W.rfind(n)) = W(n) and for every odd natural number k such that  $k \leq \operatorname{len} W$  and W(k) = W(n) holds  $k \leq W.rfind(n)$  if n is odd and  $n \leq \operatorname{len} W$ ,
  - (ii) W.rfind(n) = len W, otherwise.

Let G be a graph, let u, v be sets, and let W be a walk of G. We say that W is walk from u to v if and only if:

(Def. 23) W.first() = u and W.last() = v.

Let G be a graph and let W be a walk of G. We say that W is closed if and only if:

(Def. 24) W.first() = W.last().

We say that W is directed if and only if:

(Def. 25) For every odd natural number n such that  $n < \operatorname{len} W$  holds (the source of G)(W(n+1)) = W(n).

We say that W is trivial if and only if:

(Def. 26) W.length() = 0.

We say that W is trail-like if and only if:

(Def. 27) W.edgeSeq() is one-to-one.

Let G be a graph and let W be a walk of G. We introduce W is open as an antonym of W is closed.

Let G be a graph and let W be a walk of G. We say that W is path-like if and only if the conditions (Def. 28) are satisfied.

(Def. 28)(i) W is trail-like, and

(ii) for all odd natural numbers m, n such that m < n and  $n \le \text{len } W$  holds if W(m) = W(n), then m = 1 and n = len W.

Let G be a graph and let W be a walk of G. We say that W is vertex-distinct if and only if:

(Def. 29) For all odd natural numbers m, n such that  $m \leq \operatorname{len} W$  and  $n \leq \operatorname{len} W$ and W(m) = W(n) holds m = n.

Let G be a graph and let W be a walk of G. We say that W is circuit-like if and only if:

(Def. 30) W is closed, trail-like, and non trivial.

We say that W is cycle-like if and only if:

(Def. 31) W is closed, path-like, and non trivial.

Let G be a graph. One can verify the following observations:

- \* every walk of G which is path-like is also trail-like,
- \* every walk of G which is trivial is also path-like,
- \* every walk of G which is trivial is also vertex-distinct,
- \* every walk of G which is vertex-distinct is also path-like,
- $\ast~$  every walk of G which is circuit-like is also closed, trail-like, and non trivial, and
- \* every walk of G which is cycle-like is also closed, path-like, and non trivial.

Let G be a graph. Observe that there exists a walk of G which is closed, directed, and trivial.

Let G be a graph. Observe that there exists a walk of G which is vertexdistinct.

Let G be a graph. A trail of G is a trail-like walk of G. A path of G is a path-like walk of G.

Let G be a graph. A dwalk of G is a directed walk of G. A dtrail of G is a directed trail of G. A dpath of G is a directed path of G.

Let G be a graph and let v be a vertex of G. Note that G.walkOf(v) is closed, directed, and trivial.

Let G be a graph and let x, e, y be sets. One can check that G.walkOf(x, e, y) is path-like.

Let G be a graph and let x, e be sets. Note that G.walkOf(x, e, x) is closed.

Let G be a graph and let W be a closed walk of G. One can check that W.reverse() is closed.

Let G be a graph and let W be a trivial walk of G. One can verify that W.reverse() is trivial.

Let G be a graph and let W be a trail of G. Note that W.reverse() is trail-like.

Let G be a graph and let W be a path of G. Observe that W.reverse() is path-like.

Let G be a graph and let  $W_1$ ,  $W_2$  be closed walks of G. Note that  $W_1$ .append $(W_2)$  is closed.

Let G be a graph and let  $W_1$ ,  $W_2$  be dwalks of G. One can verify that  $W_1$ .append $(W_2)$  is directed.

Let G be a graph and let  $W_1$ ,  $W_2$  be trivial walks of G. Observe that  $W_1$ .append $(W_2)$  is trivial.

Let G be a graph, let W be a dwalk of G, and let m, n be natural numbers. Note that  $W.\operatorname{cut}(m, n)$  is directed.

Let G be a graph, let W be a trivial walk of G, and let m, n be natural numbers. Observe that  $W.\operatorname{cut}(m, n)$  is trivial.

Let G be a graph, let W be a trail of G, and let m, n be natural numbers. Note that  $W.\operatorname{cut}(m, n)$  is trail-like.

Let G be a graph, let W be a path of G, and let m, n be natural numbers. Note that  $W.\operatorname{cut}(m, n)$  is path-like.

Let G be a graph, let W be a vertex-distinct walk of G, and let m, n be natural numbers. One can verify that  $W.\operatorname{cut}(m, n)$  is vertex-distinct.

Let G be a graph, let W be a closed walk of G, and let m, n be natural numbers. One can verify that W.remove(m, n) is closed.

Let G be a graph, let W be a dwalk of G, and let m, n be natural numbers. Note that W.remove(m, n) is directed.

Let G be a graph, let W be a trivial walk of G, and let m, n be natural numbers. One can check that W.remove(m, n) is trivial.

Let G be a graph, let W be a trail of G, and let m, n be natural numbers. Observe that W.remove(m, n) is trail-like.

Let G be a graph, let W be a path of G, and let m, n be natural numbers. Observe that W.remove(m, n) is path-like.

Let G be a graph and let W be a walk of G. A walk of G is called a subwalk of W if:

(Def. 32) It is walk from W.first() to W.last() and there exists a FinSubsequence  $e_1$  of W.edgeSeq() such that it.edgeSeq() = Seq  $e_1$ .

Let G be a graph, let W be a walk of G, and let m, n be natural numbers. Then W.remove(m, n) is a subwalk of W.

Let G be a graph and let W be a walk of G. Note that there exists a subwalk of W which is trail-like and path-like.

Let G be a graph and let W be a walk of G. A trail of W is a trail-like subwalk of W. A path of W is a path-like subwalk of W.

Let G be a graph and let W be a dwalk of G. One can verify that there exists a path of W which is directed.

Let G be a graph and let W be a dwalk of G. A dwalk of W is a directed subwalk of W. A dtrail of W is a directed trail of W. A dpath of W is a directed path of W.

Let G be a graph. The functor G.allWalks() yields a non empty subset of  $((\text{the vertices of } G) \cup (\text{the edges of } G))^*$  and is defined by:

(Def. 33) G.allWalks() = {W : W ranges over walks of G}.

Let G be a graph. The functor G.allTrails() yielding a non empty subset of G.allWalks() is defined by:

(Def. 34) G.allTrails() = {W : W ranges over trails of G}.

Let G be a graph. The functor G.allPaths() yields a non empty subset of G.allTrails() and is defined as follows:

(Def. 35) G.allPaths() = {W : W ranges over paths of G}.

Let G be a graph. The functor G.allDWalks() yields a non empty subset of G.allWalks() and is defined by:

(Def. 36) G.allDWalks() = {W : W ranges over dwalks of G}.

Let G be a graph. The functor G.allDTrails() yields a non empty subset of G.allTrails() and is defined as follows:

(Def. 37) G.allDTrails() = {W : W ranges over dtrails of G}.

Let G be a graph. The functor G.allDPaths() yields a non empty subset of G.allDTrails() and is defined by:

(Def. 38) G.allDPaths() = {W : W ranges over directed paths of G }.

Let G be a finite graph. One can check that G.allTrails() is finite.

Let G be a graph and let X be a non empty subset of G.allWalks(). We see that the element of X is a walk of G.

Let G be a graph and let X be a non empty subset of G.allTrails(). We see that the element of X is a trail of G.

Let G be a graph and let X be a non empty subset of G.allPaths(). We see that the element of X is a path of G.

Let G be a graph and let X be a non empty subset of G.allDWalks(). We see that the element of X is a dwalk of G.

Let G be a graph and let X be a non empty subset of G.allDTrails(). We see that the element of X is a dtrail of G.

Let G be a graph and let X be a non empty subset of G.allDPaths(). We see that the element of X is a dpath of G.

#### 3. Walk Theorems

For simplicity, we adopt the following rules: G,  $G_1$ ,  $G_2$  are graphs, W,  $W_1$ ,  $W_2$  are walks of G, e, x, y, z are sets, v is a vertex of G, and n, m are natural numbers.

We now state a number of propositions:

- $(8)^3$  For every odd natural number n such that  $n \leq \text{len } W$  holds  $W(n) \in \text{the vertices of } G$ .
- (9) For every even natural number n such that  $n \in \text{dom } W$  holds  $W(n) \in \text{the}$  edges of G.
- (10) Let n be an even natural number. Suppose  $n \in \text{dom } W$ . Then there exists an odd natural number  $n_1$  such that  $n_1 = n 1$  and  $n 1 \in \text{dom } W$  and  $n + 1 \in \text{dom } W$  and W(n) joins  $W(n_1)$  and W(n + 1) in G.
- (11) For every odd natural number n such that n < len W holds  $W(n+1) \in (W.\text{vertexAt}(n)).\text{edgesInOut}().$
- (12) For every odd natural number n such that 1 < n and  $n \leq \text{len } W$  holds  $W(n-1) \in (W.\text{vertexAt}(n)).\text{edgesInOut}().$
- (13) For every odd natural number n such that  $n < \operatorname{len} W$  holds  $n \in \operatorname{dom} W$ and  $n + 1 \in \operatorname{dom} W$  and  $n + 2 \in \operatorname{dom} W$ .
- (14)  $\operatorname{len}(G.\operatorname{walkOf}(v)) = 1$  and  $(G.\operatorname{walkOf}(v))(1) = v$  and  $(G.\operatorname{walkOf}(v)).\operatorname{first}() = v$  and  $(G.\operatorname{walkOf}(v)).\operatorname{last}() = v$  and  $G.\operatorname{walkOf}(v)$  is walk from v to v.
- (15) If e joins x and y in G, then len(G.walkOf(x, e, y)) = 3.
- (16) If e joins x and y in G, then (G.walkOf(x, e, y)).first() = x and (G.walkOf(x, e, y)).last() = y and G.walkOf(x, e, y) is walk from x to y.
- (17) For every walk  $W_1$  of  $G_1$  and for every walk  $W_2$  of  $G_2$  such that  $W_1 = W_2$  holds  $W_1$ .first() =  $W_2$ .first() and  $W_1$ .last() =  $W_2$ .last().
- (18) W is walk from x to y iff W(1) = x and  $W(\operatorname{len} W) = y$ .
- (19) If W is walk from x to y, then x is a vertex of G and y is a vertex of G.
- (20) Let  $W_1$  be a walk of  $G_1$  and  $W_2$  be a walk of  $G_2$ . If  $W_1 = W_2$ , then  $W_1$  is walk from x to y iff  $W_2$  is walk from x to y.
- (21) For every walk  $W_1$  of  $G_1$  and for every walk  $W_2$  of  $G_2$  such that  $W_1 = W_2$ and for every natural number n holds  $W_1$ .vertexAt $(n) = W_2$ .vertexAt(n).
- (22)  $\operatorname{len} W = \operatorname{len}(W.\operatorname{reverse}())$  and  $\operatorname{dom} W = \operatorname{dom}(W.\operatorname{reverse}())$  and  $\operatorname{rng} W = \operatorname{rng}(W.\operatorname{reverse}())$ .
- (23) W.first() = W.reverse().last() and W.last() = W.reverse().first().
- (24) W is walk from x to y iff W.reverse() is walk from y to x.

<sup>&</sup>lt;sup>3</sup>The proposition (7) has been removed.

- (25) If  $n \in \text{dom } W$ , then W(n) = W.reverse()((len W n) + 1) and (len W n) + 1  $\in \text{dom}(W$ .reverse()).
- (26) If  $n \in \text{dom}(W.\text{reverse}())$ , then W.reverse()(n) = W((len W n) + 1) and  $(\text{len } W n) + 1 \in \text{dom } W.$
- (27) W.reverse().reverse() = W.
- (28) For every walk  $W_1$  of  $G_1$  and for every walk  $W_2$  of  $G_2$  such that  $W_1 = W_2$  holds  $W_1$ .reverse() =  $W_2$ .reverse().
- (29) If  $W_1.last() = W_2.first()$ , then  $len(W_1.append(W_2)) + 1 = len W_1 + len W_2$ .
- (30) If  $W_1.last() = W_2.first()$ , then  $len W_1 \leq len(W_1.append(W_2))$  and  $len W_2 \leq len(W_1.append(W_2))$ .
- (31) If  $W_1.last() = W_2.first()$ , then  $(W_1.append(W_2)).first() = W_1.first()$  and  $(W_1.append(W_2)).last() = W_2.last()$  and  $W_1.append(W_2)$  is walk from  $W_1.first()$  to  $W_2.last()$ .
- (32) If  $W_1$  is walk from x to y and  $W_2$  is walk from y to z, then  $W_1$ .append $(W_2)$  is walk from x to z.
- (33) If  $n \in \operatorname{dom} W_1$ , then  $(W_1.\operatorname{append}(W_2))(n) = W_1(n)$  and  $n \in \operatorname{dom}(W_1.\operatorname{append}(W_2))$ .
- (34) If  $W_1$ .last() =  $W_2$ .first(), then for every natural number n such that n <len  $W_2$  holds  $(W_1.append(W_2))(\text{len } W_1 + n) = W_2(n+1)$  and  $\text{len } W_1 + n \in$ dom $(W_1.append(W_2))$ .
- (35) If  $n \in \text{dom}(W_1.\text{append}(W_2))$ , then  $n \in \text{dom} W_1$  or there exists a natural number k such that  $k < \text{len } W_2$  and  $n = \text{len } W_1 + k$ .
- (36) For all walks  $W_3$ ,  $W_4$  of  $G_1$  and for all walks  $W_5$ ,  $W_6$  of  $G_2$  such that  $W_3 = W_5$  and  $W_4 = W_6$  holds  $W_3$ .append $(W_4) = W_5$ .append $(W_6)$ .
- (37) Let m, n be odd natural numbers. Suppose  $m \le n$  and  $n \le \operatorname{len} W$ . Then  $\operatorname{len}(W.\operatorname{cut}(m,n)) + m = n + 1$  and for every natural number i such that  $i < \operatorname{len}(W.\operatorname{cut}(m,n))$  holds  $(W.\operatorname{cut}(m,n))(i+1) = W(m+i)$  and  $m+i \in \operatorname{dom} W$ .
- (38) Let m, n be odd natural numbers. Suppose  $m \le n$  and  $n \le \text{len } W$ . Then  $(W.\operatorname{cut}(m,n)).\operatorname{first}() = W(m)$  and  $(W.\operatorname{cut}(m,n)).\operatorname{last}() = W(n)$  and  $W.\operatorname{cut}(m,n)$  is walk from W(m) to W(n).
- (39) For all odd natural numbers m, n, o such that  $m \le n$  and  $n \le o$  and  $o \le \operatorname{len} W$  holds  $(W.\operatorname{cut}(m, n)).\operatorname{append}((W.\operatorname{cut}(n, o))) = W.\operatorname{cut}(m, o).$
- (40)  $W.\operatorname{cut}(1, \operatorname{len} W) = W.$
- (41) For every odd natural number n such that  $n < \operatorname{len} W$  holds  $G.\operatorname{walkOf}(W(n), W(n+1), W(n+2)) = W.\operatorname{cut}(n, n+2).$
- (42) For all odd natural numbers m, n such that  $m \le n$  and  $n < \operatorname{len} W$  holds  $(W.\operatorname{cut}(m,n)).\operatorname{addEdge}(W(n+1)) = W.\operatorname{cut}(m,n+2).$

- (43) For every odd natural number n such that  $n \leq \operatorname{len} W$  holds  $W.\operatorname{cut}(n,n) = \langle W.\operatorname{vertexAt}(n) \rangle.$
- (44) If m is odd and  $m \le n$ , then  $W.\operatorname{cut}(1, n).\operatorname{cut}(1, m) = W.\operatorname{cut}(1, m)$ .
- (45) For all odd natural numbers m, n such that  $m \le n$  and  $n \le \operatorname{len} W_1$  and  $W_1.\operatorname{last}() = W_2.\operatorname{first}()$  holds  $(W_1.\operatorname{append}(W_2)).\operatorname{cut}(m, n) = W_1.\operatorname{cut}(m, n).$
- (46) For every odd natural number m such that  $m \leq \operatorname{len} W$  holds  $\operatorname{len}(W.\operatorname{cut}(1,m)) = m$ .
- (47) For every odd natural number m and for every natural number x such that  $x \in \text{dom}(W.\text{cut}(1,m))$  and  $m \leq \text{len } W$  holds (W.cut(1,m))(x) = W(x).
- (48) Let m, n be odd natural numbers and i be a natural number. If  $m \le n$  and  $n \le \operatorname{len} W$  and  $i \in \operatorname{dom}(W.\operatorname{cut}(m,n))$ , then  $(W.\operatorname{cut}(m,n))(i) = W((m+i)-1)$  and  $(m+i)-1 \in \operatorname{dom} W$ .
- (49) For every walk  $W_1$  of  $G_1$  and for every walk  $W_2$  of  $G_2$  and for all natural numbers m, n such that  $W_1 = W_2$  holds  $W_1.cut(m, n) = W_2.cut(m, n)$ .
- (50) For all odd natural numbers m, n such that  $m \le n$  and  $n \le \operatorname{len} W$  and W(m) = W(n) holds  $\operatorname{len}(W.\operatorname{remove}(m, n)) + n = \operatorname{len} W + m$ .
- (51) If W is walk from x to y, then W.remove(m, n) is walk from x to y.
- (52)  $\operatorname{len}(W.\operatorname{remove}(m, n)) \leq \operatorname{len} W.$
- (53) W.remove(m,m) = W.
- (54) For all odd natural numbers m, n such that  $m \le n$  and  $n \le \operatorname{len} W$  and W(m) = W(n) holds  $(W.\operatorname{cut}(1,m)).\operatorname{last}() = (W.\operatorname{cut}(n,\operatorname{len} W)).\operatorname{first}().$
- (55) Let m, n be odd natural numbers. Suppose  $m \le n$  and  $n \le \operatorname{len} W$ and W(m) = W(n). Let x be a natural number. If  $x \in \operatorname{Seg} m$ , then  $(W.\operatorname{remove}(m, n))(x) = W(x)$ .
- (56) Let m, n be odd natural numbers. Suppose  $m \le n$  and  $n \le \operatorname{len} W$ and W(m) = W(n). Let x be a natural number. Suppose  $m \le x$  and  $x \le \operatorname{len}(W.\operatorname{remove}(m, n))$ . Then  $(W.\operatorname{remove}(m, n))(x) = W((x - m) + n)$ and (x - m) + n is a natural number and  $(x - m) + n \le \operatorname{len} W$ .
- (57) For all odd natural numbers m, n such that  $m \le n$  and  $n \le \operatorname{len} W$  and W(m) = W(n) holds  $\operatorname{len}(W.\operatorname{remove}(m, n)) = (\operatorname{len} W + m) n$ .
- (58) For every natural number m such that W(m) = W.last() holds W.remove(m, len W) = W.cut(1, m).
- (59) For every natural number m such that W.first() = W(m) holds W.remove(1,m) = W.cut(m, len W).
- (60) (W.remove(m, n)).first() = W.first() and (W.remove(m, n)).last() = W.last().
- (61) Let m, n be odd natural numbers and x be a natural number. Suppose  $m \le n$  and  $n \le \text{len } W$  and W(m) = W(n) and  $x \in \text{dom}(W.\text{remove}(m, n))$ .

Then  $x \in \text{Seg } m$  or  $m \leq x$  and  $x \leq \text{len}(W.\text{remove}(m, n))$ .

- (62) For every walk  $W_1$  of  $G_1$  and for every walk  $W_2$  of  $G_2$  and for all natural numbers m, n such that  $W_1 = W_2$  holds  $W_1$ .remove $(m, n) = W_2$ .remove(m, n).
- (63) If e joins W.last() and x in G, then W.addEdge(e) =  $W \cap \langle e, x \rangle$ .
- (64) If e joins W.last() and x in G, then (W.addEdge(e)).first() = W.first()and (W.addEdge(e)).last() = x and W.addEdge(e) is walk from W.first() to x.
- (65) If e joins W.last() and x in G, then len(W.addEdge(e)) = len W + 2.
- (66) Suppose e joins W.last() and x in G. Then  $(W.addEdge(e))(\operatorname{len} W+1) = e$  and  $(W.addEdge(e))(\operatorname{len} W+2) = x$  and for every natural number n such that  $n \in \operatorname{dom} W$  holds (W.addEdge(e))(n) = W(n).
- (67) If W is walk from x to y and e joins y and z in G, then W.addEdge(e) is walk from x to z.
- (68)  $1 \leq \operatorname{len}(W.\operatorname{vertexSeq}()).$
- (69) For every odd natural number n such that  $n \leq \operatorname{len} W$  holds  $2 \cdot ((n+1) \div 2) 1 = n$  and  $1 \leq (n+1) \div 2$  and  $(n+1) \div 2 \leq \operatorname{len}(W.\operatorname{vertexSeq}())$ .
- (70)  $(G.walkOf(v)).vertexSeq() = \langle v \rangle.$
- (71) If e joins x and y in G, then  $(G.walkOf(x, e, y)).vertexSeq() = \langle x, y \rangle$ .
- (72) W.first() = W.vertexSeq()(1) and W.last() = W.vertexSeq()(len(W.vertexSeq())).
- (73) For every odd natural number n such that  $n \leq \operatorname{len} W$  holds  $W.\operatorname{vertexAt}(n) = W.\operatorname{vertexSeq}()((n+1) \div 2).$
- (74)  $n \in \operatorname{dom}(W.\operatorname{vertexSeq}())$  iff  $2 \cdot n 1 \in \operatorname{dom} W$ .
- (75)  $(W.\operatorname{cut}(1, n)).\operatorname{vertexSeq}() \subseteq W.\operatorname{vertexSeq}().$
- (76) If e joins W.last() and x in G, then (W.addEdge(e)).vertexSeq() =  $W.vertexSeq() \cap \langle x \rangle.$
- (77) For every walk  $W_1$  of  $G_1$  and for every walk  $W_2$  of  $G_2$  such that  $W_1 = W_2$  holds  $W_1$ .vertexSeq() =  $W_2$ .vertexSeq().
- (78) For every even natural number n such that  $1 \le n$  and  $n \le \text{len } W$  holds  $n \div 2 \in \text{dom}(W.\text{edgeSeq}())$  and  $W(n) = W.\text{edgeSeq}()(n \div 2)$ .
- (79)  $n \in \operatorname{dom}(W.\operatorname{edgeSeq}())$  iff  $2 \cdot n \in \operatorname{dom} W$ .
- (80) For every natural number n such that  $n \in \text{dom}(W.\text{edgeSeq}())$  holds  $W.\text{edgeSeq}()(n) \in \text{the edges of } G.$
- (81) There exists an even natural number  $l_1$  such that  $l_1 = \operatorname{len} W 1$  and  $\operatorname{len}(W.\operatorname{edgeSeq}()) = l_1 \div 2.$
- (82)  $(W.\operatorname{cut}(1, n)).\operatorname{edgeSeq}() \subseteq W.\operatorname{edgeSeq}().$
- (83) If e joins W.last() and x in G, then (W.addEdge(e)).edgeSeq() =  $W.edgeSeq() \cap \langle e \rangle$ .

- (84) *e* joins *x* and *y* in *G* iff (*G*.walkOf(*x*, *e*, *y*)).edgeSeq() =  $\langle e \rangle$ .
- (85) W.reverse().edgeSeq() = Rev(W.edgeSeq()).
- (86) If  $W_1.last() = W_2.first()$ , then  $(W_1.append(W_2)).edgeSeq() = W_1.edgeSeq() \cap W_2.edgeSeq()$ .
- (87) For every walk  $W_1$  of  $G_1$  and for every walk  $W_2$  of  $G_2$  such that  $W_1 = W_2$  holds  $W_1.edgeSeq() = W_2.edgeSeq()$ .
- (88)  $x \in W$ .vertices() iff there exists an odd natural number n such that  $n \leq \operatorname{len} W$  and W(n) = x.
- (89)  $W.first() \in W.vertices()$  and  $W.last() \in W.vertices()$ .
- (90) For every odd natural number n such that  $n \leq \operatorname{len} W$  holds  $W.\operatorname{vertexAt}(n) \in W.\operatorname{vertices}().$
- (91)  $(G.walkOf(v)).vertices() = \{v\}.$
- (92) If e joins x and y in G, then  $(G.walkOf(x, e, y)).vertices() = \{x, y\}.$
- (93) W.vertices() = W.reverse().vertices().
- (94) If  $W_1.last() = W_2.first()$ , then  $(W_1.append(W_2)).vertices() = W_1.vertices() \cup W_2.vertices()$ .
- (95) For all odd natural numbers m, n such that  $m \le n$  and  $n \le \text{len } W$  holds  $(W.\text{cut}(m, n)).\text{vertices}() \subseteq W.\text{vertices}().$
- (96) If e joins W.last() and x in G, then  $(W.addEdge(e)).vertices() = W.vertices() \cup \{x\}.$
- (97) Let G be a finite graph, W be a walk of G, and e, x be sets. If e joins W.last() and x in G and  $x \notin W.vertices()$ , then card((W.addEdge(e)).vertices()) = card(W.vertices()) + 1.
- (98) If  $x \in W$ .vertices() and  $y \in W$ .vertices(), then there exists a walk of G which is walk from x to y.
- (99) For every walk  $W_1$  of  $G_1$  and for every walk  $W_2$  of  $G_2$  such that  $W_1 = W_2$  holds  $W_1$ .vertices() =  $W_2$ .vertices().
- (100)  $e \in W.edges()$  iff there exists an even natural number n such that  $1 \le n$ and  $n \le \text{len } W$  and W(n) = e.
- (101)  $e \in W.edges()$  iff there exists an odd natural number n such that n < len W and W(n+1) = e.
- (102)  $\operatorname{rng} W = W.\operatorname{vertices}() \cup W.\operatorname{edges}().$
- (103) If  $W_1.last() = W_2.first()$ , then  $(W_1.append(W_2)).edges() = W_1.edges() \cup W_2.edges()$ .
- (104) Suppose  $e \in W$ .edges(). Then there exist vertices  $v_2$ ,  $v_3$  of G and there exists an odd natural number n such that  $n + 2 \leq \text{len } W$  and  $v_2 = W(n)$  and e = W(n+1) and  $v_3 = W(n+2)$  and e joins  $v_2$  and  $v_3$  in G.
- (105)  $e \in W.edges()$  iff there exists a natural number n such that  $n \in dom(W.edgeSeq())$  and W.edgeSeq()(n) = e.

- (106) If  $e \in W$ .edges() and e joins x and y in G, then  $x \in W$ .vertices() and  $y \in W$ .vertices().
- (107)  $(W.\operatorname{cut}(m, n)).\operatorname{edges}() \subseteq W.\operatorname{edges}().$
- (108) W.edges() = W.reverse().edges().
- (109) e joins x and y in G iff  $(G.walkOf(x, e, y)).edges() = \{e\}.$
- (110)  $W.edges() \subseteq G.edgesBetween(W.vertices()).$
- (111) For every walk  $W_1$  of  $G_1$  and for every walk  $W_2$  of  $G_2$  such that  $W_1 = W_2$  holds  $W_1.edges() = W_2.edges()$ .
- (112) If e joins W.last() and x in G, then  $(W.addEdge(e)).edges() = W.edges() \cup \{e\}.$
- (113)  $\operatorname{len} W = 2 \cdot W.\operatorname{length}() + 1.$
- (114)  $\operatorname{len} W_1 = \operatorname{len} W_2$  iff  $W_1.\operatorname{length}() = W_2.\operatorname{length}()$ .
- (115) For every walk  $W_1$  of  $G_1$  and for every walk  $W_2$  of  $G_2$  such that  $W_1 = W_2$  holds  $W_1$ .length() =  $W_2$ .length().
- (116) For every odd natural number n such that  $n \leq \operatorname{len} W$  holds  $W.\operatorname{find}(W(n)) \leq n$  and  $W.\operatorname{rfind}(W(n)) \geq n$ .
- (117) For every walk  $W_1$  of  $G_1$  and for every walk  $W_2$  of  $G_2$  and for every set v such that  $W_1 = W_2$  holds  $W_1.find(v) = W_2.find(v)$  and  $W_1.rfind(v) = W_2.rfind(v)$ .
- (118) For every odd natural number n such that  $n \leq \operatorname{len} W$  holds  $W.\operatorname{find}(n) \leq n$  and  $W.\operatorname{rfind}(n) \geq n$ .
- (119) W is closed iff  $W(1) = W(\operatorname{len} W)$ .
- (120) W is closed iff there exists a set x such that W is walk from x to x.
- (121) W is closed iff W.reverse() is closed.
- (122) For every walk  $W_1$  of  $G_1$  and for every walk  $W_2$  of  $G_2$  such that  $W_1 = W_2$ and  $W_1$  is closed holds  $W_2$  is closed.
- (123) W is directed if and only if for every odd natural number n such that n < len W holds W(n+1) joins W(n) to W(n+2) in G.
- (124) Suppose W is directed and walk from x to y and e joins y to z in G. Then W.addEdge(e) is directed and W.addEdge(e) is walk from x to z.
- (125) For every dwalk W of G and for all natural numbers m, n holds  $W.\operatorname{cut}(m,n)$  is directed.
- (126) W is non trivial iff  $3 \leq \operatorname{len} W$ .
- (127) W is non trivial iff len  $W \neq 1$ .
- (128) If  $W.first() \neq W.last()$ , then W is non trivial.
- (129) W is trivial iff there exists a vertex v of G such that W = G.walkOf(v).
- (130) W is trivial iff W.reverse() is trivial.
- (131) If  $W_2$  is trivial, then  $W_1$ .append $(W_2) = W_1$ .

- (132) For all odd natural numbers m, n such that  $m \le n$  and  $n \le \text{len } W$  holds W.cut(m, n) is trivial iff m = n.
- (133) If e joins W.last() and x in G, then W.addEdge(e) is non trivial.
- (134) If W is non trivial, then there exists an odd natural number  $l_2$  such that  $l_2 = \text{len } W 2$  and  $(W.\text{cut}(1, l_2)).\text{addEdge}(W(l_2 + 1)) = W.$
- (135) If  $W_2$  is non trivial and  $W_2$ .edges()  $\subseteq W_1$ .edges(), then  $W_2$ .vertices()  $\subseteq W_1$ .vertices().
- (136) If W is non trivial, then for every vertex v of G such that  $v \in W$ .vertices() holds v is not isolated.
- (137) W is trivial iff  $W.edges() = \emptyset$ .
- (138) For every walk  $W_1$  of  $G_1$  and for every walk  $W_2$  of  $G_2$  such that  $W_1 = W_2$ and  $W_1$  is trivial holds  $W_2$  is trivial.
- (139) W is trail-like iff for all even natural numbers m, n such that  $1 \le m$  and m < n and  $n \le \operatorname{len} W$  holds  $W(m) \ne W(n)$ .
- (140) If len  $W \leq 3$ , then W is trail-like.
- (141) W is trail-like iff W.reverse() is trail-like.
- (142) For every trail W of G and for all natural numbers m, n holds  $W.\operatorname{cut}(m,n)$  is trail-like.
- (143) For every trail W of G and for every set e such that  $e \in W.\text{last}().\text{edgesInOut}()$  and  $e \notin W.\text{edges}()$  holds W.addEdge(e) is trail-like.
- (144) For every trail W of G and for every vertex v of G such that  $v \in W$ .vertices() and v is endvertex holds v = W.first() or v = W.last().
- (145) For every finite graph G and for every trail W of G holds  $len(W.edgeSeq()) \leq G.size().$
- (146) If len  $W \leq 3$ , then W is path-like.
- (147) If for all odd natural numbers m, n such that  $m \leq \operatorname{len} W$  and  $n \leq \operatorname{len} W$ and W(m) = W(n) holds m = n, then W is path-like.
- (148) Let W be a path of G. Suppose W is open. Let m, n be odd natural numbers. If m < n and  $n \leq \text{len } W$ , then  $W(m) \neq W(n)$ .
- (149) W is path-like iff W.reverse() is path-like.
- (150) For every path W of G and for all natural numbers m, n holds  $W.\operatorname{cut}(m, n)$  is path-like.
- (151) Let W be a path of G and e, v be sets. Suppose that
  - (i) e joins W.last() and v in G,
  - (ii)  $e \notin W.edges(),$
  - (iii) W is trivial or open, and
  - (iv) for every odd natural number n such that 1 < n and  $n \leq \text{len } W$  holds  $W(n) \neq v$ .

Then W.addEdge(e) is path-like.

- (152) Let W be a path of G and e, v be sets. Suppose e joins W.last() and v in G and  $v \notin W$ .vertices() and W is trivial or open. Then W.addEdge(e) is path-like.
- (153) If for every odd natural number n such that  $n \leq \operatorname{len} W$  holds  $W.\operatorname{find}(W(n)) = W.\operatorname{rfind}(W(n))$ , then W is path-like.
- (154) If for every odd natural number n such that  $n \leq \operatorname{len} W$  holds W.rfind(n) = n, then W is path-like.
- (155) For every finite graph G and for every path W of G holds  $len(W.vertexSeq()) \le G.order() + 1.$
- (156) Let G be a graph, W be a vertex-distinct walk of G, and e, v be sets. If e joins W.last() and v in G and  $v \notin W$ .vertices(), then W.addEdge(e) is vertex-distinct.
- (157) If e joins x and x in G, then G.walkOf(x, e, x) is cycle-like.
- (158) Suppose e joins x and y in G and  $e \in W_1.edges()$  and  $W_1$  is cycle-like. Then there exists a walk  $W_2$  of G such that  $W_2$  is walk from x to y and  $e \notin W_2.edges()$ .
- (159) W is a subwalk of W.
- (160) For every walk  $W_1$  of G and for every subwalk  $W_2$  of  $W_1$  holds every subwalk of  $W_2$  is a subwalk of  $W_1$ .
- (161) If  $W_1$  is a subwalk of  $W_2$ , then  $W_1$  is walk from x to y iff  $W_2$  is walk from x to y.
- (162) If  $W_1$  is a subwalk of  $W_2$ , then  $W_1$ .first() =  $W_2$ .first() and  $W_1$ .last() =  $W_2$ .last().
- (163) If  $W_1$  is a subwalk of  $W_2$ , then  $\operatorname{len} W_1 \leq \operatorname{len} W_2$ .
- (164) If  $W_1$  is a subwalk of  $W_2$ , then  $W_1.edges() \subseteq W_2.edges()$  and  $W_1.vertices() \subseteq W_2.vertices()$ .
- (165) Suppose  $W_1$  is a subwalk of  $W_2$ . Let m be an odd natural number. Suppose  $m \leq \operatorname{len} W_1$ . Then there exists an odd natural number n such that  $m \leq n$  and  $n \leq \operatorname{len} W_2$  and  $W_1(m) = W_2(n)$ .
- (166) Suppose  $W_1$  is a subwalk of  $W_2$ . Let m be an even natural number. Suppose  $1 \le m$  and  $m \le \operatorname{len} W_1$ . Then there exists an even natural number n such that  $m \le n$  and  $n \le \operatorname{len} W_2$  and  $W_1(m) = W_2(n)$ .
- (167) For every trail  $W_1$  of G such that  $W_1$  is non trivial holds there exists a path of  $W_1$  which is non trivial.
- (168) For every graph  $G_1$  and for every subgraph  $G_2$  of  $G_1$  holds every walk of  $G_2$  is a walk of  $G_1$ .
- (169) Let  $G_1$  be a graph,  $G_2$  be a subgraph of  $G_1$ , and W be a walk of  $G_1$ . If W is trivial and W.first()  $\in$  the vertices of  $G_2$ , then W is a walk of  $G_2$ .

- (170) Let  $G_1$  be a graph,  $G_2$  be a subgraph of  $G_1$ , and W be a walk of  $G_1$ . If W is non trivial and W.edges()  $\subseteq$  the edges of  $G_2$ , then W is a walk of  $G_2$ .
- (171) Let  $G_1$  be a graph,  $G_2$  be a subgraph of  $G_1$ , and W be a walk of  $G_1$ . Suppose W.vertices()  $\subseteq$  the vertices of  $G_2$  and W.edges()  $\subseteq$  the edges of  $G_2$ . Then W is a walk of  $G_2$ .
- (172) Let  $G_1$  be a non trivial graph, W be a walk of  $G_1$ , v be a vertex of  $G_1$ , and  $G_2$  be a subgraph of  $G_1$  with vertex v removed. If  $v \notin W$ .vertices(), then W is a walk of  $G_2$ .
- (173) Let  $G_1$  be a graph, W be a walk of  $G_1$ , e be a set, and  $G_2$  be a subgraph of  $G_1$  with edge e removed. If  $e \notin W$ .edges(), then W is a walk of  $G_2$ .
- (174) Let  $G_1$  be a graph,  $G_2$  be a subgraph of  $G_1$ , and x, y, e be sets. If e joins x and y in  $G_2$ , then  $G_1$ .walkOf $(x, e, y) = G_2$ .walkOf(x, e, y).
- (175) Let  $G_1$  be a graph,  $G_2$  be a subgraph of  $G_1$ ,  $W_1$  be a walk of  $G_1$ ,  $W_2$  be a walk of  $G_2$ , and e be a set. If  $W_1 = W_2$  and  $e \in W_2$ .last().edgesInOut(), then  $W_1$ .addEdge $(e) = W_2$ .addEdge(e).
- (176) Let  $G_1$  be a graph,  $G_2$  be a subgraph of  $G_1$ , and W be a walk of  $G_2$ . Then
  - (i) if W is closed, then W is a closed walk of  $G_1$ ,
  - (ii) if W is directed, then W is a directed walk of  $G_1$ ,
  - (iii) if W is trivial, then W is a trivial walk of  $G_1$ ,
  - (iv) if W is trail-like, then W is a trail-like walk of  $G_1$ ,
  - (v) if W is path-like, then W is a path-like walk of  $G_1$ , and
  - (vi) if W is vertex-distinct, then W is a vertex-distinct walk of  $G_1$ .
- (177) Let  $G_1$  be a graph,  $G_2$  be a subgraph of  $G_1$ ,  $W_1$  be a walk of  $G_1$ , and  $W_2$  be a walk of  $G_2$  such that  $W_1 = W_2$ . Then
  - (i)  $W_1$  is closed iff  $W_2$  is closed,
  - (ii)  $W_1$  is directed iff  $W_2$  is directed,
  - (iii)  $W_1$  is trivial iff  $W_2$  is trivial,
  - (iv)  $W_1$  is trail-like iff  $W_2$  is trail-like,
  - (v)  $W_1$  is path-like iff  $W_2$  is path-like, and
  - (vi)  $W_1$  is vertex-distinct iff  $W_2$  is vertex-distinct.
- (178) If  $G_1 =_G G_2$  and x is a vertex sequence of  $G_1$ , then x is a vertex sequence of  $G_2$ .
- (179) If  $G_1 =_G G_2$  and x is a edge sequence of  $G_1$ , then x is a edge sequence of  $G_2$ .
- (180) If  $G_1 =_G G_2$  and x is a walk of  $G_1$ , then x is a walk of  $G_2$ .
- (181) If  $G_1 =_G G_2$ , then  $G_1$ .walkOf $(x, e, y) = G_2$ .walkOf(x, e, y).
- (182) Let  $W_1$  be a walk of  $G_1$  and  $W_2$  be a walk of  $G_2$  such that  $G_1 =_G G_2$ and  $W_1 = W_2$ . Then

- (i)  $W_1$  is closed iff  $W_2$  is closed,
- (ii)  $W_1$  is directed iff  $W_2$  is directed,
- (iii)  $W_1$  is trivial iff  $W_2$  is trivial,
- (iv)  $W_1$  is trail-like iff  $W_2$  is trail-like,
- (v)  $W_1$  is path-like iff  $W_2$  is path-like, and
- (vi)  $W_1$  is vertex-distinct iff  $W_2$  is vertex-distinct.

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# Trees and Graph Components<sup>1</sup>

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**Summary.** In the graph framework of [11] we define connected and acyclic graphs, components of a graph, and define the notion of cut-vertex (articulation point).

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The articles [15], [8], [14], [17], [12], [18], [6], [1], [16], [7], [3], [4], [5], [9], [2], [11], [10], and [13] provide the terminology and notation for this paper.

## 1. Preliminaries

Let X be a finite set. Observe that  $2^X$  is finite. The following proposition is true

(1) For every finite set X such that  $1 < \operatorname{card} X$  there exist sets  $x_1, x_2$  such that  $x_1 \in X$  and  $x_2 \in X$  and  $x_1 \neq x_2$ .

## 2. Definitions

Let G be a graph. We say that G is connected if and only if:

(Def. 1) For all vertices u, v of G holds there exists a walk of G which is walk from u to v.

Let G be a graph. We say that G is acyclic if and only if:

(Def. 2) There exists no walk of G which is cycle-like.

Let G be a graph. We say that G is tree-like if and only if:

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(Def. 3) G is acyclic and connected.

One can verify that every graph which is trivial is also connected.

Let us note that every graph which is trivial and loopless is also tree-like.

Let us note that every graph which is acyclic is also simple.

Let us observe that every graph which is tree-like is also acyclic and connected.

Let us observe that every graph which is acyclic and connected is also treelike.

Let G be a graph and let v be a vertex of G. Observe that every subgraph of G induced by  $\{v\}$  and  $\emptyset$  is tree-like.

Let G be a graph and let v be a set. We say that G is dtree rooted at v if and only if:

(Def. 4) G is tree-like and for every vertex x of G holds there exists a dwalk of G which is walk from v to x.

Let us observe that there exists a graph which is trivial, finite, and tree-like and there exists a graph which is non trivial, finite, and tree-like.

Let G be a graph. Note that there exists a subgraph of G which is trivial, finite, and tree-like.

Let G be an acyclic graph. Observe that every subgraph of G is acyclic.

Let G be a graph and let v be a vertex of G. The functor G.reachableFrom(v) yields a non empty subset of the vertices of G and is defined as follows:

(Def. 5) For every set x holds  $x \in G$ .reachableFrom(v) iff there exists a walk of G which is walk from v to x.

Let G be a graph and let v be a vertex of G. The functor G.reachableDFrom(v) yielding a non empty subset of the vertices of G is defined by:

(Def. 6) For every set x holds  $x \in G$ .reachableDFrom(v) iff there exists a dwalk of G which is walk from v to x.

Let  $G_1$  be a graph and let  $G_2$  be a subgraph of  $G_1$ . We say that  $G_2$  is component-like if and only if:

(Def. 7)  $G_2$  is connected and it is not true that there exists a connected subgraph  $G_3$  of  $G_1$  such that  $G_2 \subset G_3$ .

Let G be a graph. Note that every subgraph of G which is component-like is also connected.

Let G be a graph and let v be a vertex of G. Note that every subgraph of G induced by G.reachableFrom(v) is component-like.

Let G be a graph. Observe that there exists a subgraph of G which is component-like.

Let G be a graph. A component of G is a component-like subgraph of G.

Let G be a graph. The functor G.componentSet() yielding a non empty family of subsets of the vertices of G is defined as follows:

(Def. 8) For every set x holds  $x \in G$ .componentSet() iff there exists a vertex v of G such that x = G.reachableFrom(v).

Let G be a graph and let X be an element of G.componentSet(). Observe that every subgraph of G induced by X is component-like.

Let G be a graph. The functor G.numComponents() yielding a cardinal number is defined by:

(Def. 9)  $G.numComponents() = \overline{G.componentSet()}$ .

Let G be a finite graph. Then G.numComponents() is a non empty natural number.

Let G be a graph and let v be a vertex of G. We say that v is cut-vertex if and only if:

(Def. 10) For every subgraph  $G_2$  of G with vertex v removed holds  $G.numComponents() < G_2.numComponents().$ 

Let G be a finite graph and let v be a vertex of G. Let us observe that v is cut-vertex if and only if:

(Def. 11) For every subgraph  $G_2$  of G with vertex v removed holds  $G.numComponents() < G_2.numComponents().$ 

Let G be a non trivial finite connected graph. Observe that there exists a vertex of G which is non cut-vertex.

Let G be a non trivial finite tree-like graph. One can check that there exists a vertex of G which is endvertex.

Let G be a non trivial finite tree-like graph and let v be an endvertex vertex of G. Observe that every subgraph of G with vertex v removed is tree-like.

Let  $G_4$  be a graph sequence. We say that  $G_4$  is connected if and only if:

(Def. 12) For every natural number n holds  $G_4 \rightarrow n$  is connected.

We say that  $G_4$  is acyclic if and only if:

(Def. 13) For every natural number n holds  $G_4 \rightarrow n$  is acyclic.

We say that  $G_4$  is tree-like if and only if:

(Def. 14) For every natural number n holds  $G_4 \rightarrow n$  is tree-like.

One can check the following observations:

- \* every graph sequence which is trivial is also connected,
- \* every graph sequence which is trivial and loopless is also tree-like,
- \* every graph sequence which is acyclic is also simple,
- \* every graph sequence which is tree-like is also acyclic and connected, and
- \* every graph sequence which is acyclic and connected is also tree-like.

Let us note that there exists a graph sequence which is halting, finite, and tree-like.

Let  $G_4$  be a connected graph sequence and let n be a natural number. Note that  $G_4 \rightarrow n$  is connected.

Let  $G_4$  be an acyclic graph sequence and let n be a natural number. Observe that  $G_4 \rightarrow n$  is acyclic.

Let  $G_4$  be a tree-like graph sequence and let n be a natural number. Note that  $G_4 \rightarrow n$  is tree-like.

#### 3. Theorems

For simplicity, we use the following convention:  $G, G_1, G_2$  are graphs, e, x, y are sets,  $v, v_1, v_2$  are vertices of G, and W is a walk of G.

We now state a number of propositions:

- (2) For every non trivial connected graph G and for every vertex v of G holds v is not isolated.
- (3) Let  $G_1$  be a non trivial graph, v be a vertex of  $G_1$ , and  $G_2$  be a subgraph of  $G_1$  with vertex v removed. Suppose  $G_2$  is connected and there exists a set e such that  $e \in v$ .edgesInOut() and e does not join v and v in  $G_1$ . Then  $G_1$  is connected.
- (4) Let  $G_1$  be a non trivial connected graph, v be a vertex of  $G_1$ , and  $G_2$  be a subgraph of  $G_1$  with vertex v removed. If v is endvertex, then  $G_2$  is connected.
- (5) Let  $G_1$  be a connected graph, W be a walk of  $G_1$ , e be a set, and  $G_2$  be a subgraph of  $G_1$  with edge e removed. If W is cycle-like and  $e \in W$ .edges(), then  $G_2$  is connected.
- (6) If there exists a vertex  $v_1$  of G such that for every vertex  $v_2$  of G holds there exists a walk of G which is walk from  $v_1$  to  $v_2$ , then G is connected.
- (7) Every trivial graph is connected.
- (8) If  $G_1 =_G G_2$  and  $G_1$  is connected, then  $G_2$  is connected.
- (9)  $v \in G.$ reachableFrom(v).
- (10) If  $x \in G$ .reachableFrom $(v_1)$  and e joins x and y in G, then  $y \in G$ .reachableFrom $(v_1)$ .
- (11) G.edgesBetween(G.reachableFrom(v)) = G.edgesInOut(G.reachableFrom(v)).
- (12) If  $v_1 \in G$ .reachableFrom $(v_2)$ , then G.reachableFrom $(v_1) = G$ .reachableFrom $(v_2)$ .
- (13) If  $v \in W$ .vertices(), then W.vertices()  $\subseteq G$ .reachableFrom(v).
- (14) Let  $G_1$  be a graph,  $G_2$  be a subgraph of  $G_1$ ,  $v_1$  be a vertex of  $G_1$ , and  $v_2$  be a vertex of  $G_2$ . If  $v_1 = v_2$ , then  $G_2$ .reachableFrom $(v_2) \subseteq G_1$ .reachableFrom $(v_1)$ .
- (15) If there exists a vertex v of G such that G.reachableFrom(v) = the vertices of G, then G is connected.

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- (16) If G is connected, then for every vertex v of G holds G.reachableFrom(v) = the vertices of G.
- (17) For every vertex  $v_1$  of  $G_1$  and for every vertex  $v_2$  of  $G_2$  such that  $G_1 =_G G_2$  and  $v_1 = v_2$  holds  $G_1$ .reachableFrom $(v_1) = G_2$ .reachableFrom $(v_2)$ .
- (18)  $v \in G.$ reachableDFrom(v).
- (19) If  $x \in G$ .reachableDFrom $(v_1)$  and e joins x to y in G, then  $y \in G$ .reachableDFrom $(v_1)$ .
- (20)  $G.reachableDFrom(v) \subseteq G.reachableFrom(v).$
- (21) Let  $G_1$  be a graph,  $G_2$  be a subgraph of  $G_1$ ,  $v_1$  be a vertex of  $G_1$ , and  $v_2$  be a vertex of  $G_2$ . If  $v_1 = v_2$ , then  $G_2$ .reachableDFrom $(v_2) \subseteq G_1$ .reachableDFrom $(v_1)$ .
- (22) For every vertex  $v_1$  of  $G_1$  and for every vertex  $v_2$  of  $G_2$  such that  $G_1 =_G G_2$  and  $v_1 = v_2$  holds  $G_1$ .reachableDFrom $(v_1) = G_2$ .reachableDFrom $(v_2)$ .
- (23) For every graph  $G_1$  and for every connected subgraph  $G_2$  of  $G_1$  such that  $G_2$  is spanning holds  $G_1$  is connected.
- (24)  $\bigcup$  (*G*.componentSet()) = the vertices of *G*.
- (25) G is connected iff G.componentSet() = {the vertices of G}.
- (26) If  $G_1 =_G G_2$ , then  $G_1$ .componentSet() =  $G_2$ .componentSet().
- (27) If  $x \in G$ .componentSet(), then x is a non empty subset of the vertices of G.
- (28) G is connected iff G.numComponents() = 1.
- (29) If  $G_1 =_G G_2$ , then  $G_1$ .numComponents() =  $G_2$ .numComponents().
- (30) G is a component of G iff G is connected.
- (31) For every component C of G holds the edges of C = G.edgesBetween(the vertices of C).
- (32) For all components  $C_1$ ,  $C_2$  of G holds the vertices of  $C_1$  = the vertices of  $C_2$  iff  $C_1 =_G C_2$ .
- (33) Let C be a component of G and v be a vertex of G. Then  $v \in$  the vertices of C if and only if the vertices of C = G.reachableFrom(v).
- (34) Let  $C_1, C_2$  be components of G and v be a set. If  $v \in$  the vertices of  $C_1$  and  $v \in$  the vertices of  $C_2$ , then  $C_1 =_G C_2$ .
- (35) Let G be a connected graph and v be a vertex of G. Then v is non cutvertex if and only if for every subgraph  $G_2$  of G with vertex v removed holds  $G_2$ .numComponents()  $\leq G$ .numComponents().
- (36) Let G be a connected graph, v be a vertex of G, and  $G_2$  be a subgraph of G with vertex v removed. If v is not cut-vertex, then  $G_2$  is connected.
- (37) Let G be a non trivial finite connected graph. Then there exist vertices  $v_1, v_2$  of G such that  $v_1 \neq v_2$  and  $v_1$  is not cut-vertex and  $v_2$  is not cut-vertex.

- (38) If v is cut-vertex, then G is non trivial.
- (39) Let  $v_1$  be a vertex of  $G_1$  and  $v_2$  be a vertex of  $G_2$ . If  $G_1 =_G G_2$  and  $v_1 = v_2$ , then if  $v_1$  is cut-vertex, then  $v_2$  is cut-vertex.
- (40) For every finite connected graph G holds  $G.order() \le G.size() + 1$ .
- (41) Every acyclic graph is simple.
- (42) Let G be an acyclic graph, W be a path of G, and e be a set. If  $e \notin W.edges()$  and  $e \in W.last().edgesInOut()$ , then W.addEdge(e) is path-like.
- (43) Let G be a non trivial finite acyclic graph. Suppose the edges of  $G \neq \emptyset$ . Then there exist vertices  $v_1$ ,  $v_2$  of G such that  $v_1 \neq v_2$  and  $v_1$  is endvertex and  $v_2$  is endvertex and  $v_2 \in G$ .reachableFrom $(v_1)$ .
- (44) If  $G_1 =_G G_2$  and  $G_1$  is acyclic, then  $G_2$  is acyclic.
- (45) Let G be a non trivial finite tree-like graph. Then there exist vertices  $v_1, v_2$  of G such that  $v_1 \neq v_2$  and  $v_1$  is endvertex and  $v_2$  is endvertex.
- (46) For every finite graph G holds G is tree-like iff G is acyclic and G.order() = G.size() + 1.
- (47) For every finite graph G holds G is tree-like iff G is connected and G.order() = G.size() + 1.
- (48) If  $G_1 =_G G_2$  and  $G_1$  is tree-like, then  $G_2$  is tree-like.
- (49) If G is dtree rooted at x, then x is a vertex of G.
- (50) If  $G_1 =_G G_2$  and  $G_1$  is dtree rooted at x, then  $G_2$  is dtree rooted at x.

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# Weighted and Labeled Graphs<sup>1</sup>

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**Summary.** In the graph framework of [17] we introduce new selectors: weights for edges and labels for both edges and vertices. We introduce also a number of tools for accessing and modifying these new fields.

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The articles [20], [19], [22], [14], [23], [9], [6], [15], [1], [18], [21], [7], [12], [10], [11], [3], [24], [4], [13], [2], [5], [8], [17], and [16] provide the notation and terminology for this paper.

#### 1. Preliminaries

Let D be a set, let  $f_1$  be a finite sequence of elements of D, and let  $f_2$  be a FinSubsequence of  $f_1$ . Then Seq  $f_2$  is a finite sequence of elements of D.

Let F be a real-yielding binary relation and let X be a set. One can check that  $F \upharpoonright X$  is real-yielding.

Next we state two propositions:

- (1) Let  $x_1, x_2, x_3, x_4, x_5, x_6, x_7, x_8, x_9, x_{10}$  be sets and p be a finite sequence. Suppose  $p = \langle x_1 \rangle ^{\frown} \langle x_2 \rangle ^{\frown} \langle x_3 \rangle ^{\frown} \langle x_4 \rangle ^{\frown} \langle x_5 \rangle ^{\frown} \langle x_6 \rangle ^{\frown} \langle x_7 \rangle ^{\frown} \langle x_8 \rangle ^{\frown} \langle x_9 \rangle ^{\frown} \langle x_{10} \rangle$ . Then len p = 10 and  $p(1) = x_1$  and  $p(2) = x_2$  and  $p(3) = x_3$  and  $p(4) = x_4$  and  $p(5) = x_5$  and  $p(6) = x_6$  and  $p(7) = x_7$  and  $p(8) = x_8$  and  $p(9) = x_9$  and  $p(10) = x_{10}$ .
- (2) Let  $f_1$  be a finite sequence of elements of  $\mathbb{R}$  and  $f_2$  be a FinSubsequence of  $f_1$ . If for every natural number i such that  $i \in \text{dom } f_1$  holds  $0 \leq f_1(i)$ , then  $\sum \text{Seq } f_2 \leq \sum f_1$ .

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#### 2. Definitions

The natural number WeightSelector is defined by:

(Def. 1) WeightSelector = 5.

The natural number ELabelSelector is defined as follows:

(Def. 2) ELabelSelector = 6.

The natural number VLabelSelector is defined as follows:

(Def. 3) VLabelSelector = 7.

Let G be a graph structure. We say that G is weighted if and only if:

(Def. 4) WeightSelector  $\in \text{dom } G$  and G(WeightSelector) is a many sorted set indexed by the edges of G.

We say that G is elabeled if and only if:

(Def. 5) ELabelSelector  $\in$  dom G and there exists a function f such that G(ELabelSelector) = f and dom  $f \subseteq$  the edges of G.

We say that G is vlabeled if and only if:

(Def. 6) VLabelSelector  $\in$  dom G and there exists a function f such that G(VLabelSelector) = f and dom  $f \subseteq$  the vertices of G.

Let us mention that there exists a graph structure which is graph-like, weighted, elabeled, and vlabeled.

A w-graph is a weighted graph. A e-graph is a elabeled graph. A v-graph is a vlabeled graph. A we-graph is a weighted elabeled graph. A wv-graph is a weighted vlabeled graph. A ev-graph is a elabeled vlabeled graph. A wev-graph is a weighted elabeled vlabeled graph.

Let G be a w-graph. The weight of G yielding a many sorted set indexed by the edges of G is defined by:

(Def. 7) The weight of G = G(WeightSelector).

Let G be a e-graph. The elabel of G yields a function and is defined by:

(Def. 8) The elabel of G = G(ELabelSelector).

Let G be a v-graph. The vlabel of G yielding a function is defined by:

(Def. 9) The vlabel of G = G(VLabelSelector).

Let G be a graph and let X be a set. One can check the following observations:

\* G.set(WeightSelector, X) is graph-like,

\* G.set(ELabelSelector, X) is graph-like, and

\* G.set(VLabelSelector, X) is graph-like.

Let G be a finite graph and let X be a set. One can check the following observations:

\* G.set(WeightSelector, X) is finite,

\* G.set(ELabelSelector, X) is finite, and

\* G.set(VLabelSelector, X) is finite.

Let G be a loopless graph and let X be a set. One can check the following observations:

\* G.set(WeightSelector, X) is loopless,

\* G.set(ELabelSelector, X) is loopless, and

\* G.set(VLabelSelector, X) is loopless.

Let G be a trivial graph and let X be a set. One can check the following observations:

- \* G.set(WeightSelector, X) is trivial,
- \* G.set(ELabelSelector, X) is trivial, and
- \* G.set(VLabelSelector, X) is trivial.

Let G be a non trivial graph and let X be a set. One can verify the following observations:

- \* G.set(WeightSelector, X) is non trivial,
- \* G.set(ELabelSelector, X) is non trivial, and
- \* G.set(VLabelSelector, X) is non trivial.

Let G be a non-multi graph and let X be a set. One can check the following observations:

- \* G.set(WeightSelector, X) is non-multi,
- \* G.set(ELabelSelector, X) is non-multi, and
- \* G.set(VLabelSelector, X) is non-multi.

Let G be a non-directed-multi graph and let X be a set. One can verify the following observations:

- \* G.set(WeightSelector, X) is non-directed-multi,
- \* G.set(ELabelSelector, X) is non-directed-multi, and
- \* G.set(VLabelSelector, X) is non-directed-multi.

Let G be a connected graph and let X be a set. One can check the following observations:

- \* G.set(WeightSelector, X) is connected,
- \* G.set(ELabelSelector, X) is connected, and
- \* G.set(VLabelSelector, X) is connected.

Let G be an acyclic graph and let X be a set. One can verify the following observations:

- \* G.set(WeightSelector, X) is acyclic,
- \* G.set(ELabelSelector, X) is acyclic, and
- \* G.set(VLabelSelector, X) is acyclic.

Let G be a w-graph and let X be a set. Observe that G.set(ELabelSelector, X) is weighted and G.set(VLabelSelector, X) is weighted.

Let G be a graph and let X be a many sorted set indexed by the edges of G. Note that G.set(WeightSelector, X) is weighted.

Let G be a graph, let  $W_1$  be a non empty set, and let W be a function from the edges of G into  $W_1$ . Note that G.set(WeightSelector, W) is weighted.

Let G be a e-graph and let X be a set. Note that G.set(WeightSelector, X) is elabeled and G.set(VLabelSelector, X) is elabeled.

Let G be a graph, let Y be a set, and let X be a partial function from the edges of G to Y. One can check that G.set(ELabelSelector, X) is elabeled.

Let G be a graph and let X be a many sorted set indexed by the edges of G. One can verify that G.set(ELabelSelector, X) is elabeled.

Let G be a v-graph and let X be a set. Note that G.set(WeightSelector, X) is vlabeled and G.set(ELabelSelector, X) is vlabeled.

Let G be a graph, let Y be a set, and let X be a partial function from the vertices of G to Y. Note that G.set(VLabelSelector, X) is vlabeled.

Let G be a graph and let X be a many sorted set indexed by the vertices of G. One can verify that G.set(VLabelSelector, X) is vlabeled.

Let G be a graph. Note that  $G.set(ELabelSelector, \emptyset)$  is elabeled and  $G.set(VLabelSelector, \emptyset)$  is vlabeled.

Let G be a graph. Note that there exists a subgraph of G which is weighted, elabeled, and vlabeled.

Let G be a w-graph and let  $G_2$  be a weighted subgraph of G. We say that  $G_2$  inherits weight if and only if:

(Def. 10) The weight of  $G_2 = (\text{the weight of } G) \upharpoonright (\text{the edges of } G_2).$ 

Let G be a e-graph and let  $G_2$  be a elabeled subgraph of G. We say that  $G_2$  inherits elabel if and only if:

(Def. 11) The elabel of  $G_2 = (\text{the elabel of } G) \upharpoonright (\text{the edges of } G_2).$ 

Let G be a v-graph and let  $G_2$  be a vlabeled subgraph of G. We say that  $G_2$  inherits vlabel if and only if:

(Def. 12) The vlabel of  $G_2 = (\text{the vlabel of } G) \upharpoonright (\text{the vertices of } G_2).$ 

Let G be a w-graph. Observe that there exists a weighted subgraph of G which inherits weight.

Let G be a e-graph. One can check that there exists a elabeled subgraph of G which inherits elabel.

Let G be a v-graph. One can verify that there exists a vlabeled subgraph of G which inherits vlabel.

Let G be a we-graph. Note that there exists a weighted elabeled subgraph of G which inherits weight and elabel.

Let G be a wv-graph. Observe that there exists a weighted vlabeled subgraph of G which inherits weight and vlabel.

Let G be a ev-graph. Observe that there exists a elabeled vlabeled subgraph of G which inherits elabel and vlabel.

Let G be a wev-graph. One can verify that there exists a weighted elabeled vlabeled subgraph of G which inherits weight, elabel, and vlabel.

Let G be a w-graph. A w-subgraph of G is a weighted subgraph of G inheriting weight.

Let G be a e-graph. A e-subgraph of G is a elabeled subgraph of G inheriting elabel.

Let G be a v-graph. A v-subgraph of G is a vlabeled subgraph of G inheriting vlabel.

Let G be a we-graph. A we-subgraph of G is a weighted elabeled subgraph of G inheriting weight and elabel.

Let G be a wv-graph. A wv-subgraph of G is a weighted vlabeled subgraph of G inheriting weight and vlabel.

Let G be a ev-graph. A ev-subgraph of G is a elabeled vlabeled subgraph of G inheriting elabel and vlabel.

Let G be a wev-graph. A wev-subgraph of G is a weighted elabeled vlabeled subgraph of G inheriting weight, elabel, and vlabel.

Let G be a graph and let V, E be sets. One can verify that there exists a subgraph of G induced by V and E which is weighted, elabeled, and vlabeled.

Let G be a w-graph and let V, E be sets. One can verify that there exists a weighted subgraph of G induced by V and E which inherits weight.

Let G be a e-graph and let V, E be sets. One can verify that there exists a elabeled subgraph of G induced by V and E which inherits elabel.

Let G be a v-graph and let V, E be sets. One can verify that there exists a vlabeled subgraph of G induced by V and E which inherits vlabel.

Let G be a we-graph and let V, E be sets. Note that there exists a weighted elabeled subgraph of G induced by V and E which inherits weight and elabel.

Let G be a wv-graph and let V, E be sets. Observe that there exists a weighted vlabeled subgraph of G induced by V and E which inherits weight and vlabel.

Let G be a ev-graph and let V, E be sets. Note that there exists a elabeled vlabeled subgraph of G induced by V and E which inherits elabel and vlabel.

Let G be a wev-graph and let V, E be sets. Observe that there exists a weighted elabeled vlabeled subgraph of G induced by V and E which inherits weight, elabel, and vlabel.

Let G be a w-graph and let V, E be sets. A induced w-subgraph of G, V, E is a weighted subgraph of G induced by V and E inheriting weight.

Let G be a e-graph and let V, E be sets. A induced e-subgraph of G, V, E is a elabeled subgraph of G induced by V and E inheriting elabel.

Let G be a v-graph and let V, E be sets. A induced v-subgraph of G, V, E is a vlabeled subgraph of G induced by V and E inheriting vlabel.

Let G be a we-graph and let V, E be sets. A induced we-subgraph of G, V, E is a weighted elabeled subgraph of G induced by V and E inheriting weight and elabel.

Let G be a wv-graph and let V, E be sets. A induced wv-subgraph of G, V, E is a weighted vlabeled subgraph of G induced by V and E inheriting weight and vlabel.

Let G be a ev-graph and let V, E be sets. A induced ev-subgraph of G, V, E is a elabeled vlabeled subgraph of G induced by V and E inheriting elabel and vlabel.

Let G be a wev-graph and let V, E be sets. A induced wev-subgraph of G, V, E is a weighted elabeled vlabeled subgraph of G induced by V and E inheriting weight, elabel, and vlabel.

Let G be a w-graph and let V be a set. A induced w-subgraph of G, V is a induced w-subgraph of G, V, G.edgesBetween(V).

Let G be a e-graph and let V be a set. A induced e-subgraph of G, V is a induced e-subgraph of G, V, G.edgesBetween(V).

Let G be a v-graph and let V be a set. A induced v-subgraph of G, V is a induced v-subgraph of G, V, G.edgesBetween(V).

Let G be a we-graph and let V be a set. A induced we-subgraph of G, V is a induced we-subgraph of G, V, G.edgesBetween(V).

Let G be a wv-graph and let V be a set. A induced wv-subgraph of G, V is a induced wv-subgraph of G, V, G.edgesBetween(V).

Let G be a ev-graph and let V be a set. A induced ev-subgraph of G, V is a induced ev-subgraph of G, V, G.edgesBetween(V).

Let G be a wev-graph and let V be a set. A induced wev-subgraph of G, V is a induced wev-subgraph of G, V, G.edgesBetween(V).

Let G be a w-graph. We say that G is real-weighted if and only if:

(Def. 13) The weight of G is real-yielding.

Let G be a w-graph. We say that G is nonnegative-weighted if and only if:

(Def. 14) rng (the weight of G)  $\subseteq \mathbb{R}_{>0}$ .

Let us note that every w-graph which is nonnegative-weighted is also real-weighted.

Let G be a e-graph. We say that G is real-elabeled if and only if:

(Def. 15) The elabel of G is real-yielding.

Let G be a v-graph. We say that G is real-vlabeled if and only if:

(Def. 16) The vlabel of G is real-yielding.

Let G be a wev-graph. We say that G is real-wev if and only if:

(Def. 17) G is real-weighted, real-elabeled, and real-vlabeled.

Let us note that every wev-graph which is real-wev is also real-weighted, real-elabeled, and real-vlabeled and every wev-graph which is real-weighted, real-elabeled, and real-vlabeled is also real-wev.

Let G be a graph and let X be a function from the edges of G into  $\mathbb{R}$ . Note that G.set(WeightSelector, X) is real-weighted.

Let G be a graph and let X be a partial function from the edges of G to  $\mathbb{R}$ . One can verify that G.set(ELabelSelector, X) is real-elabeled.

Let G be a graph and let X be a real-yielding many sorted set indexed by the edges of G. One can verify that G.set(ELabelSelector, X) is real-elabeled.

Let G be a graph and let X be a partial function from the vertices of G to  $\mathbb{R}$ . Observe that G.set(VLabelSelector, X) is real-vlabeled.

Let G be a graph and let X be a real-yielding many sorted set indexed by the vertices of G. One can verify that G.set(VLabelSelector, X) is real-vlabeled.

Let G be a graph. Observe that  $G.set(ELabelSelector, \emptyset)$  is real-elabeled and  $G.set(VLabelSelector, \emptyset)$  is real-vlabeled.

Let G be a graph, let v be a vertex of G, and let  $v_1$  be a real number. Note that  $G.set(VLabelSelector, v \mapsto v_1)$  is vlabeled.

Let G be a graph, let v be a vertex of G, and let  $v_1$  be a real number. One can verify that  $G.set(VLabelSelector, v \mapsto v_1)$  is real-vlabeled.

One can check that there exists a wev-graph which is finite, trivial, tree-like, nonnegative-weighted, and real-wev and there exists a wev-graph which is finite, non trivial, tree-like, nonnegative-weighted, and real-wev.

Let G be a finite w-graph. Note that the weight of G is finite.

Let G be a finite e-graph. Note that the elabel of G is finite.

Let G be a finite v-graph. Note that the vlabel of G is finite.

Let G be a real-weighted w-graph. Observe that the weight of G is real-yielding.

Let G be a real-elabeled e-graph. One can verify that the elabel of G is real-yielding.

Let G be a real-vlabeled v-graph. Observe that the vlabel of G is real-yielding.

Let G be a real-weighted w-graph and let X be a set. Observe that G.set(ELabelSelector, X) is real-weighted and G.set(VLabelSelector, X) is real-weighted.

Let G be a nonnegative-weighted w-graph and let X be a set. One can check that G.set(ELabelSelector, X) is nonnegative-weighted and G.set(VLabelSelector, X) is nonnegative-weighted.

Let G be a real-elabeled e-graph and let X be a set. One can verify that G.set(WeightSelector, X) is real-elabeled and G.set(VLabelSelector, X) is real-elabeled.

Let G be a real-vlabeled v-graph and let X be a set. Observe that G.set(WeightSelector, X) is real-vlabeled and G.set(ELabelSelector, X) is real-vlabeled.

Let G be a w-graph and let W be a walk of G. The functor W.weightSeq() yielding a finite sequence is defined as follows:

(Def. 18)  $\operatorname{len}(W.\operatorname{weightSeq}()) = \operatorname{len}(W.\operatorname{edgeSeq}())$  and for every natural number n such that  $1 \leq n$  and  $n \leq \operatorname{len}(W.\operatorname{weightSeq}())$  holds  $W.\operatorname{weightSeq}()(n) =$  (the weight of G)( $W.\operatorname{edgeSeq}()(n)$ ).

Let G be a real-weighted w-graph and let W be a walk of G. Then W.weightSeq() is a finite sequence of elements of  $\mathbb{R}$ .

Let G be a real-weighted w-graph and let W be a walk of G. The functor W.cost() yielding a real number is defined as follows:

(Def. 19)  $W.cost() = \sum (W.weightSeq()).$ 

Let G be a e-graph. The functor G.labeledE() yields a subset of the edges of G and is defined as follows:

(Def. 20) G.labeledE() = dom (the elabel of G).

Let G be a e-graph and let e, x be sets. The functor G.labelEdge(e, x) yielding a e-graph is defined as follows:

 $(\text{Def. 21}) \quad G.\text{labelEdge}(e, x) = \begin{cases} G.\text{set}(\text{ELabelSelector}, (\text{the elabel of } G) + \cdot (e \vdash \to x)), \\ \text{if } e \in \text{the edges of } G, \\ G, \text{ otherwise.} \end{cases}$ 

Let G be a finite e-graph and let e, x be sets. Note that G.labelEdge(e, x) is finite.

Let G be a loopless e-graph and let e, x be sets. Observe that G.labelEdge(e, x) is loopless.

Let G be a trivial e-graph and let e, x be sets. One can check that G.labelEdge(e, x) is trivial.

Let G be a non trivial e-graph and let e, x be sets. One can verify that G.labelEdge(e, x) is non trivial.

Let G be a non-multi e-graph and let e, x be sets. Observe that G.labelEdge(e, x) is non-multi.

Let G be a non-directed-multi e-graph and let e, x be sets. One can check that G.labelEdge(e, x) is non-directed-multi.

Let G be a connected e-graph and let e, x be sets. Observe that G.labelEdge(e, x) is connected.

Let G be an acyclic e-graph and let e, x be sets. Observe that G.labelEdge(e, x) is acyclic.

Let G be a we-graph and let e, x be sets. Observe that G.labelEdge(e, x) is weighted.

Let G be a ev-graph and let e, x be sets. Note that G.labelEdge(e, x) is vlabeled.

Let G be a real-weighted we-graph and let e, x be sets. Observe that G.labelEdge(e, x) is real-weighted.

Let G be a nonnegative-weighted we-graph and let e, x be sets. Observe that G.labelEdge(e, x) is nonnegative-weighted.

Let G be a real-elabeled e-graph, let e be a set, and let x be a real number. Observe that G.labelEdge(e, x) is real-elabeled.

Let G be a real-vlabeled ev-graph and let e, x be sets. Note that G.labelEdge(e, x) is real-vlabeled.

Let G be a v-graph and let v, x be sets. The functor G.labelVertex(v, x) yielding a v-graph is defined as follows:

 $(\text{Def. 22}) \quad G.\text{labelVertex}(v, x) = \begin{cases} G.\text{set}(\text{VLabelSelector}, \\ (\text{the vlabel of } G) + \cdot (v \mapsto x)), \\ \text{if } v \in \text{the vertices of } G, \\ G, \text{ otherwise.} \end{cases}$ 

Let G be a v-graph. The functor G.labeledV() yielding a subset of the vertices of G is defined as follows:

(Def. 23) G.labeledV() = dom (the vlabel of G).

Let G be a finite v-graph and let v, x be sets. One can check that G.labelVertex(v, x) is finite.

Let G be a loopless v-graph and let v, x be sets. One can check that G.labelVertex(v, x) is loopless.

Let G be a trivial v-graph and let v, x be sets. One can check that G.labelVertex(v, x) is trivial.

Let G be a non trivial v-graph and let v, x be sets. Observe that G.labelVertex(v, x) is non trivial.

Let G be a non-multi v-graph and let v, x be sets. Note that G.labelVertex(v, x) is non-multi.

Let G be a non-directed-multi v-graph and let v, x be sets. One can verify that G.labelVertex(v, x) is non-directed-multi.

Let G be a connected v-graph and let v, x be sets. Observe that G.labelVertex(v, x) is connected.

Let G be an acyclic v-graph and let v, x be sets. Note that G.labelVertex(v, x) is acyclic.

Let G be a wv-graph and let v, x be sets. One can check that G.labelVertex(v, x) is weighted.

Let G be a ev-graph and let v, x be sets. Observe that G.labelVertex(v, x) is elabeled.

Let G be a real-weighted wv-graph and let v, x be sets. Observe that G.labelVertex(v, x) is real-weighted.

Let G be a nonnegative-weighted wv-graph and let v, x be sets. Note that G.labelVertex(v, x) is nonnegative-weighted.

Let G be a real-elabeled ev-graph and let v, x be sets. Observe that G.labelVertex(v, x) is real-elabeled.

Let G be a real-vlabeled v-graph, let v be a set, and let x be a real number. Note that G.labelVertex(v, x) is real-vlabeled.

Let G be a real-weighted w-graph. Observe that every w-subgraph of G is real-weighted.

Let G be a nonnegative-weighted w-graph. Observe that every w-subgraph of G is nonnegative-weighted.

Let G be a real-elabeled e-graph. Observe that every e-subgraph of G is real-elabeled.

Let G be a real-vlabeled v-graph. Observe that every v-subgraph of G is real-vlabeled.

Let  $G_1$  be a graph sequence. We say that  $G_1$  is weighted if and only if:

(Def. 24) For every natural number x holds  $G_1 \rightarrow x$  is weighted.

We say that  $G_1$  is elabeled if and only if:

(Def. 25) For every natural number x holds  $G_1 \rightarrow x$  is elabeled.

We say that  $G_1$  is vlabeled if and only if:

(Def. 26) For every natural number x holds  $G_1 \rightarrow x$  is vlabeled.

Let us mention that there exists a graph sequence which is weighted, elabeled, and vlabeled.

A w-graph sequence is a weighted graph sequence. A e-graph sequence is a elabeled graph sequence. A v-graph sequence is a vlabeled graph sequence. A we-graph sequence is a weighted elabeled graph sequence. A wv-graph sequence is a weighted vlabeled graph sequence. A ev-graph sequence is a elabeled vlabeled graph sequence. A wev-graph sequence is a weighted elabeled vlabeled graph sequence. A wev-graph sequence is a weighted elabeled vlabeled graph sequence.

Let  $G_1$  be a w-graph sequence and let x be a natural number. One can check that  $G_1 \rightarrow x$  is weighted.

Let  $G_1$  be a e-graph sequence and let x be a natural number. One can check that  $G_1 \rightarrow x$  is elabeled.

Let  $G_1$  be a v-graph sequence and let x be a natural number. Observe that  $G_1 \rightarrow x$  is vlabeled.

Let  $G_1$  be a w-graph sequence. We say that  $G_1$  is real-weighted if and only if:

(Def. 27) For every natural number x holds  $G_1 \rightarrow x$  is real-weighted.

We say that  $G_1$  is nonnegative-weighted if and only if:

(Def. 28) For every natural number x holds  $G_1 \rightarrow x$  is nonnegative-weighted.

Let  $G_1$  be a e-graph sequence. We say that  $G_1$  is real-elabeled if and only if:

(Def. 29) For every natural number x holds  $G_1 \rightarrow x$  is real-elabeled.

Let  $G_1$  be a v-graph sequence. We say that  $G_1$  is real-vlabeled if and only if:

(Def. 30) For every natural number x holds  $G_1 \rightarrow x$  is real-vlabeled.

Let  $G_1$  be a wev-graph sequence. We say that  $G_1$  is real-wev if and only if:

(Def. 31) For every natural number x holds  $G_1 \rightarrow x$  is real-wev.

Let us note that every wev-graph sequence which is real-wev is also realweighted, real-elabeled, and real-vlabeled and every wev-graph sequence which is real-weighted, real-elabeled, and real-vlabeled is also real-wev.

Let us observe that there exists a wev-graph sequence which is halting, finite, loopless, trivial, non-multi, simple, real-wev, nonnegative-weighted, and tree-like.

Let  $G_1$  be a real-weighted w-graph sequence and let x be a natural number. One can check that  $G_1 \rightarrow x$  is real-weighted.

Let  $G_1$  be a nonnegative-weighted w-graph sequence and let x be a natural number. Observe that  $G_1 \rightarrow x$  is nonnegative-weighted.

Let  $G_1$  be a real-elabeled e-graph sequence and let x be a natural number. Note that  $G_1 \rightarrow x$  is real-elabeled.

Let  $G_1$  be a real-vlabeled v-graph sequence and let x be a natural number. One can verify that  $G_1 \rightarrow x$  is real-vlabeled.

# 3. Theorems

The following propositions are true:

- (3) WeightSelector = 5 and ELabelSelector = 6 and VLabelSelector = 7.
- (4)(i) For every w-graph G holds the weight of G = G(WeightSelector),
- (ii) for every e-graph G holds the elabel of G = G(ELabelSelector), and
- (iii) for every v-graph G holds the vlabel of G = G(VLabelSelector).
- $(6)^3$  For every e-graph G holds dom (the elabel of G)  $\subseteq$  the edges of G.
- (7) For every v-graph G holds dom (the vlabel of G)  $\subseteq$  the vertices of G.
- (8) For every graph G and for every set X holds  $G =_G G.set(WeightSelector, X)$  and  $G =_G G.set(ELabelSelector, X)$  and  $G =_G G.set(VLabelSelector, X).$
- (9) For every e-graph G and for every set X holds the elabel of G = the elabel of G.set(WeightSelector, X).
- (10) For every v-graph G and for every set X holds the vlabel of G = the vlabel of G.set(WeightSelector, X).
- (11) For every w-graph G and for every set X holds the weight of G = the weight of G.set(ELabelSelector, X).
- (12) For every v-graph G and for every set X holds the vlabel of G = the vlabel of G.set(ELabelSelector, X).

<sup>&</sup>lt;sup>3</sup>The proposition (5) has been removed.

- (13) For every w-graph G and for every set X holds the weight of G = the weight of G.set(VLabelSelector, X).
- (14) For every e-graph G and for every set X holds the elabel of G = the elabel of G.set(VLabelSelector, X).
- (15) Let  $G_3$ ,  $G_2$  be w-graphs and  $G_4$  be a w-graph. Suppose  $G_3 =_G G_2$  and the weight of  $G_3$  = the weight of  $G_2$  and  $G_3$  is a w-subgraph of  $G_4$ . Then  $G_2$  is a w-subgraph of  $G_4$ .
- (16) For every w-graph  $G_3$  and for every w-subgraph  $G_2$  of  $G_3$  holds every w-subgraph of  $G_2$  is a w-subgraph of  $G_3$ .
- (17) Let  $G_3$ ,  $G_2$  be w-graphs and  $G_4$  be a w-subgraph of  $G_3$ . Suppose  $G_3 =_G G_2$  and the weight of  $G_3$  = the weight of  $G_2$ . Then  $G_4$  is a w-subgraph of  $G_2$ .
- (18) Let  $G_3$  be a w-graph,  $G_2$  be a w-subgraph of  $G_3$ , and x be a set. If  $x \in$  the edges of  $G_2$ , then (the weight of  $G_2$ )(x) = (the weight of  $G_3$ )(x).
- (19) For every w-graph G and for every walk W of G such that W is trivial holds W.weightSeq() =  $\emptyset$ .
- (20) For every w-graph G and for every walk W of G holds len(W.weightSeq()) = W.length().
- (21) For every w-graph G and for all sets x, y, e such that e joins x and y in G holds (G.walkOf(x, e, y)).weightSeq() =  $\langle (\text{the weight of } G)(e) \rangle$ .
- (22) For every w-graph G and for every walk W of G holds W.reverse().weightSeq() = Rev(W.weightSeq()).
- (23) For every w-graph G and for all walks  $W_2$ ,  $W_3$  of G such that  $W_2$ .last() =  $W_3$ .first() holds ( $W_2$ .append( $W_3$ )).weightSeq() =  $W_2$ .weightSeq()  $\cap W_3$ .weightSeq().
- (24) Let G be a w-graph, W be a walk of G, and e be a set. If  $e \in W.last().edgesInOut()$ , then  $(W.addEdge(e)).weightSeq() = W.weightSeq() \cap \langle (the weight of G)(e) \rangle$ .
- (25) Let G be a real-weighted w-graph,  $W_2$  be a walk of G, and  $W_3$  be a subwalk of  $W_2$ . Then there exists a FinSubsequence  $w_1$  of  $W_2$ .weightSeq() such that  $W_3$ .weightSeq() = Seq  $w_1$ .
- (26) Let  $G_3$ ,  $G_2$  be w-graphs,  $W_2$  be a walk of  $G_3$ , and  $W_3$  be a walk of  $G_2$ . If  $W_2 = W_3$  and the weight of  $G_3 =$  the weight of  $G_2$ , then  $W_2$ .weightSeq() =  $W_3$ .weightSeq().
- (27) Let  $G_3$  be a w-graph,  $G_2$  be a w-subgraph of  $G_3$ ,  $W_2$  be a walk of  $G_3$ , and  $W_3$  be a walk of  $G_2$ . If  $W_2 = W_3$ , then  $W_2$ .weightSeq() =  $W_3$ .weightSeq().
- (28) For every real-weighted w-graph G and for every walk W of G such that W is trivial holds W.cost() = 0.
- (29) Let G be a real-weighted w-graph,  $v_2$ ,  $v_3$  be vertices of G, and e be a set.

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If e joins  $v_2$  and  $v_3$  in G, then  $(G.walkOf(v_2, e, v_3)).cost() = (the weight of <math>G)(e)$ .

- (30) For every real-weighted w-graph G and for every walk W of G holds W.cost() = W.reverse().cost().
- (31) For every real-weighted w-graph G and for all walks  $W_2$ ,  $W_3$  of G such that  $W_2$ .last() =  $W_3$ .first() holds ( $W_2$ .append( $W_3$ )).cost() =  $W_2$ .cost() +  $W_3$ .cost().
- (32) Let G be a real-weighted w-graph, W be a walk of G, and e be a set. If  $e \in W.last().edgesInOut()$ , then (W.addEdge(e)).cost() = W.cost()+(the weight of G)(e).
- (33) Let  $G_3$ ,  $G_2$  be real-weighted w-graphs,  $W_2$  be a walk of  $G_3$ , and  $W_3$  be a walk of  $G_2$ . If  $W_2 = W_3$  and the weight of  $G_3$  = the weight of  $G_2$ , then  $W_2.\text{cost}() = W_3.\text{cost}()$ .
- (34) Let  $G_3$  be a real-weighted w-graph,  $G_2$  be a w-subgraph of  $G_3$ ,  $W_2$  be a walk of  $G_3$ , and  $W_3$  be a walk of  $G_2$ . If  $W_2 = W_3$ , then  $W_2.cost() = W_3.cost()$ .
- (35) Let G be a nonnegative-weighted w-graph, W be a walk of G, and n be a natural number. If  $n \in \text{dom}(W.\text{weightSeq}())$ , then  $0 \leq W.\text{weightSeq}()(n)$ .
- (36) For every nonnegative-weighted w-graph G and for every walk W of G holds  $0 \le W.\operatorname{cost}()$ .
- (37) For every nonnegative-weighted w-graph G and for every walk  $W_2$  of G and for every subwalk  $W_3$  of  $W_2$  holds  $W_3.cost() \le W_2.cost()$ .
- (38) Let G be a nonnegative-weighted w-graph and e be a set. If  $e \in$  the edges of G, then  $0 \leq$  (the weight of G)(e).
- (39) Let G be a e-graph and e, x be sets. Suppose  $e \in$  the edges of G. Then the elabel of G.labelEdge(e, x) = (the elabel of G)+ $\cdot(e \mapsto x)$ .
- (40) For every e-graph G and for all sets e, x such that  $e \in$  the edges of G holds (the elabel of G.labelEdge(e, x))(e) = x.
- (41) For every e-graph G and for all sets e, x holds  $G =_G G$ .labelEdge(e, x).
- (42) For every we-graph G and for all sets e, x holds the weight of G = the weight of G.labelEdge(e, x).
- (43) For every ev-graph G and for all sets e, x holds the vlabel of G = the vlabel of G.labelEdge(e, x).
- (44) For every e-graph G and for all sets  $e_1$ ,  $e_2$ , x such that  $e_1 \neq e_2$  holds (the elabel of G.labelEdge $(e_1, x)$ ) $(e_2) =$  (the elabel of G) $(e_2)$ .
- (45) Let G be a v-graph and v, x be sets. Suppose  $v \in$  the vertices of G. Then the vlabel of G.labelVertex(v, x) = (the vlabel of G)+ $\cdot(v \mapsto x)$ .
- (46) For every v-graph G and for all sets v, x such that  $v \in$  the vertices of G holds (the vlabel of G.labelVertex(v, x))(v) = x.

- (47) For every v-graph G and for all sets v, x holds  $G =_G G$ .labelVertex(v, x).
- (48) For every wv-graph G and for all sets v, x holds the weight of G = the weight of G.labelVertex(v, x).
- (49) For every ev-graph G and for all sets v, x holds the elabel of G = the elabel of G.labelVertex(v, x).
- (50) For every v-graph G and for all sets  $v_2$ ,  $v_3$ , x such that  $v_2 \neq v_3$  holds (the vlabel of G.labelVertex $(v_2, x)$ ) $(v_3) =$  (the vlabel of G) $(v_3)$ .
- (51) For all e-graphs  $G_3$ ,  $G_2$  such that the elabel of  $G_3$  = the elabel of  $G_2$  holds  $G_3$ .labeledE() =  $G_2$ .labeledE().
- (52) For every e-graph G and for all sets e, x such that  $e \in$  the edges of G holds (G.labelEdge(e, x)).labeledE() = G.labeledE $() \cup \{e\}$ .
- (53) For every e-graph G and for all sets e, x such that  $e \in$  the edges of G holds G.labeledE()  $\subseteq$  (G.labelEdge(e, x)).labeledE().
- (54) For every finite e-graph G and for all sets e, x such that  $e \in$  the edges of G and  $e \notin G.$ labeledE() holds card((G.labelEdge(e, x)).labeledE()) = card(G.labeledE()) + 1.
- (55) For every e-graph G and for all sets  $e_1$ ,  $e_2$ , x such that  $e_2 \notin G$ .labeledE() and  $e_2 \in (G.labelEdge(e_1, x)).labeledE()$  holds  $e_1 = e_2$  and  $e_1 \in$  the edges of G.
- (56) For every ev-graph G and for all sets v, x holds G.labeledE() = (G.labelVertex(v, x)).labeledE().
- (57) For every e-graph G and for all sets e, x such that  $e \in$  the edges of G holds  $e \in (G.labelEdge(e, x)).labeledE().$
- (58) For all v-graphs  $G_3$ ,  $G_2$  such that the vlabel of  $G_3$  = the vlabel of  $G_2$  holds  $G_3$ .labeledV() =  $G_2$ .labeledV().
- (59) For every v-graph G and for all sets v, x such that  $v \in$  the vertices of G holds (G.labelVertex(v, x)).labeledV() = G.labeledV()  $\cup \{v\}$ .
- (60) For every v-graph G and for all sets v, x such that  $v \in$  the vertices of G holds G.labeledV()  $\subseteq$  (G.labelVertex(v, x)).labeledV().
- (61) For every finite v-graph G and for all sets v, x such that  $v \in$  the vertices of G and  $v \notin G$ .labeledV() holds card((G.labelVertex(v, x)).labeledV()) = card(G.labeledV()) + 1.
- (62) For every v-graph G and for all sets  $v_2$ ,  $v_3$ , x such that  $v_3 \notin G$ .labeledV() and  $v_3 \in (G.labelVertex(v_2, x)).labeledV()$  holds  $v_2 = v_3$  and  $v_2 \in$  the vertices of G.
- (63) For every ev-graph G and for all sets e, x holds G.labeledV() = (G.labelEdge(e, x)).labeledV().
- (64) For every v-graph G and for every vertex v of G and for every set x holds  $v \in (G.labelVertex(v, x)).labeledV().$

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# Correctness of Dijkstra's Shortest Path and Prim's Minimum Spanning Tree Algorithms<sup>1</sup>

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**Summary.** We prove correctness for Dijkstra's shortest path algorithm and Prim's minimum weight spanning tree algorithm at the level of graph manipulations.

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The notation and terminology used in this paper are introduced in the following articles: [25], [11], [24], [22], [28], [23], [13], [30], [10], [7], [4], [6], [14], [1], [26], [29], [8], [3], [27], [21], [19], [12], [2], [5], [9], [18], [16], [15], [20], and [17].

#### 1. Preliminaries

One can prove the following propositions:

- (1) For all functions f, g holds  $\operatorname{support}(f + g) \subseteq \operatorname{support} f \cup \operatorname{support} g$ .
- (2) For every function f and for all sets x, y holds  $\operatorname{support}(f + (x \mapsto y)) \subseteq \operatorname{support} f \cup \{x\}.$
- (3) Let A, B be sets, b be a real bag over A,  $b_1$  be a real bag over B, and  $b_2$  be a real bag over  $A \setminus B$ . If  $b = b_1 + b_2$ , then  $\sum b = \sum b_1 + \sum b_2$ .
- (4) For all sets X, x and for every real bag b over X such that dom  $b = \{x\}$  holds  $\sum b = b(x)$ .

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- (5) For every set A and for all real bags  $b_1$ ,  $b_2$  over A such that for every set x such that  $x \in A$  holds  $b_1(x) \leq b_2(x)$  holds  $\sum b_1 \leq \sum b_2$ .
- (6) For every set A and for all real bags  $b_1$ ,  $b_2$  over A such that for every set x such that  $x \in A$  holds  $b_1(x) = b_2(x)$  holds  $\sum b_1 = \sum b_2$ .
- (7) For all sets  $A_1$ ,  $A_2$  and for every real bag  $b_1$  over  $A_1$  and for every real bag  $b_2$  over  $A_2$  such that  $b_1 = b_2$  holds  $\sum b_1 = \sum b_2$ .
- (8) For all sets X, x and for every real bag b over X and for every real number y such that  $b = \text{EmptyBag } X + (x \mapsto y)$  holds  $\sum b = y$ .
- (9) Let X, x be sets,  $b_1$ ,  $b_2$  be real bags over X, and y be a real number. If  $b_2 = b_1 + (x \mapsto y)$ , then  $\sum b_2 = (\sum b_1 + y) b_1(x)$ .

# 2. DIJKSTRA'S SHORTEST PATH ALGORITHM: DEFINITIONS

Let  $G_1$  be a real-weighted w-graph, let  $G_2$  be a w-subgraph of  $G_1$ , and let v be a set. We say that  $G_2$  is mincost d-tree rooted at v if and only if the conditions (Def. 1) are satisfied.

# (Def. 1)(i) $G_2$ is tree-like, and

(ii) for every vertex x of  $G_2$  there exists a dpath  $W_2$  of  $G_2$  such that  $W_2$  is walk from v to x and for every dpath  $W_1$  of  $G_1$  such that  $W_1$  is walk from v to x holds  $W_2.cost() \leq W_1.cost()$ .

Let G be a real-weighted w-graph, let W be a dpath of G, and let x, y be sets. We say that W is mincost d-path from x to y if and only if:

(Def. 2) W is walk from x to y and for every dpath  $W_2$  of G such that  $W_2$  is walk from x to y holds  $W.cost() \le W_2.cost()$ .

Let G be a finite real-weighted w-graph and let x, y be sets. The G .mincostd-path(x, y) yielding a real number is defined as follows:

- (Def. 3)(i) There exists a dpath W of G such that W is mincost d-path from x to y and the G-mincost-d-path(x, y) = W.cost() if there exists a dwalk of G which is walk from x to y,
  - (ii) the G.mincost-d-path(x, y) = 0, otherwise.

Let G be a real-wev wev-graph. The functor DIJK : NextBestEdges(G) yielding a subset of the edges of G is defined by the condition (Def. 4).

- (Def. 4) Let  $e_1$  be a set. Then  $e_1 \in \text{DIJK}$ : NextBestEdges(G) if and only if the following conditions are satisfied:
  - (i)  $e_1$  joins a vertex from G.labeledV() to a vertex from (the vertices of  $G \setminus G.$ labeledV() in G, and
  - (ii) for every set  $e_2$  such that  $e_2$  joins a vertex from G.labeledV() to a vertex from (the vertices of G) \ G.labeledV() in G holds (the vlabel of G)((the source of G)( $e_1$ ) + (the weight of G)( $e_1$ )  $\leq$  (the vlabel of G)((the source of G)( $e_2$ )) + (the weight of G)( $e_2$ ).

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Let G be a real-wev wev-graph. The functor DIJK : Step(G) yields a realwev wev-graph and is defined by:

 $(\text{Def. 5}) \quad \text{DIJK}: \text{Step}(G) = \begin{cases} G, \text{ if } \text{DIJK}: \text{NextBestEdges}(G) = \emptyset, \\ (G.\text{labelEdge}(e, 1)).\text{labelVertex}((\text{the target of } G)(e), \\ (\text{the vlabel of } G)((\text{the source of } G)(e)) + \\ (\text{the weight of } G)(e)), \text{ otherwise.} \end{cases}$ 

Let G be a finite real-wev wev-graph. One can verify that DIJK : Step(G) is finite.

Let G be a nonnegative-weighted real-wev wev-graph. Observe that DIJK : Step(G) is nonnegative-weighted.

Let G be a real-weighted w-graph and let  $s_1$  be a vertex of G. The functor DIJK : Init $(G, s_1)$  yielding a real-wev wev-graph is defined by:

(Def. 6) DIJK : Init( $G, s_1$ ) = G.set(ELabelSelector,  $\emptyset$ ).set(VLabelSelector,  $s_1 \mapsto 0$ ).

Let G be a real-weighted w-graph and let  $s_1$  be a vertex of G. The functor DIJK : CompSeq $(G, s_1)$  yielding a real-wev wev-graph sequence is defined as follows:

(Def. 7) DIJK : CompSeq $(G, s_1) \rightarrow 0$  = DIJK : Init $(G, s_1)$  and for every natural number n holds DIJK : CompSeq $(G, s_1) \rightarrow (n + 1)$  = DIJK : Step $((DIJK : CompSeq<math>(G, s_1) \rightarrow n))$ .

Let G be a finite real-weighted w-graph and let  $s_1$  be a vertex of G. Observe that DIJK : CompSeq $(G, s_1)$  is finite.

Let G be a nonnegative-weighted w-graph and let  $s_1$  be a vertex of G. One can verify that DIJK : CompSeq $(G, s_1)$  is nonnegative-weighted.

Let G be a real-weighted w-graph and let  $s_1$  be a vertex of G. The functor DIJK :  $SSSP(G, s_1)$  yields a real-wev wev-graph and is defined by:

(Def. 8)  $\text{DIJK} : \text{SSSP}(G, s_1) = (\text{DIJK} : \text{CompSeq}(G, s_1)).\text{Result}().$ 

Let G be a finite real-weighted w-graph and let  $s_1$  be a vertex of G. One can check that DIJK :  $SSSP(G, s_1)$  is finite.

# 3. DIJKSTRA'S SHORTEST PATH ALGORITHM: THEOREMS

The following propositions are true:

- (10) Let G be a finite nonnegative-weighted w-graph, W be a dpath of G, x, y be sets, and m, n be natural numbers. Suppose W is mincost d-path from x to y. Then  $W.\operatorname{cut}(m, n)$  is mincost d-path from  $(W.\operatorname{cut}(m, n)).\operatorname{first}()$  to  $(W.\operatorname{cut}(m, n)).\operatorname{last}().$
- (11) Let G be a finite real-weighted w-graph,  $W_1$ ,  $W_2$  be dpaths of G, and x, y be sets. Suppose  $W_1$  is mincost d-path from x to y and  $W_2$  is mincost d-path from x to y. Then  $W_1.cost() = W_2.cost()$ .

- (12) Let G be a finite real-weighted w-graph, W be a dpath of G, and x, y be sets. Suppose W is mincost d-path from x to y. Then the G.mincost-d-path(x, y) = W.cost().
- (13) Let G be a finite real-wev wev-graph. Then
  - (i)  $\operatorname{card}((\operatorname{DIJK} : \operatorname{Step}(G)).\operatorname{labeledV}()) = \operatorname{card}(G.\operatorname{labeledV}())$  iff  $\operatorname{DIJK} : \operatorname{NextBestEdges}(G) = \emptyset$ , and
- (ii)  $\operatorname{card}((\operatorname{DIJK} : \operatorname{Step}(G)).\operatorname{labeledV}()) = \operatorname{card}(G.\operatorname{labeledV}()) + 1$  iff  $\operatorname{DIJK} : \operatorname{NextBestEdges}(G) \neq \emptyset.$
- (14) For every real-wev wev-graph G holds  $G =_G$  DIJK : Step(G) and the weight of G = the weight of DIJK : Step(G) and  $G.labeledE() \subseteq (DIJK : Step<math>(G)).labeledE()$  and  $G.labeledV() \subseteq (DIJK : Step<math>(G)).labeledV()$ .
- (15) For every real-weighted w-graph G and for every vertex  $s_1$  of G holds (DIJK : Init(G,  $s_1$ )).labeledV() = { $s_1$ }.
- (16) Let G be a real-weighted w-graph,  $s_1$  be a vertex of G, and i, j be natural numbers. If  $i \leq j$ , then (DIJK : CompSeq $(G, s_1) \rightarrow i$ ).labeledV()  $\subseteq$  (DIJK : CompSeq $(G, s_1) \rightarrow j$ ).labeledV() and (DIJK : CompSeq $(G, s_1) \rightarrow i$ ).labeledE()  $\subseteq$  (DIJK : CompSeq $(G, s_1) \rightarrow j$ ).labeledE().
- (17) Let G be a real-weighted w-graph,  $s_1$  be a vertex of G, and n be a natural number. Then  $G =_G \text{DIJK} : \text{CompSeq}(G, s_1) \rightarrow n$  and the weight of G = the weight of  $\text{DIJK} : \text{CompSeq}(G, s_1) \rightarrow n$ .
- (18) Let G be a finite real-weighted w-graph,  $s_1$  be a vertex of G, and n be a natural number. Then (DIJK : CompSeq(G,  $s_1$ ). $\rightarrow n$ ).labeledV()  $\subseteq$  G.reachableDFrom( $s_1$ ).
- (19) Let G be a finite real-weighted w-graph, s<sub>1</sub> be a vertex of G, and n be a natural number.
  Then DIJK : NextBestEdges((DIJK : CompSeq(G, s<sub>1</sub>)→n)) = Ø if and only if (DIJK : CompSeq(G, s<sub>1</sub>)→n).labeledV() = G.reachableDFrom(s<sub>1</sub>).
- (20) Let G be a finite real-weighted w-graph,  $s_1$  be a vertex of G, and n be a natural number. Then  $\overline{(\text{DIJK}: \text{CompSeq}(G, s_1) \rightarrow n).\text{labeledV}()} = \min(n+1, \operatorname{card}(G.\text{reachableDFrom}(s_1))).$
- (21) Let G be a finite real-weighted w-graph,  $s_1$  be a vertex of G, and n be a natural number. Then (DIJK : CompSeq $(G, s_1) \rightarrow n$ ).labeledE()  $\subseteq$  (DIJK : CompSeq $(G, s_1) \rightarrow n$ ).edgesBetween((DIJK : CompSeq $(G, s_1) \rightarrow n$ ).labeledV()).
- (22) Let G be a finite nonnegative-weighted w-graph,  $s_1$  be a vertex of G, n be a natural number, and  $G_2$  be a induced w-subgraph of G, (DIJK : CompSeq $(G, s_1) \rightarrow n$ ).labeledV(), (DIJK : CompSeq $(G, s_1) \rightarrow n$ ).labeledE(). Then
  - (i)  $G_2$  is mincost d-tree rooted at  $s_1$ , and

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- (ii) for every vertex v of G such that  $v \in (\text{DIJK} : \text{CompSeq}(G, s_1) \rightarrow n)$ .labeledV() holds the G.mincost-d-path( $s_1, v$ ) = (the vlabel of DIJK : CompSeq( $G, s_1$ ). $\rightarrow n$ )(v).
- (23) For every finite real-weighted w-graph G and for every vertex  $s_1$  of G holds DIJK : CompSeq $(G, s_1)$  is halting.

Let G be a finite real-weighted w-graph and let  $s_1$  be a vertex of G. Observe that DIJK : CompSeq $(G, s_1)$  is halting.

One can prove the following three propositions:

- (24) For every finite real-weighted w-graph G and for every vertex  $s_1$  of G holds (DIJK : CompSeq $(G, s_1)$ ).Lifespan $() + 1 = card(G.reachableDFrom(s_1)).$
- (25) For every finite real-weighted w-graph G and for every vertex  $s_1$  of G holds (DIJK : SSSP(G,  $s_1$ )).labeledV() = G.reachableDFrom( $s_1$ ).
- (26) Let G be a finite nonnegative-weighted w-graph,  $s_1$  be a vertex of G, and  $G_2$  be a induced w-subgraph of G, (DIJK :  $SSSP(G, s_1)$ ).labeledV(), (DIJK :  $SSSP(G, s_1)$ ).labeledE(). Then
  - (i)  $G_2$  is mincost d-tree rooted at  $s_1$ , and
  - (ii) for every vertex v of G such that  $v \in G$ .reachableDFrom $(s_1)$  holds  $v \in$  the vertices of  $G_2$  and the G.mincost-d-path $(s_1, v) =$  (the vlabel of DIJK : SSSP $(G, s_1))(v)$ .

4. PRIM'S ALGORITHM: PRELIMINARIES

The non empty finite subset WGraphSelectors of  $\mathbb N$  is defined as follows:

(Def. 9) WGraphSelectors =

{VertexSelector, EdgeSelector, SourceSelector, TargetSelector, WeightSelector}.

Let G be a w-graph. One can check that G.strict(WGraphSelectors) is graph-like and weighted.

Let G be a w-graph. The functor G.allWSubgraphs() yields a non empty set and is defined as follows:

(Def. 10) For every set x holds  $x \in G$ .allWSubgraphs() iff there exists a wsubgraph  $G_2$  of G such that  $x = G_2$  and dom  $G_2$  = WGraphSelectors.

Let G be a finite w-graph. One can check that G.allWSubgraphs() is finite. Let G be a w-graph and let X be a non empty subset of G.allWSubgraphs(). We see that the element of X is a w-subgraph of G.

Let G be a finite real-weighted w-graph. The functor G.cost() yields a real number and is defined by:

(Def. 11)  $G.cost() = \sum (the weight of G).$ The following propositions are true:

- (27) For every set x holds  $x \in$  WGraphSelectors iff x = VertexSelector or x = EdgeSelector or x = SourceSelector or x = TargetSelector or x = WeightSelector.
- (28) For every w-graph G holds WGraphSelectors  $\subseteq \text{dom } G$ .
- (29) For every w-graph G holds  $G =_G G$ .strict(WGraphSelectors) and the weight of G = the weight of G.strict(WGraphSelectors).
- (30) For every w-graph G holds dom(G.strict(WGraphSelectors)) = WGraphSelectors.
- (31) For every finite real-weighted w-graph G such that the edges of  $G = \emptyset$  holds G.cost() = 0.
- (32) Let  $G_1$ ,  $G_2$  be finite real-weighted w-graphs. Suppose the edges of  $G_1$  = the edges of  $G_2$  and the weight of  $G_1$  = the weight of  $G_2$ . Then  $G_1$ .cost() =  $G_2$ .cost().
- (33) Let  $G_1$  be a finite real-weighted w-graph, e be a set, and  $G_2$  be a weighted subgraph of  $G_1$  with edge e removed inheriting weight. If  $e \in$  the edges of  $G_1$ , then  $G_1.cost() = G_2.cost() +$  (the weight of  $G_1)(e)$ .
- (34) Let G be a finite real-weighted w-graph,  $V_1$  be a non empty subset of the vertices of G,  $E_1$  be a subset of G.edgesBetween $(V_1)$ ,  $G_1$  be a induced w-subgraph of G,  $V_1$ ,  $E_1$ , e be a set, and  $G_2$  be a induced w-subgraph of G,  $V_1$ ,  $E_1 \cup \{e\}$ . If  $e \notin E_1$  and  $e \in G$ .edgesBetween $(V_1)$ , then  $G_1.cost() + (the$  $weight of G)(e) = G_2.cost().$
- 5. PRIM'S MINIMUM WEIGHT SPANNING TREE ALGORITHM: DEFINITIONS

Let G be a real-weighted wv-graph. The functor PRIM : NextBestEdges(G) yields a subset of the edges of G and is defined by the condition (Def. 12).

- (Def. 12) Let  $e_1$  be a set. Then  $e_1 \in PRIM : NextBestEdges(G)$  if and only if the following conditions are satisfied:
  - (i)  $e_1$  joins a vertex from G.labeledV() and a vertex from (the vertices of  $G \setminus G.$ labeledV() in G, and
  - (ii) for every set  $e_2$  such that  $e_2$  joins a vertex from G.labeledV() and a vertex from (the vertices of G) \ G.labeledV() in G holds (the weight of G) $(e_1) \leq$  (the weight of G) $(e_2)$ .

Let G be a real-weighted w-graph. The functor PRIM : Init(G) yields a real-wev wev-graph and is defined by:

(Def. 13) PRIM : Init(G) = G.set(VLabelSelector, choose(the vertices of G) $\mapsto 1).set(ELabelSelector, \emptyset).$ 

Let G be a real-wev wev-graph. The functor PRIM : Step(G) yielding a real-wev wev-graph is defined by:

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$$(\text{Def. 14}) \quad \text{PRIM}: \text{Step}(G) = \begin{cases} G, \text{ if PRIM}: \text{NextBestEdges}(G) = \emptyset, \\ (G.\text{labelEdge}(e, 1)).\text{labelVertex}((\text{the target of } G) \\ (e), 1), \text{ if PRIM}: \text{NextBestEdges}(G) \neq \emptyset \text{ and} \\ (\text{the source of } G)(e) \in G.\text{labeledV}(), \\ (G.\text{labelEdge}(e, 1)).\text{labelVertex}((\text{the source of } G) \\ (e), 1), \text{ otherwise.} \end{cases}$$

Let G be a real-weighted w-graph. The functor PRIM : CompSeq(G) yields a real-wev wev-graph sequence and is defined by:

(Def. 15) PRIM :  $\text{CompSeq}(G) \rightarrow 0 = \text{PRIM} : \text{Init}(G)$  and for every natural number n holds PRIM :  $\text{CompSeq}(G) \rightarrow (n+1) = \text{PRIM} : \text{Step}((\text{PRIM} : \text{CompSeq}(G) \rightarrow n)).$ 

Let G be a finite real-weighted w-graph. One can check that PRIM : CompSeq(G) is finite.

Let G be a real-weighted w-graph. The functor PRIM : MST(G) yielding a real-wev wev-graph is defined as follows:

(Def. 16) PRIM : MST(G) = (PRIM : CompSeq(G)).Result().

Let G be a finite real-weighted w-graph. Observe that PRIM : MST(G) is finite.

Let  $G_1$  be a finite real-weighted w-graph and let n be a natural number. Observe that every subgraph of  $G_1$  induced by (PRIM : CompSeq $(G_1) \rightarrow n$ ).labeledV() is connected.

Let  $G_1$  be a finite real-weighted w-graph and let n be a natural number. Note that every subgraph of  $G_1$  induced by (PRIM : CompSeq $(G_1) \rightarrow n$ ).labeledV() and (PRIM : CompSeq $(G_1) \rightarrow n$ ).labeledE() is connected.

Let G be a finite connected real-weighted w-graph. Observe that there exists a w-subgraph of G which is spanning and tree-like.

Let  $G_1$  be a finite connected real-weighted w-graph and let  $G_2$  be a spanning tree-like w-subgraph of  $G_1$ . We say that  $G_2$  is min-cost if and only if:

(Def. 17) For every spanning tree-like w-subgraph  $G_3$  of  $G_1$  holds  $G_2.cost() \leq G_3.cost()$ .

Let  $G_1$  be a finite connected real-weighted w-graph. One can check that there exists a spanning tree-like w-subgraph of  $G_1$  which is min-cost.

Let G be a finite connected real-weighted w-graph. A minimum spanning tree of G is a min-cost spanning tree-like w-subgraph of G.

# 6. PRIM'S MINIMUM WEIGHT SPANNING TREE ALGORITHM: THEOREMS

One can prove the following propositions:

(35) Let  $G_1$ ,  $G_2$  be finite connected real-weighted w-graphs and  $G_3$  be a w-subgraph of  $G_1$ . Suppose  $G_3$  is a minimum spanning tree of  $G_1$  and

 $G_1 =_G G_2$  and the weight of  $G_1$  = the weight of  $G_2$ . Then  $G_3$  is a minimum spanning tree of  $G_2$ .

- (36) Let G be a finite connected real-weighted w-graph,  $G_1$  be a minimum spanning tree of G, and  $G_2$  be a w-graph. Suppose  $G_1 =_G G_2$  and the weight of  $G_1$  = the weight of  $G_2$ . Then  $G_2$  is a minimum spanning tree of G.
- (37) Let G be a real-weighted w-graph. Then
  - (i)  $G =_G \text{PRIM} : \text{Init}(G),$
  - (ii) the weight of G = the weight of PRIM : Init(G),
- (iii) the elabel of PRIM :  $Init(G) = \emptyset$ , and
- (iv) the vlabel of PRIM :  $Init(G) = choose(the vertices of G) \mapsto 1$ .
- (38) For every real-weighted w-graph G holds  $(PRIM : Init(G)).labeledV() = \{choose(the vertices of G)\}$  and  $(PRIM : Init(G)).labeledE() = \emptyset$ .
- (39) For every real-wev wev-graph G such that PRIM : NextBestEdges $(G) \neq \emptyset$  there exists a vertex v of G such that  $v \notin G$ .labeledV() and PRIM : Step(G) =

(G.labelEdge(choose(PRIM : NextBestEdges(G)), 1)).labelVertex(v, 1).

- (40) For every real-wev wev-graph G holds  $G =_G \operatorname{PRIM} : \operatorname{Step}(G)$ and the weight of G = the weight of  $\operatorname{PRIM} : \operatorname{Step}(G)$  and  $G.\operatorname{labeledE}() \subseteq (\operatorname{PRIM} : \operatorname{Step}(G)).\operatorname{labeledE}()$  and  $G.\operatorname{labeledV}() \subseteq (\operatorname{PRIM} : \operatorname{Step}(G)).\operatorname{labeledV}()$ .
- (41) Let G be a finite real-weighted w-graph and n be a natural number. Then  $G =_G \text{PRIM}: \text{CompSeq}(G) \rightarrow n$  and the weight of  $\text{PRIM}: \text{CompSeq}(G) \rightarrow n = \text{the weight of } G.$
- (42) Let G be a finite real-weighted w-graph and n be a natural number. Then (PRIM : CompSeq(G)  $\rightarrow n$ ).labeledV() is a non empty subset of the vertices of G and (PRIM : CompSeq(G)  $\rightarrow n$ ).labeledE()  $\subseteq$  G.edgesBetween((PRIM : CompSeq(G)  $\rightarrow n$ ).labeledV()).
- (43) For every finite real-weighted w-graph  $G_1$  and for every natural number n holds every subgraph of  $G_1$  induced by PRIM : CompSeq $(G_1)$ . $\rightarrow n$ .labeledV() and PRIM : CompSeq $(G_1)$ . $\rightarrow n$ .labeledE() is connected.
- (44) For every finite real-weighted w-graph  $G_1$  and for every natural number n holds every subgraph of  $G_1$  induced by PRIM : CompSeq $(G_1) \rightarrow n$ .labeledV() is connected.
- (45) For every finite real-weighted w-graph G and for every natural number n holds (PRIM : CompSeq(G)  $\rightarrow$  n).labeledV()  $\subseteq$  G.reachableFrom(choose(the vertices of G)).
- (46) Let G be a finite real-weighted w-graph and i, j be natural numbers. If  $i \leq j$ , then (PRIM : CompSeq(G) $\rightarrow i$ ).labeledV()  $\subseteq$  (PRIM : CompSeq(G) $\rightarrow j$ ).labeledV() and (PRIM : CompSeq(G) $\rightarrow i$ )

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 $.labeledE() \subseteq (PRIM : CompSeq(G) \rightarrow j).labeledE().$ 

- (47) Let G be a finite real-weighted w-graph and n be a natural number. Then PRIM : NextBestEdges((PRIM : CompSeq(G)  $\rightarrow n$ )) =  $\emptyset$  if and only if (PRIM : CompSeq(G)  $\rightarrow n$ ).labeledV() = G.reachableFrom(choose(the vertices of G)).
- (48) Let G be a finite real-weighted w-graph and n be a natural number. Then  $\operatorname{card}((\operatorname{PRIM}:\operatorname{CompSeq}(G)\to n).\operatorname{labeledV}()) = \min(n + 1, \operatorname{card}(G.\operatorname{reachableFrom}(\operatorname{choose}(\operatorname{the vertices of } G)))).$
- (49) For every finite real-weighted w-graph G holds PRIM : CompSeq(G) is halting and (PRIM : CompSeq(G)).Lifespan() + 1 =card(G.reachableFrom(choose(the vertices of G))).
- (50) For every finite real-weighted w-graph  $G_1$  and for every natural number n holds every subgraph of  $G_1$  induced by PRIM : CompSeq $(G_1)$ . $\rightarrow n$ .labeledV() and PRIM : CompSeq $(G_1)$ . $\rightarrow n$ .labeledE() is tree-like.
- (51) For every finite connected real-weighted w-graph G holds (PRIM : MST(G)).labeledV() = the vertices of G.
- (52) For every finite connected real-weighted w-graph G and for every natural number n holds (PRIM : CompSeq(G) $\rightarrow n$ ).labeledE()  $\subseteq$  (PRIM : MST(G)).labeledE().
- (53) For every finite connected real-weighted w-graph  $G_1$  holds every induced w-subgraph of  $G_1$ , PRIM : MST $(G_1)$ .labeledV(), PRIM : MST $(G_1)$ .labeledE() is a minimum spanning tree of  $G_1$ .

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# $\begin{array}{c} {\bf Correctnesss \ of \ Ford-Fulkerson's \ Maximum} \\ {\bf Flow \ Algorithm^1} \end{array}$

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**Summary.** We define and prove correctness of Ford-Fulkerson's maximum network flow algorithm at the level of graph manipulations.

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The articles [23], [21], [25], [22], [11], [27], [9], [7], [5], [13], [1], [24], [26], [8], [3], [4], [20], [18], [28], [10], [2], [6], [17], [12], [16], [14], [19], and [15] provide the notation and terminology for this paper.

## **1. PRELIMINARY THEOREMS**

Let x be a set and let y be a real number. One can verify that  $x \mapsto y$  is real-yielding.

Let x be a set and let y be a natural number. One can verify that  $x \mapsto y$  is natural-yielding.

Let f, g be real-yielding functions. Observe that f + g is real-yielding.

2. Preliminary Definitions for Ford-Fulkerson Flow Algorithm

Let G be a e-graph. We say that G is complete-elabeled if and only if: (Def. 1) dom (the elabel of G) = the edges of G.

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Let G be a graph and let X be a many sorted set indexed by the edges of G. Observe that G.set(ELabelSelector, X) is complete-elabeled.

Let G be a graph, let Y be a non empty set, and let X be a function from the edges of G into Y. One can check that G.set(ELabelSelector, X) is complete-elabeled.

Let  $G_1$  be a e-graph sequence. We say that  $G_1$  is complete-elabeled if and only if:

(Def. 2) For every natural number x holds  $G_1 \rightarrow x$  is complete-elabeled.

Let G be a w-graph. We say that G is natural-weighted if and only if:

(Def. 3) The weight of G is natural-yielding.

Let G be a e-graph. We say that G is natural-elabeled if and only if:

(Def. 4) The elabel of G is natural-yielding.

Let  $G_1$  be a w-graph sequence. We say that  $G_1$  is natural-weighted if and only if:

(Def. 5) For every natural number x holds  $G_1 \rightarrow x$  is natural-weighted.

Let  $G_1$  be a e-graph sequence. We say that  $G_1$  is natural-elabeled if and only if:

(Def. 6) For every natural number x holds  $G_1 \rightarrow x$  is natural-elabeled.

One can verify that every w-graph which is natural-weighted is also nonnegative-weighted.

Let us observe that every e-graph which is natural-elabeled is also realelabeled.

One can verify that there exists a wev-graph which is finite, trivial, tree-like, natural-weighted, natural-elabeled, complete-elabeled, and real-vlabeled.

One can verify that there exists a wev-graph sequence which is finite, naturalweighted, real-wev, natural-elabeled, and complete-elabeled.

Let  $G_1$  be a complete-elabeled e-graph sequence and let x be a natural number. Note that  $G_1 \rightarrow x$  is complete-elabeled.

Let  $G_1$  be a natural-elabeled e-graph sequence and let x be a natural number. One can verify that  $G_1 \rightarrow x$  is natural-elabeled.

Let  $G_1$  be a natural-weighted w-graph sequence and let x be a natural number. One can verify that  $G_1 \rightarrow x$  is natural-weighted.

Let G be a natural-weighted w-graph. One can check that the weight of G is natural-yielding.

Let G be a natural-elabeled e-graph. Note that the elabel of G is natural-yielding.

Let G be a complete-elabeled e-graph. Then the elabel of G is a many sorted set indexed by the edges of G.

Let G be a natural-weighted w-graph and let X be a set. Note that G.set(ELabelSelector, X) is natural-weighted and G.set(VLabelSelector, X) is
natural-weighted.

Let G be a graph and let X be a natural-yielding many sorted set indexed by the edges of G. Observe that G.set(ELabelSelector, X) is natural-elabeled.

Let G be a finite real-weighted real-elabeled complete-elabeled we-graph and let  $s_1$ ,  $s_2$  be sets. We say that G has valid flow from  $s_1$  to  $s_2$  if and only if the conditions (Def. 7) are satisfied.

- (Def. 7)(i)  $s_1$  is a vertex of G,
  - (ii)  $s_2$  is a vertex of G,
  - (iii) for every set e such that  $e \in$  the edges of G holds  $0 \leq$  (the elabel of G)(e) and (the elabel of G)(e)  $\leq$  (the weight of G)(e), and
  - (iv) for every vertex v of G such that  $v \neq s_1$  and  $v \neq s_2$  holds  $\sum((\text{the elabel of } G) \upharpoonright v.\text{edgesIn}()) = \sum((\text{the elabel of } G) \upharpoonright v.\text{edgesOut}()).$

Let G be a finite real-weighted real-elabeled complete-elabeled we-graph and let  $s_1$ ,  $s_2$  be sets. Let us assume that G has valid flow from  $s_1$  to  $s_2$ . The functor G.flow $(s_1, s_2)$  yields a real number and is defined as follows:

(Def. 8)  $G.\operatorname{flow}(s_1, s_2) = \sum ((\text{the elabel of } G) \upharpoonright G.\operatorname{edgesInto}(\{s_2\})) - \sum ((\text{the elabel of } G) \upharpoonright G.\operatorname{edgesOutOf}(\{s_2\})).$ 

Let G be a finite real-weighted real-elabeled complete-elabeled we-graph and let  $s_1$ ,  $s_2$  be sets. We say that G has maximum flow from  $s_1$  to  $s_2$  if and only if the conditions (Def. 9) are satisfied.

- (Def. 9)(i) G has valid flow from  $s_1$  to  $s_2$ , and
  - (ii) for every finite real-weighted real-elabeled complete-elabeled we-graph  $G_2$  such that  $G_2 =_G G$  and the weight of G = the weight of  $G_2$  and  $G_2$  has valid flow from  $s_1$  to  $s_2$  holds  $G_2$ .flow $(s_1, s_2) \leq G$ .flow $(s_1, s_2)$ .

Let G be a real-weighted real-elabeled wev-graph and let e be a set. We say that e is forward labeling in G if and only if the conditions (Def. 10) are satisfied.

- (Def. 10)(i)  $e \in$  the edges of G,
  - (ii) (the source of G) $(e) \in G$ .labeledV(),
  - (iii) (the target of G) $(e) \notin G$ .labeledV(), and
  - (iv) (the elabel of G)(e) < (the weight of G)(e).

Let G be a real-elabeled ev-graph and let e be a set. We say that e is backward labeling in G if and only if:

(Def. 11)  $e \in$  the edges of G and (the target of G) $(e) \in G$ .labeledV() and (the source of G) $(e) \notin G$ .labeledV() and 0 < (the elabel of G)(e).

Let G be a real-weighted real-elabeled we-graph and let W be a walk of G. We say that W is augmenting if and only if the condition (Def. 12) is satisfied.

- (Def. 12) Let n be an odd natural number such that n < len W. Then
  - (i) if W(n+1) joins W(n) to W(n+2) in G, then (the elabel of G)(W(n+1)) < (the weight of G)(W(n+1)), and

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(ii) if W(n+1) does not join W(n) to W(n+2) in G, then 0 < (the elabel of G)(W(n+1)).

Let G be a real-weighted real-elabeled we-graph. One can check that every walk of G which is trivial is also augmenting.

Let G be a real-weighted real-elabeled we-graph. Note that there exists a path of G which is vertex-distinct and augmenting.

Let G be a real-weighted real-elabeled we-graph, let W be an augmenting walk of G, and let m, n be natural numbers. Note that  $W.\operatorname{cut}(m, n)$  is augmenting.

Next we state two propositions:

- (1) Let  $G_3$ ,  $G_2$  be real-weighted real-elabeled we-graphs,  $W_1$  be a walk of  $G_3$ , and  $W_2$  be a walk of  $G_2$ . Suppose that
- (i)  $W_1$  is augmenting,
- (ii)  $G_3 =_G G_2$ ,
- (iii) the weight of  $G_3$  = the weight of  $G_2$ ,
- (iv) the elabel of  $G_3$  = the elabel of  $G_2$ , and
- (v)  $W_1 = W_2$ .

Then  $W_2$  is augmenting.

- (2) Let G be a real-weighted real-elabeled we-graph, W be an augmenting walk of G, and e, v be sets. Suppose that
- (i)  $v \notin W$ .vertices(), and
- (ii) e joins W.last() to v in G and (the elabel of G)(e) < (the weight of G)(e) or e joins v to W.last() in G and 0 < (the elabel of G)(e). Then W.addEdge(e) is augmenting.

3. Algorithm for Finding Augmenting Path in a Graph

Let G be a real-weighted real-elabeled wev-graph. The functor AP: NextBestEdges(G) yielding a subset of the edges of G is defined as follows:

(Def. 13) For every set e holds  $e \in AP$ : NextBestEdges(G) iff e is forward labeling in G or backward labeling in G.

Let G be a real-weighted real-elabeled wev-graph. The functor AP : Step(G) yields a real-weighted real-elabeled wev-graph and is defined by:

$$(\text{Def. 14}) \quad \text{AP}: \text{Step}(G) = \begin{cases} G, \text{ if AP}: \text{NextBestEdges}(G) = \emptyset, \\ G.\text{labelVertex}((\text{the source of } G)(e), e), \\ \text{ if AP}: \text{NextBestEdges}(G) \neq \emptyset \text{ and (the source of } G) \\ (e) \notin G.\text{labeledV}(), \\ G.\text{labelVertex}((\text{the target of } G)(e), e), \text{ otherwise.} \end{cases}$$

Let G be a finite real-weighted real-elabeled wev-graph. One can check that AP : Step(G) is finite.

Let G be a real-weighted real-elabeled we-graph and let  $s_1$  be a vertex of G. The functor AP : CompSeq $(G, s_1)$  yielding a real-weighted real-elabeled wevgraph sequence is defined as follows:

(Def. 15) AP : CompSeq $(G, s_1) \rightarrow 0 = G$ .set(VLabelSelector,  $s_1 \rightarrow 1$ ) and for every natural number n holds AP : CompSeq $(G, s_1) \rightarrow (n + 1) =$ AP : Step((AP : CompSeq $(G, s_1) \rightarrow n$ )).

Let G be a finite real-weighted real-elabeled we-graph and let  $s_1$  be a vertex of G. One can check that AP : CompSeq $(G, s_1)$  is finite.

The following three propositions are true:

- (3) Let G be a real-weighted real-elabeled we-graph and  $s_1$  be a vertex of G. Then
- (i)  $G =_G AP : CompSeq(G, s_1) \rightarrow 0$ ,
- (ii) the weight of G = the weight of AP : CompSeq $(G, s_1) \rightarrow 0$ ,
- (iii) the elabel of G = the elabel of AP : CompSeq $(G, s_1) \rightarrow 0$ , and
- (iv) (AP : CompSeq( $G, s_1$ ) $\rightarrow 0$ ).labeledV() = { $s_1$ }.
- (4) Let G be a real-weighted real-elabeled we-graph,  $s_1$  be a vertex of G, and i, j be natural numbers. If  $i \leq j$ , then  $(AP : CompSeq(G, s_1) \rightarrow i).labeledV() \subseteq (AP : CompSeq(G, s_1) \rightarrow j).labeledV().$
- (5) Let G be a real-weighted real-elabeled we-graph,  $s_1$  be a vertex of G, and n be a natural number. Then  $G =_G AP : CompSeq(G, s_1) \rightarrow n$  and the weight of G = the weight of AP : CompSeq(G,  $s_1$ ) $\rightarrow n$  and the elabel of G = the elabel of AP : CompSeq(G,  $s_1$ ) $\rightarrow n$ .

Let G be a real-weighted real-elabeled we-graph and let  $s_1$  be a vertex of G. The functor AP : FindAugPath $(G, s_1)$  yielding a real-weighted real-elabeled wev-graph is defined as follows:

(Def. 16) AP : FindAugPath $(G, s_1) = (AP : CompSeq(G, s_1))$ .Result().

We now state two propositions:

- (6) For every finite real-weighted real-elabeled we-graph G and for every vertex  $s_1$  of G holds AP : CompSeq $(G, s_1)$  is halting.
- (7) Let G be a finite real-weighted real-elabeled we-graph,  $s_1$  be a vertex of G, n be a natural number, and v be a set. If  $v \in (AP : CompSeq(G, s_1) \rightarrow n).labeledV()$ , then (the vlabel of  $AP : CompSeq(G, s_1) \rightarrow n)(v) = (the vlabel of AP : FindAugPath(G, s_1))(v).$

Let G be a finite real-weighted real-elabeled we-graph and let  $s_1$ ,  $s_2$  be vertices of G. The functor AP : GetAugPath $(G, s_1, s_2)$  yielding a vertex-distinct augmenting path of G is defined by:

(Def. 17)(i) AP : GetAugPath $(G, s_1, s_2)$  is walk from  $s_1$  to  $s_2$  and for every even natural number n such that  $n \in \text{dom AP}$  : GetAugPath $(G, s_1, s_2)$  holds (AP : GetAugPath $(G, s_1, s_2)$ ) $(n) = (\text{the vlabel of AP} : \text{FindAugPath}(G, s_1))$ 

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 $((AP : GetAugPath(G, s_1, s_2))(n + 1))$  if  $s_2 \in (AP : FindAugPath(G, s_1))$ .labeledV(),

(ii) AP : GetAugPath $(G, s_1, s_2) = G$ .walkOf $(s_1)$ , otherwise.

Next we state three propositions:

- (8) Let G be a real-weighted real-elabeled we-graph,  $s_1$  be a vertex of G, n be a natural number, and v be a set. Suppose  $v \in$  $(AP : CompSeq(G, s_1) \rightarrow n).labeledV()$ . Then there exists a path P of G such that P is augmenting and walk from  $s_1$  to v and P.vertices()  $\subseteq$  $(AP : CompSeq(G, s_1) \rightarrow n).labeledV().$
- (9) Let G be a finite real-weighted real-elabeled we-graph,  $s_1$  be a vertex of G, and v be a set. Then  $v \in (AP : FindAugPath(G, s_1)).labeledV()$  if and only if there exists a path of G which is augmenting and walk from  $s_1$  to v.
- (10) Let G be a finite real-weighted real-elabeled we-graph and  $s_1$  be a vertex of G. Then  $s_1 \in (AP : FindAugPath(G, s_1)).labeledV()$  and  $G =_G AP : FindAugPath(G, s_1)$  and the weight of G = the weight of AP : FindAugPath(G, s\_1) and the elabel of G = the elabel of AP : FindAugPath(G, s\_1).

## 4. Definition of Ford-Fulkerson Maximum Flow Algorithm

Let G be a real-weighted real-elabeled we-graph and let W be an augmenting walk of G. The functor W.flowSeq() yields a finite sequence of elements of  $\mathbb{R}$  and is defined by the conditions (Def. 18).

- $(Def. 18)(i) \quad dom(W.flowSeq()) = dom(W.edgeSeq()), and$ 
  - (ii) for every natural number n such that  $n \in \text{dom}(W.\text{flowSeq}())$  holds if  $W(2 \cdot n)$  joins  $W(2 \cdot n-1)$  to  $W(2 \cdot n+1)$  in G, then  $W.\text{flowSeq}()(n) = (\text{the weight of } G)(W(2 \cdot n)) (\text{the elabel of } G)(W(2 \cdot n))$  and if  $W(2 \cdot n)$  does not join  $W(2 \cdot n-1)$  to  $W(2 \cdot n+1)$  in G, then  $W.\text{flowSeq}()(n) = (\text{the elabel of } G)(W(2 \cdot n))$ .

Let G be a real-weighted real-elabeled we-graph and let W be an augmenting walk of G. The functor W.tolerance() yielding a real number is defined as follows:

(Def. 19)(i) W.tolerance()  $\in$  rng(W.flowSeq()) and for every real number k such that  $k \in$  rng(W.flowSeq()) holds W.tolerance()  $\leq k$  if W is non trivial,

(ii) W.tolerance() = 0, otherwise.

Let G be a natural-weighted natural-elabeled we-graph and let W be an augmenting walk of G. Then W.tolerance() is a natural number.

Let G be a real-weighted real-elabeled we-graph and let P be an augmenting path of G. The functor FF : PushFlow(G, P) yielding a many sorted set indexed by the edges of G is defined by the conditions (Def. 20).

- (Def. 20)(i) For every set e such that  $e \in$  the edges of G and  $e \notin P$ .edges() holds (FF : PushFlow(G, P))(e) = (the elabel of G)(e), and
  - (ii) for every odd natural number n such that  $n < \ln P$  holds if P(n+1) joins P(n) to P(n+2) in G, then (FF : PushFlow(G, P))(P(n+1)) = (the elabel of G)(P(n+1)) + P.tolerance() and if P(n+1) does not join P(n) to P(n+2) in G, then (FF : PushFlow(G, P))(P(n+1)) = (the elabel of G)(P(n+1)) P.tolerance().

Let G be a real-weighted real-elabeled we-graph and let P be an augmenting path of G. Observe that FF: PushFlow(G, P) is real-yielding.

Let G be a natural-weighted natural-elabeled we-graph and let P be an augmenting path of G. Note that FF: PushFlow(G, P) is natural-yielding.

Let G be a real-weighted real-elabeled we-graph and let P be an augmenting path of G. The functor FF: AugmentPath(G, P) yielding a real-weighted realelabeled complete-elabeled we-graph is defined as follows:

(Def. 21) FF : AugmentPath(G, P) = G.set(ELabelSelector, FF : PushFlow<math>(G, P)). Let G be a finite real-weighted real-elabeled we-graph and let P be an augmenting path of G. Observe that FF : AugmentPath(G, P) is finite.

Let G be a finite nonnegative-weighted real-elabeled we-graph and let P be an augmenting path of G. Note that FF: AugmentPath(G, P) is nonnegativeweighted.

Let G be a finite natural-weighted natural-elabeled we-graph and let P be an augmenting path of G. Note that FF: AugmentPath(G, P) is natural-weighted and natural-elabeled.

Let G be a finite real-weighted real-elabeled complete-elabeled we-graph and let  $s_2$ ,  $s_1$  be vertices of G. The functor FF :  $\text{Step}(G, s_1, s_2)$  yields a finite realweighted real-elabeled complete-elabeled we-graph and is defined by:

(Def. 22) FF : Step(
$$G, s_1, s_2$$
) =   

$$\begin{cases}
FF : AugmentPath(G, AP : GetAugPath(G, s_1, s_2)), & \text{if } s_2 \in (AP : FindAugPath(G, s_1)) \\ .labeledV(), \\ G, & \text{otherwise.} \end{cases}$$

Let G be a finite nonnegative-weighted real-elabeled complete-elabeled wegraph and let  $s_1$ ,  $s_2$  be vertices of G. One can check that FF :  $\text{Step}(G, s_1, s_2)$  is nonnegative-weighted.

Let G be a finite natural-weighted natural-elabeled complete-elabeled wegraph and let  $s_1$ ,  $s_2$  be vertices of G. One can verify that FF :  $\text{Step}(G, s_1, s_2)$  is natural-weighted and natural-elabeled.

Let G be a finite real-weighted w-graph and let  $s_1$ ,  $s_2$  be vertices of G. The functor FF : CompSeq $(G, s_1, s_2)$  yields a finite real-weighted real-elabeled complete-elabeled we-graph sequence and is defined by the conditions (Def. 23).

(Def. 23)(i) FF : CompSeq $(G, s_1, s_2) \rightarrow 0 = G$ .set(ELabelSelector, (the edges of  $G) \rightarrow 0$ ), and

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- (ii) for every natural number n there exist vertices  $s'_1$ ,  $s'_2$  of FF: CompSeq $(G, s_1, s_2) \rightarrow n$  such that  $s'_1 = s_1$  and  $s'_2 = s_2$  and FF: CompSeq $(G, s_1, s_2) \rightarrow (n+1) =$ 
  - $FF: Step(FF: CompSeq(G, s_1, s_2) \rightarrow n, s'_1, s'_2).$

Let G be a finite nonnegative-weighted w-graph and let  $s_2$ ,  $s_1$  be vertices of G. One can verify that FF : CompSeq $(G, s_1, s_2)$  is nonnegative-weighted.

Let G be a finite natural-weighted w-graph and let  $s_2$ ,  $s_1$  be vertices of G. One can check that FF : CompSeq $(G, s_1, s_2)$  is natural-weighted and naturalelabeled.

Let G be a finite real-weighted w-graph and let  $s_2$ ,  $s_1$  be vertices of G. The functor FF : MaxFlow $(G, s_1, s_2)$  yields a finite real-weighted real-elabeled complete-elabeled we-graph and is defined by:

(Def. 24) FF : MaxFlow( $G, s_1, s_2$ ) = (FF : CompSeq( $G, s_1, s_2$ )).Result().

## 5. Theorems for Ford-Fulkerson Maximum Flow Algorithm

One can prove the following propositions:

- (11) Let G be a finite real-weighted real-elabeled complete-elabeled we-graph,  $s_1, s_2$  be sets, and V be a subset of the vertices of G. Suppose G has valid flow from  $s_1$  to  $s_2$  and  $s_1 \in V$  and  $s_2 \notin V$ . Then  $G.\operatorname{flow}(s_1, s_2) = \sum((\text{the elabel of } G) \upharpoonright G.\text{edgesDBetween}(V, (\text{the vertices of } G) \setminus V)) - \sum((\text{the elabel of } G) \upharpoonright G.\text{edgesDBetween}((\text{the vertices of } G) \setminus V, V)).$
- (12) Let G be a finite real-weighted real-elabeled complete-elabeled we-graph,  $s_1, s_2$  be sets, and V be a subset of the vertices of G. Suppose G has valid flow from  $s_1$  to  $s_2$  and  $s_1 \in V$  and  $s_2 \notin V$ . Then  $G.\operatorname{flow}(s_1, s_2) \leq \sum ((\text{the weight of } G) \upharpoonright G.\operatorname{elgesDBetween}(V, (\text{the vertices of } G) \setminus V)).$
- (13) Let G be a real-weighted real-elabeled we-graph and P be an augmenting path of G. Then  $G =_G FF$ : AugmentPath(G, P) and the weight of G = the weight of FF : AugmentPath(G, P).
- (14) Let G be a finite real-weighted real-elabeled we-graph and W be an augmenting walk of G. If W is non trivial, then 0 < W.tolerance().
- (15) Let G be a finite real-weighted real-elabeled complete-elabeled we-graph,  $s_1$ ,  $s_2$  be sets, and P be an augmenting path of G. Suppose  $s_1 \neq s_2$  and G has valid flow from  $s_1$  to  $s_2$  and P is walk from  $s_1$  to  $s_2$ . Then FF : AugmentPath(G, P) has valid flow from  $s_1$  to  $s_2$ .
- (16) Let G be a finite real-weighted real-elabeled complete-elabeled we-graph,  $s_1, s_2$  be sets, and P be an augmenting path of G. Suppose  $s_1 \neq s_2$ and G has valid flow from  $s_1$  to  $s_2$  and P is walk from  $s_1$  to  $s_2$ . Then  $(G.flow(s_1, s_2)) + P.tolerance() = FF : AugmentPath(G, P).flow(s_1, s_2).$

- (17) Let G be a finite real-weighted w-graph,  $s_1, s_2$  be vertices of G, and n be a natural number. Then FF : CompSeq $(G, s_1, s_2) \rightarrow n =_G G$  and the weight of G = the weight of FF : CompSeq $(G, s_1, s_2) \rightarrow n$ .
- (18) Let G be a finite nonnegative-weighted w-graph,  $s_1$ ,  $s_2$  be vertices of G, and n be a natural number. If  $s_1 \neq s_2$ , then FF : CompSeq $(G, s_1, s_2) \rightarrow n$  has valid flow from  $s_1$  to  $s_2$ .
- (19) For every finite natural-weighted w-graph G and for all vertices  $s_1$ ,  $s_2$  of G such that  $s_1 \neq s_2$  holds FF : CompSeq $(G, s_1, s_2)$  is halting.
- (20) Let G be a finite real-weighted real-elabeled complete-elabeled we-graph and  $s_1$ ,  $s_2$  be sets such that  $s_1 \neq s_2$  and G has valid flow from  $s_1$  to  $s_2$  and there exists no augmenting path of G which is walk from  $s_1$  to  $s_2$ . Then G has maximum flow from  $s_1$  to  $s_2$ .
- (21) Let G be a finite real-weighted w-graph and  $s_1$ ,  $s_2$  be vertices of G. Then  $G =_G FF$ : MaxFlow $(G, s_1, s_2)$  and the weight of G = the weight of FF : MaxFlow $(G, s_1, s_2)$ .
- (22) Let G be a finite natural-weighted w-graph and  $s_1$ ,  $s_2$  be vertices of G. If  $s_2 \neq s_1$ , then FF : MaxFlow $(G, s_1, s_2)$  has maximum flow from  $s_1$  to  $s_2$ .

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# Properties of Connected Subsets of the Real Line

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The papers [31], [36], [3], [37], [27], [18], [9], [38], [10], [22], [14], [4], [34], [5], [39], [1], [33], [30], [2], [23], [21], [6], [20], [35], [29], [24], [28], [40], [17], [13], [12], [26], [15], [8], [11], [16], [19], [25], [32], and [7] provide the notation and terminology for this paper.

## 1. Preliminaries

Let X be a non empty set. Observe that  $\Omega_X$  is non empty.

Let us observe that every subspace of the metric space of real numbers is real-membered.

Let S be a real-membered 1-sorted structure. One can check that the carrier of S is real-membered.

One can check that there exists a real-membered set which is non empty, finite, lower bounded, and upper bounded.

We now state three propositions:

- (1) For every non empty lower bounded real-membered set X and for every closed subset Y of  $\mathbb{R}$  such that  $X \subseteq Y$  holds inf  $X \in Y$ .
- (2) For every non empty upper bounded real-membered set X and for every closed subset Y of  $\mathbb{R}$  such that  $X \subseteq Y$  holds sup  $X \in Y$ .
- (3) For all subsets X, Y of  $\mathbb{R}$  holds  $\overline{X \cup Y} = \overline{X} \cup \overline{Y}$ .

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### 2. Intervals

In the sequel a, b, r, s are real numbers.

Let us consider r, s. One can check the following observations:

- \* [r, s] is bounded,
- \* ]r, s] is bounded, and
- \* ]r, s[ is bounded.

Let us consider r, s. One can verify the following observations:

- \* [r, s] is connected,
- \* [r, s[ is connected,
- \* [r, s] is connected, and
- \* ]r, s[ is connected.

Let us observe that there exists a subset of  $\mathbb R$  which is open, bounded, connected, and non empty.

One can prove the following propositions:

- (4) If r < s, then  $\inf[r, s] = r$ .
- (5) If r < s, then  $\sup[r, s] = s$ .
- (6) If r < s, then  $\inf[r, s] = r$ .
- (7) If r < s, then  $\sup[r, s] = s$ .
- (8) If  $a \le b$  or  $r \le s$  and if [a, b] = [r, s], then a = r and b = s.
- (9) If a < b or r < s and if [a, b] = [r, s], then a = r and b = s.
- (10) If a < b or r < s and if [a, b] = [r, s], then a = r and b = s.
- (11) If a < b or r < s and if [a, b] = [r, s], then a = r and b = s.
- (12) If a < b and  $[a, b] \subseteq [r, s]$ , then  $r \leq a$  and  $b \leq s$ .
- (13) If a < b and  $[a, b] \subseteq [r, s]$ , then  $r \leq a$  and  $b \leq s$ .
- (14) If a < b and  $[a, b] \subseteq [r, s]$ , then  $r \leq a$  and  $b \leq s$ .
- (15) If a < b and  $]a, b] \subseteq ]r, s]$ , then  $r \leq a$  and  $b \leq s$ .

## 3. Halflines

One can prove the following propositions:

- (16)  $[a,b]^{c} = ]-\infty, a[\cup]b, +\infty[.$
- (17)  $]a, b[^{c} = ]-\infty, a] \cup [b, +\infty[.$
- (18)  $[a, b]^{c} = ]-\infty, a[ \cup [b, +\infty[.$
- (19)  $[a,b]^{c} = ]-\infty, a] \cup [b,+\infty[.$
- (20) If  $a \le b$ , then  $[a, b] \cap (]-\infty, a] \cup [b, +\infty[) = \{a, b\}$ .

Let us consider a. One can verify the following observations:

- \*  $]-\infty, a]$  is non lower bounded, upper bounded, and connected,
- \*  $]-\infty, a[$  is non lower bounded, upper bounded, and connected,
- \*  $[a, +\infty]$  is lower bounded, non upper bounded, and connected, and
- \*  $]a, +\infty[$  is lower bounded, non upper bounded, and connected.

The following propositions are true:

- (21)  $\sup \left[-\infty, a\right] = a.$
- (22)  $\sup \left[-\infty, a\right] = a.$
- (23)  $\inf[a, +\infty] = a.$
- (24)  $\inf ]a, +\infty [=a]$

## 4. Connectedness

Let us observe that  $\Omega_{\mathbb{R}}$  is connected, non lower bounded, and non upper bounded.

One can prove the following propositions:

- (25) For every bounded connected subset X of  $\mathbb{R}$  such that  $\inf X \in X$  and  $\sup X \in X$  holds  $X = [\inf X, \sup X]$ .
- (26) For every bounded subset X of  $\mathbb{R}$  such that  $\inf X \notin X$  holds  $X \subseteq [\inf X, \sup X]$ .
- (27) For every bounded connected subset X of  $\mathbb{R}$  such that  $\inf X \notin X$  and  $\sup X \in X$  holds  $X = [\inf X, \sup X]$ .
- (28) For every bounded subset X of  $\mathbb{R}$  such that  $\sup X \notin X$  holds  $X \subseteq [\inf X, \sup X[.$
- (29) For every bounded connected subset X of  $\mathbb{R}$  such that  $\inf X \in X$  and  $\sup X \notin X$  holds  $X = [\inf X, \sup X[.$
- (30) For every bounded subset X of  $\mathbb{R}$  such that  $\inf X \notin X$  and  $\sup X \notin X$  holds  $X \subseteq \inf X, \sup X[$ .
- (31) For every non empty bounded connected subset X of  $\mathbb{R}$  such that  $\inf X \notin X$  and  $\sup X \notin X$  holds  $X = \inf X, \sup X[$ .
- (32) For every subset X of  $\mathbb{R}$  such that X is upper bounded holds  $X \subseteq ]-\infty, \sup X]$ .
- (33) For every connected subset X of  $\mathbb{R}$  such that X is not lower bounded and X is upper bounded and  $\sup X \in X$  holds  $X = ]-\infty, \sup X]$ .
- (34) For every subset X of  $\mathbb{R}$  such that X is upper bounded and  $\sup X \notin X$  holds  $X \subseteq ]-\infty$ ,  $\sup X[$ .
- (35) For every non empty connected subset X of  $\mathbb{R}$  such that X is not lower bounded and X is upper bounded and  $\sup X \notin X$  holds  $X = ]-\infty, \sup X[$ .
- (36) For every subset X of  $\mathbb{R}$  such that X is lower bounded holds  $X \subseteq [\inf X, +\infty]$ .

- (37) For every connected subset X of  $\mathbb{R}$  such that X is lower bounded and X is not upper bounded and  $inf X \in X$  holds  $X = [inf X, +\infty]$ .
- (38) For every subset X of  $\mathbb{R}$  such that X is lower bounded and  $\inf X \notin X$  holds  $X \subseteq [\inf X, +\infty]$ .
- (39) For every non empty connected subset X of  $\mathbb{R}$  such that X is lower bounded and X is not upper bounded and  $\inf X \notin X$  holds  $X = [\inf X, +\infty[.$
- (40) For every connected subset X of  $\mathbb{R}$  such that X is not upper bounded and X is not lower bounded holds  $X = \mathbb{R}$ .
- (41) Let X be a connected subset of  $\mathbb{R}$ . Then X is empty or  $X = \mathbb{R}$  or there exists a such that  $X = ]-\infty, a]$  or there exists a such that  $X = ]-\infty, a[$  or there exists a such that  $X = [a, +\infty[$  or there exists a such that  $X = ]a, +\infty[$  or there exist a, b such that  $a \le b$  and X = [a, b] or there exist a, b such that a < b and X = [a, b] or there exist a, b such that a < b and X = [a, b] or there exist a, b such that a < b and X = [a, b] or there exist a, b such that a < b and X = [a, b] or there exist a, b such that a < b and X = [a, b].
- (42) For every non empty connected subset X of  $\mathbb{R}$  such that  $r \notin X$  holds  $r \leq \inf X$  or  $\sup X \leq r$ .
- (43) Let X, Y be non empty bounded connected subsets of  $\mathbb{R}$ . Suppose  $\inf X \leq \inf Y$  and  $\sup Y \leq \sup X$  and if  $\inf X = \inf Y$  and  $\inf Y \in Y$ , then  $\inf X \in X$  and if  $\sup X = \sup Y$  and  $\sup Y \in Y$ , then  $\sup X \in X$ . Then  $Y \subseteq X$ .

Let us observe that there exists a subset of  $\mathbb{R}$  which is open, closed, connected, non empty, and non bounded.

## 5. $\mathbb{R}^1$

Next we state several propositions:

- (44) For every subset X of  $\mathbb{R}^1$  such that  $a \leq b$  and X = [a, b] holds Fr  $X = \{a, b\}$ .
- (45) For every subset X of  $\mathbb{R}^1$  such that a < b and X = ]a, b[ holds Fr  $X = \{a, b\}$ .
- (46) For every subset X of  $\mathbb{R}^1$  such that a < b and X = [a, b] holds Fr  $X = \{a, b\}$ .
- (47) For every subset X of  $\mathbb{R}^1$  such that a < b and X = ]a, b] holds Fr  $X = \{a, b\}$ .
- (48) For every subset X of  $\mathbb{R}^1$  such that X = [a, b] holds Int X = [a, b].
- (49) For every subset X of  $\mathbb{R}^1$  such that X = [a, b] holds Int X = [a, b].
- (50) For every subset X of  $\mathbb{R}^1$  such that X = [a, b] holds Int X = [a, b].
- (51) For every subset X of  $\mathbb{R}^1$  such that X = [a, b] holds Int X = [a, b].

Let X be a convex subset of  $\mathbb{R}^1$ . Observe that  $\mathbb{R}^1 \upharpoonright X$  is convex.

Let A be a connected subset of  $\mathbb{R}$ . One can check that  $R^1A$  is convex. We now state the proposition

(52) Let X be a subset of  $\mathbb{R}^1$  and Y be a subset of  $\mathbb{R}$ . If X = Y, then X is connected iff Y is connected.

## 6. TOPOLOGY OF CLOSED INTERVALS

Let us consider r. Note that  $[r, r]_{T}$  is trivial.

The following four propositions are true:

- (53) If  $r \leq s$ , then every subset of  $[r, s]_{T}$  is a bounded subset of  $\mathbb{R}$ .
- (54) If  $r \leq s$ , then for every subset X of  $[r, s]_T$  such that X = [a, b] and r < a and  $b \leq s$  holds Int X = [a, b].
- (55) If  $r \leq s$ , then for every subset X of  $[r, s]_T$  such that X = ]a, b] and  $r \leq a$  and b < s holds Int X = ]a, b[.
- (56) Let X be a subset of  $[r, s]_T$  and Y be a subset of  $\mathbb{R}$ . If X = Y, then X is connected iff Y is connected.

Let T be a topological space. Observe that there exists a subset of T which is open, closed, and connected.

Let T be a non empty connected topological space. Observe that there exists a subset of T which is non empty, open, closed, and connected.

We now state the proposition

- (57) Suppose  $r \leq s$ . Let X be an open connected subset of  $[r, s]_{T}$ . Then
  - (i) X is empty, or
  - (ii) X = [r, s], or
- (iii) there exists a real number a such that r < a and  $a \le s$  and X = [r, a[, or
- (iv) there exists a real number a such that  $r \leq a$  and a < s and X = ]a, s], or
- (v) there exist real numbers a, b such that  $r \leq a$  and a < b and  $b \leq s$  and X = ]a, b[.

## 7. MINIMAL COVER OF INTERVALS

Next we state three propositions:

- (58) Let T be a 1-sorted structure and F be a family of subsets of T. Then F is a cover of T if and only if F is a cover of  $\Omega_T$ .
- (59) Let T be a 1-sorted structure, F be a finite family of subsets of T, and  $F_1$  be a family of subsets of T. Suppose F is a cover of T and  $F_1 = F \setminus \{X; X\}$

ranges over subsets of  $T: X \in F \land \bigvee_{Y: \text{subset of } T} (Y \in F \land X \subseteq Y \land X \neq Y)$ . Then  $F_1$  is a cover of T.

(60) Let S be a trivial non empty 1-sorted structure, s be a point of S, and F be a family of subsets of S. If F is a cover of S, then  $\{s\} \in F$ .

Let T be a topological structure and let F be a family of subsets of T. We say that F is connected if and only if:

(Def. 1) For every subset X of T such that  $X \in F$  holds X is connected.

Let T be a topological space. Note that there exists a family of subsets of T which is non empty, open, closed, and connected.

In the sequel n, m are natural numbers and F is a family of subsets of  $[r, s]_{T}$ . The following two propositions are true:

- (61) Let L be a topological space and G,  $G_1$  be families of subsets of L. Suppose G is a cover of L and finite. Let  $A_1$  be a set such that  $G_1 = G \setminus \{X; X \text{ ranges over subsets of } L: X \in G \land \bigvee_{Y: \text{subset of } L} (Y \in G \land X \subseteq Y \land X \neq Y)\}$  and  $A_1 = \{C; C \text{ ranges over families of subsets of } L: C \text{ is a cover of } L \land C \subseteq G_1\}$ . Then  $A_1$  has the lower Zorn property w.r.t.  $\subseteq_{(A_1)}$ .
- (62) Let L be a topological space and G,  $A_1$  be sets. Suppose  $A_1 = \{C; C \text{ ranges over families of subsets of } L: C \text{ is a cover of } L \land C \subseteq G\}$ . Let M be a set. Suppose M is minimal in  $\subseteq_{(A_1)}$  and  $M \in \text{field}(\subseteq_{(A_1)})$ . Let  $A_4$  be a subset of L. Suppose  $A_4 \in M$ . Then it is not true that there exist subsets  $A_2$ ,  $A_3$  of L such that  $A_2 \in M$  and  $A_3 \in M$  and  $A_4 \subseteq A_2 \cup A_3$  and  $A_4 \neq A_2$  and  $A_4 \neq A_3$ .

Let r, s be real numbers and let F be a family of subsets of  $[r, s]_{\mathrm{T}}$ . Let us assume that F is a cover of  $[r, s]_{\mathrm{T}} F$  is open F is connected and  $r \leq s$ . A finite sequence of elements of  $2^{\mathbb{R}}$  is said to be an interval cover of F if it satisfies the conditions (Def. 2).

- (Def. 2)(i) rng it  $\subseteq F$ ,
  - (ii)  $\bigcup$  rng it = [r, s],
  - (iii) for every natural number n such that  $1 \leq n$  holds if  $n \leq \text{len it}$ , then it<sub>n</sub> is non empty and if  $n + 1 \leq \text{len it}$ , then  $\inf(\text{it}_n) \leq \inf(\text{it}_{n+1})$  and  $\sup(\text{it}_n) \leq \sup(\text{it}_{n+1})$  and  $\inf(\text{it}_{n+1}) < \sup(\text{it}_n)$  and if  $n + 2 \leq \text{len it}$ , then  $\sup(\text{it}_n) \leq \inf(\text{it}_{n+2})$ ,
  - (iv) if  $[r, s] \in F$ , then it =  $\langle [r, s] \rangle$ , and
  - (v) if  $[r, s] \notin F$ , then there exists a real number p such that r < p and  $p \leq s$  and it(1) = [r, p[ and there exists a real number p such that  $r \leq p$  and p < s and it(len it) = ]p, s] and for every natural number n such that 1 < n and n < len it there exist real numbers p, q such that  $r \leq p$  and p < q and  $q \leq s$  and it(n) = ]p, q[.

We now state the proposition

(63) If F is a cover of  $[r, s]_T$ , open, and connected and  $r \leq s$  and  $[r, s] \in F$ , then  $\langle [r, s] \rangle$  is an interval cover of F.

In the sequel C denotes an interval cover of F.

One can prove the following propositions:

- (64) Let F be a family of subsets of  $[r, r]_T$  and C be an interval cover of F. If F is a cover of  $[r, r]_T$ , open, and connected, then  $C = \langle [r, r] \rangle$ .
- (65) If F is a cover of  $[r, s]_T$ , open, and connected and  $r \leq s$ , then  $1 \leq \text{len } C$ .
- (66) If F is a cover of  $[r, s]_{\mathrm{T}}$ , open, and connected and  $r \leq s$  and len C = 1, then  $C = \langle [r, s] \rangle$ .
- (67) If F is a cover of  $[r, s]_{\mathrm{T}}$ , open, and connected and  $r \leq s$  and  $n \in \mathrm{dom}\,C$ and  $m \in \mathrm{dom}\,C$  and n < m, then  $\inf(C_n) \leq \inf(C_m)$ .
- (68) If F is a cover of  $[r, s]_{\mathrm{T}}$ , open, and connected and  $r \leq s$  and  $n \in \mathrm{dom}\,C$ and  $m \in \mathrm{dom}\,C$  and n < m, then  $\sup(C_n) \leq \sup(C_m)$ .
- (69) If F is a cover of  $[r, s]_{\mathrm{T}}$ , open, and connected and  $r \leq s$  and  $1 \leq n$  and  $n+1 \leq \mathrm{len} C$ , then  $\mathrm{linf}(C_{n+1}), \mathrm{sup}(C_n)[$  is non empty.
- (70) If F is a cover of  $[r, s]_{T}$ , open, and connected and  $r \leq s$ , then  $\inf(C_1) = r$ .
- (71) If F is a cover of  $[r, s]_{T}$ , open, and connected and  $r \leq s$ , then  $r \in C_1$ .
- (72) If F is a cover of  $[r, s]_{\mathrm{T}}$ , open, and connected and  $r \leq s$ , then  $\sup(C_{\mathrm{len}\,C}) = s$ .
- (73) If F is a cover of  $[r, s]_{\mathrm{T}}$ , open, and connected and  $r \leq s$ , then  $s \in C_{\mathrm{len} C}$ .

Let r, s be real numbers, let F be a family of subsets of  $[r, s]_{T}$ , and let C be an interval cover of F. Let us assume that F is a cover of  $[r, s]_{T}$  F is open F is connected and  $r \leq s$ . A finite sequence of elements of  $\mathbb{R}$  is said to be a chain of rivets in interval cover C if it satisfies the conditions (Def. 3).

- $(Def. 3)(i) \quad len it = len C + 1,$ 
  - (ii) it(1) = r,
  - (iii) it(len it) = s, and
  - (iv) for every natural number n such that  $1 \le n$  and n+1 < len it holds $\operatorname{it}(n+1) \in \operatorname{jinf}(C_{n+1}), \sup(C_n)[.$

In the sequel G denotes a chain of rivets in interval cover C.

One can prove the following propositions:

- (74) If F is a cover of  $[r, s]_{T}$ , open, and connected and  $r \leq s$ , then  $2 \leq \text{len } G$ .
- (75) If F is a cover of  $[r, s]_{\mathrm{T}}$ , open, and connected and  $r \leq s$  and len C = 1, then  $G = \langle r, s \rangle$ .
- (76) If F is a cover of  $[r, s]_{T}$ , open, and connected and  $r \leq s$  and  $1 \leq n$  and  $n+1 < \operatorname{len} G$ , then  $G(n+1) < \sup(C_n)$ .
- (77) If F is a cover of  $[r, s]_T$ , open, and connected and  $r \leq s$  and 1 < n and  $n \leq \text{len } C$ , then  $\inf(C_n) < G(n)$ .

- (78) If F is a cover of  $[r, s]_T$ , open, and connected and  $r \leq s$  and  $1 \leq n$  and  $n < \operatorname{len} C$ , then  $G(n) \leq \operatorname{inf}(C_{n+1})$ .
- (79) If F is a cover of  $[r, s]_{T}$ , open, and connected and r < s, then G is increasing.
- (80) If F is a cover of  $[r, s]_T$ , open, and connected and  $r \le s$  and  $1 \le n$  and  $n < \operatorname{len} G$ , then  $[G(n), G(n+1)] \subseteq C(n)$ .

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## The Fundamental Group of the Circle

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**Summary.** The article formalizes a proof of the theorem counting the fundamental group of a circle taken from [18]. The last result describes an isomorphism between the additive group of integers and the fundamental group of a simple closed curve.

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The notation and terminology used in this paper have been introduced in the following articles: [38], [10], [44], [2], [45], [33], [7], [1], [46], [9], [27], [8], [6], [40], [12], [3], [37], [19], [41], [26], [4], [34], [28], [32], [42], [36], [43], [20], [35], [39], [11], [30], [31], [29], [22], [21], [14], [13], [5], [15], [47], [16], [17], [25], [23], and [24].

### 1. Preliminaries

Let us observe that every element of  $\mathbb{Z}^+$  is integer.

Let us note that  $\mathbb{Z}^+$  is infinite.

Let S be an infinite 1-sorted structure. Note that the carrier of S is infinite. In the sequel a, r, s denote real numbers.

One can prove the following propositions:

(1) If  $r \leq s$  and 0 < a, then for every point p of  $[r, s]_{M}$  holds Ball(p, a) = [r, s] or Ball(p, a) = [r, p+a[ or Ball(p, a) = ]p-a, s] or Ball(p, a) = ]p-a, p+a[.

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- (2) Suppose  $r \leq s$ . Then there exists a basis B of  $[r, s]_T$  such that
- (i) there exists a many sorted set f indexed by  $[r, s]_T$  such that for every point y of  $[r, s]_M$  holds  $f(y) = \{\text{Ball}(y, \frac{1}{n}); n \text{ ranges over natural numbers:} n \neq 0\}$  and  $B = \bigcup f$ , and
- (ii) for every subset X of  $[r, s]_T$  such that  $X \in B$  holds X is connected.
- (3) For every topological structure T and for every subset A of T and for every point t of T such that  $t \in A$  holds  $skl(t, A) \subseteq A$ .

Let T be a topological space and let A be an open subset of T. Observe that  $T \upharpoonright A$  is open.

Next we state several propositions:

- (4) Let T be a topological space, S be a subspace of T, A be a subset of T, and B be a subset of S. If A = B, then  $T \upharpoonright A = S \upharpoonright B$ .
- (5) Let S, T be topological spaces, A, B be subsets of T, and C, D be subsets of S. Suppose that
- (i) the topological structure of S = the topological structure of T,
- (ii) A = C,
- (iii) B = D, and
- (iv) A and B are separated. Then C and D are separated.
- (6) Let S, T be topological spaces. Suppose the topological structure of S = the topological structure of T and S is connected. Then T is connected.
- (7) Let S, T be topological spaces, A be a subset of S, and B be a subset of T. Suppose the topological structure of S = the topological structure of T and A = B and A is connected. Then B is connected.
- (8) Let S, T be non empty topological spaces, s be a point of S, t be a point of T, and A be a neighbourhood of s. Suppose the topological structure of S = the topological structure of T and s = t. Then A is a neighbourhood of t.
- (9) Let S, T be non empty topological spaces, A be a subset of S, B be a subset of T, and N be a neighbourhood of A. Suppose the topological structure of S = the topological structure of T and A = B. Then N is a neighbourhood of B.
- (10) Let S, T be non empty topological spaces, A, B be subsets of T, and f be a map from S into T. Suppose f is a homeomorphism and A is a component of B. Then  $f^{-1}(A)$  is a component of  $f^{-1}(B)$ .

#### 2. Local Connectedness

The following propositions are true:

- (11) Let T be a non empty topological space, S be a non empty subspace of T, A be a non empty subset of T, and B be a non empty subset of S. If A = B and A is locally connected, then B is locally connected.
- (12) Let S, T be non empty topological spaces. Suppose the topological structure of S = the topological structure of T and S is locally connected. Then T is locally connected.
- (13) For every non empty topological space T holds T is locally connected iff  $\Omega_T$  is locally connected.
- (14) Let T be a non empty topological space and S be a non empty open subspace of T. If T is locally connected, then S is locally connected.
- (15) Let S, T be non empty topological spaces. Suppose S and T are homeomorphic and S is locally connected. Then T is locally connected.
- (16) Let T be a non empty topological space. Given a basis B of T such that let X be a subset of T. If  $X \in B$ , then X is connected. Then T is locally connected.
- (17) If  $r \leq s$ , then  $[r, s]_{T}$  is locally connected.

Let us mention that  $\mathbb{I}$  is locally connected.

Let A be a non empty open subset of  $\mathbb{I}$ . Observe that  $\mathbb{I} \upharpoonright A$  is locally connected.

## 3. Some Useful Functions

Let r be a real number. The functor ExtendInt r yielding a map from  $\mathbb I$  into  $\mathbb R^1$  is defined as follows:

- (Def. 1) For every point x of I holds  $(\text{ExtendInt } r)(x) = r \cdot x$ . Let r be a real number. One can check that ExtendInt r is continuous. Let r be a real number. Then ExtendInt r is a path from  $R^{10}$  to  $R^{1}r$ . Let S, T, Y be non empty topological spaces, let H be a map from [S, T]into Y, and let t be a point of T. The functor Prj1(t, H) yields a map from S into Y and is defined by:
- (Def. 2) For every point s of S holds (Prj1(t, H))(s) = H(s, t).

Let S, T, Y be non empty topological spaces, let H be a map from [S, T] into Y, and let s be a point of S. The functor Prj2(s, H) yields a map from T into Y and is defined as follows:

(Def. 3) For every point t of T holds (Prj2(s, H))(t) = H(s, t).

Let S, T, Y be non empty topological spaces, let H be a continuous map from [S, T] into Y, and let t be a point of T. Note that Prj1(t, H) is continuous.

Let S, T, Y be non empty topological spaces, let H be a continuous map from [S, T] into Y, and let s be a point of S. One can check that Prj2(s, H)is continuous.

One can prove the following two propositions:

- (18) Let T be a non empty topological space, a, b be points of T, P, Q be paths from a to b, H be a homotopy between P and Q, and t be a point of I. If H is continuous, then Prj1(t, H) is continuous.
- (19) Let T be a non empty topological space, a, b be points of T, P, Q be paths from a to b, H be a homotopy between P and Q, and s be a point of I. If H is continuous, then Prj2(s, H) is continuous.

Let r be a real number. The functor cLoop r yielding a map from  $\mathbb{I}$  into TopUnitCircle 2 is defined as follows:

- (Def. 4) For every point x of I holds  $(c \operatorname{Loop} r)(x) = [\cos(2 \cdot \pi \cdot r \cdot x), \sin(2 \cdot \pi \cdot r \cdot x)]$ . The following proposition is true
  - (20)  $\operatorname{cLoop} r = \operatorname{CircleMap} \cdot \operatorname{ExtendInt} r.$

Let n be an integer. Then cLoop n is a loop of c[10].

4. Main Theorems

Next we state four propositions:

- (21) Let  $U_1$  be a family of subsets of TopUnitCircle 2. Suppose  $U_1$  is a cover of TopUnitCircle 2 and open. Let Y be a non empty topological space, F be a continuous map from  $[Y, \mathbb{I}]$  into TopUnitCircle 2, and y be a point of Y. Then there exists a non empty finite sequence T of elements of  $\mathbb{R}$ such that
  - (i) T(1) = 0,
- (ii)  $T(\operatorname{len} T) = 1$ ,
- (iii) T is increasing, and
- (iv) there exists an open subset N of Y such that  $y \in N$  and for every natural number i such that  $i \in \text{dom } T$  and  $i + 1 \in \text{dom } T$  there exists a non empty subset  $U_2$  of TopUnitCircle 2 such that  $U_2 \in U_1$  and  $F^{\circ}[N, [T(i), T(i+1)]] \subseteq U_2$ .
- (22) Let Y be a non empty topological space, F be a map from  $[Y, \mathbb{I}]$  into TopUnitCircle 2, and  $F_1$  be a map from  $[Y, Sspace(0_{\mathbb{I}})]$  into  $\mathbb{R}^1$ . Suppose F is continuous and  $F_1$  is continuous and  $F \upharpoonright [$  the carrier of Y,  $\{0\} ] =$ CircleMap  $\cdot F_1$ . Then there exists a map G from  $[Y, \mathbb{I}]$  into  $\mathbb{R}^1$  such that
- (i) G is continuous,
- (ii)  $F = \text{CircleMap} \cdot G$ ,
- (iii)  $G \upharpoonright [$  the carrier of  $Y, \{0\} ] = F_1$ , and
- (iv) for every map H from  $[Y, \mathbb{I}]$  into  $\mathbb{R}^1$  such that H is continuous and  $F = \text{CircleMap} \cdot H$  and  $H \upharpoonright$  the carrier of  $Y, \{0\} \ddagger = F_1$  holds G = H.

- (23) Let  $x_0, y_0$  be points of TopUnitCircle 2,  $x_1$  be a point of  $\mathbb{R}^1$ , and f be a path from  $x_0$  to  $y_0$ . Suppose  $x_1 \in \text{CircleMap}^{-1}(\{x_0\})$ . Then there exists a map  $f_1$  from  $\mathbb{I}$  into  $\mathbb{R}^1$  such that
  - (i)  $f_1(0) = x_1$ ,
  - (ii)  $f = \operatorname{CircleMap} \cdot f_1,$
- (iii)  $f_1$  is continuous, and
- (iv) for every map  $f_2$  from I into  $\mathbb{R}^1$  such that  $f_2$  is continuous and  $f = \text{CircleMap} \cdot f_2$  and  $f_2(0) = x_1$  holds  $f_1 = f_2$ .
- (24) Let  $x_0, y_0$  be points of TopUnitCircle 2, P, Q be paths from  $x_0$  to  $y_0$ , F be a homotopy between P and Q, and  $x_1$  be a point of  $\mathbb{R}^1$ . Suppose P, Q are homotopic and  $x_1 \in \text{CircleMap}^{-1}(\{x_0\})$ . Then there exists a point  $y_1$  of  $\mathbb{R}^1$  and there exist paths  $P_1, Q_1$  from  $x_1$  to  $y_1$  and there exists a homotopy  $F_1$  between  $P_1$  and  $Q_1$  such that  $P_1, Q_1$  are homotopic and  $F = \text{CircleMap} \cdot F_1$  and  $y_1 \in \text{CircleMap}^{-1}(\{y_0\})$  and for every homotopy  $F_2$  between  $P_1$  and  $Q_1$  such that  $F = \text{CircleMap} \cdot F_2$  holds  $F_1 = F_2$ .

The map Ciso from  $\mathbb{Z}^+$  into  $\pi_1$  (TopUnitCircle 2, c[10]) is defined by:

- (Def. 5) For every integer n holds  $(Ciso)(n) = [cLoop n]_{EqRel(TopUnitCircle 2, c[10])}$ . One can prove the following proposition
  - (25) For every integer *i* and for every path *f* from  $R^{10}$  to  $R^{1}i$  holds (Ciso)(*i*) = [CircleMap  $\cdot f$ ]<sub>EqRel(TopUnitCircle 2, c[10])</sub>. Ciso is a homomorphism from  $\mathbb{Z}^+$  to  $\pi_1$ (TopUnitCircle 2, c[10]). Let us mention that Ciso is one-to-one and onto.

We now state two propositions:

- (26) Ciso is isomorphism.
- (27) Let S be a subspace of  $\mathcal{E}_{T}^{2}$  satisfying conditions of simple closed curve and x be a point of S. Then  $\mathbb{Z}^{+}$  and  $\pi_{1}(S, x)$  are isomorphic.

Let S be a subspace of  $\mathcal{E}_{T}^{2}$  satisfying conditions of simple closed curve and let x be a point of S. Note that  $\pi_{1}(S, x)$  is infinite.

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# Brouwer Fixed Point Theorem for Disks on the Plane

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**Summary.** The article formalizes the proof of Brouwer's Fixed Point Theorem for 2-dimensional disks. Assuming, on the contrary, that the theorem is false, we show that a circle is a retract of a disk. Next, using the retraction, we prove that any loop in the circle is homotopic to the constant loop what contradicts with infiniteness of the fundamental group of a circle, see [15].

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The terminology and notation used in this paper are introduced in the following papers: [26], [9], [29], [2], [22], [28], [30], [6], [8], [7], [5], [4], [12], [3], [25], [16], [23], [21], [20], [27], [11], [13], [14], [18], [17], [19], [10], [1], and [24].

In this paper n is a natural number, a, r are real numbers, and x is a point of  $\mathcal{E}^n_{\mathrm{T}}$ .

Let S, T be non empty topological spaces. The functor DiffElems(S,T) yielding a subset of [S, T] is defined by:

(Def. 1) DiffElems $(S, T) = \{ \langle s, t \rangle; s \text{ ranges over points of } S, t \text{ ranges over points of } T: s \neq t \}.$ 

One can prove the following proposition

(1) Let S, T be non empty topological spaces and x be a set. Then  $x \in \text{DiffElems}(S,T)$  if and only if there exists a point s of S and there exists a point t of T such that  $x = \langle s, t \rangle$  and  $s \neq t$ .

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C 2005 University of Białystok ISSN 1426-2630 Let S be a non trivial non empty topological space and let T be a non empty topological space. One can check that DiffElems(S,T) is non empty.

Let S be a non empty topological space and let T be a non trivial non empty topological space. Note that DiffElems(S,T) is non empty.

We now state the proposition

(2)  $Ball(x,0) = \{x\}.$ 

Let *n* be a natural number, let *x* be a point of  $\mathcal{E}^n_{\mathrm{T}}$ , and let *r* be a real number. The functor  $\mathrm{Tdisk}(x, r)$  yields a subspace of  $\mathcal{E}^n_{\mathrm{T}}$  and is defined by:

(Def. 2)  $\operatorname{Tdisk}(x, r) = (\mathcal{E}_{\mathrm{T}}^n) \upharpoonright \overline{\operatorname{Ball}}(x, r).$ 

Let n be a natural number, let x be a point of  $\mathcal{E}_{T}^{n}$ , and let r be a non negative real number. Note that Tdisk(x, r) is non empty.

We now state the proposition

(3) The carrier of  $Tdisk(x, r) = \overline{Ball}(x, r)$ .

Let n be a natural number, let x be a point of  $\mathcal{E}^n_{\mathrm{T}}$ , and let r be a real number. Note that  $\mathrm{Tdisk}(x,r)$  is convex.

We adopt the following convention: n denotes a natural number, r denotes a non negative real number, and s, t, x denote points of  $\mathcal{E}^n_{\mathrm{T}}$ .

One can prove the following two propositions:

- (4) If  $s \neq t$  and s is a point of  $\operatorname{Tdisk}(x, r)$  and s is not a point of  $\operatorname{Tcircle}(x, r)$ , then there exists a point e of  $\operatorname{Tcircle}(x, r)$  such that  $\{e\} = \operatorname{halfline}(s, t) \cap \operatorname{Sphere}(x, r)$ .
- (5) Suppose  $s \neq t$  and  $s \in$  the carrier of Tcircle(x, r) and t is a point of Tdisk(x, r). Then there exists a point e of Tcircle(x, r) such that  $e \neq s$  and  $\{s, e\} = \text{halfline}(s, t) \cap \text{Sphere}(x, r)$ .

Let *n* be a non empty natural number, let *o* be a point of  $\mathcal{E}_{\mathrm{T}}^{n}$ , let *s*, *t* be points of  $\mathcal{E}_{\mathrm{T}}^{n}$ , and let *r* be a non negative real number. Let us assume that *s* is a point of  $\mathrm{Tdisk}(o, r)$ , and *t* is a point of  $\mathrm{Tdisk}(o, r)$  and  $s \neq t$ . The functor  $\mathrm{HC}(s, t, o, r)$  yields a point of  $\mathcal{E}_{\mathrm{T}}^{n}$  and is defined as follows:

(Def. 3)  $\operatorname{HC}(s, t, o, r) \in \operatorname{halfline}(s, t) \cap \operatorname{Sphere}(o, r) \text{ and } \operatorname{HC}(s, t, o, r) \neq s.$ 

In the sequel n is a non empty natural number and s, t, o are points of  $\mathcal{E}_{\mathrm{T}}^{n}$ . We now state three propositions:

- (6) If s is a point of  $\operatorname{Tdisk}(o, r)$  and t is a point of  $\operatorname{Tdisk}(o, r)$  and  $s \neq t$ , then  $\operatorname{HC}(s, t, o, r)$  is a point of  $\operatorname{Tcircle}(o, r)$ .
- (7) Let S, T, O be elements of  $\mathcal{R}^n$ . Suppose S = s and T = t and O = o. Suppose s is a point of  $\mathrm{Tdisk}(o, r)$  and t is a point of  $\mathrm{Tdisk}(o, r)$  and  $s \neq t$  and  $a = \frac{-|(t-s,s-o)| + \sqrt{|(t-s,s-o)|^2 \sum^2 (T-S) \cdot (\sum^2 (S-O) r^2)}}{\sum^2 (T-S)}$ . Then  $\mathrm{HC}(s,t,o,r) = (1-a) \cdot s + a \cdot t$ .
- (8) Let  $r_1$ ,  $r_2$ ,  $s_1$ ,  $s_2$  be real numbers and s, t, o be points of  $\mathcal{E}_{\mathrm{T}}^2$ . Suppose that s is a point of  $\mathrm{Tdisk}(o, r)$  and t is a point of  $\mathrm{Tdisk}(o, r)$  and

 $s \neq t$  and  $r_1 = t_1 - s_1$  and  $r_2 = t_2 - s_2$  and  $s_1 = s_1 - o_1$  and  $s_2 = s_2 - o_2$  and  $a = \frac{-(s_1 \cdot r_1 + s_2 \cdot r_2) + \sqrt{(s_1 \cdot r_1 + s_2 \cdot r_2)^2 - (r_1^2 + r_2^2) \cdot ((s_1^2 + s_2^2) - r^2)}}{r_1^2 + r_2^2}$ . Then  $\operatorname{HC}(s, t, o, r) = [s_1 + a \cdot r_1, s_2 + a \cdot r_2]$ .

Let n be a non empty natural number, let o be a point of  $\mathcal{E}_{\mathrm{T}}^{n}$ , let r be a non negative real number, let x be a point of  $\mathrm{Tdisk}(o,r)$ , and let f be a map from  $\mathrm{Tdisk}(o,r)$  into  $\mathrm{Tdisk}(o,r)$ . Let us assume that x is not a fixpoint of f. The functor  $\mathrm{HC}(x, f)$  yielding a point of  $\mathrm{Tcircle}(o, r)$  is defined as follows:

(Def. 4) There exist points y, z of  $\mathcal{E}^n_T$  such that y = x and z = f(x) and  $\operatorname{HC}(x, f) = \operatorname{HC}(z, y, o, r)$ .

The following two propositions are true:

- (9) Let x be a point of Tdisk(o, r) and f be a map from Tdisk(o, r) into Tdisk(o, r). If x is not a fixpoint of f and x is a point of Tcircle(o, r), then HC(x, f) = x.
- (10) Let r be a positive real number, o be a point of  $\mathcal{E}_{\mathrm{T}}^2$ , and Y be a non empty subspace of Tdisk(o, r). If  $Y = \mathrm{Tcircle}(o, r)$ , then Y is not a retract of Tdisk(o, r).

Let n be a non empty natural number, let r be a non negative real number, let o be a point of  $\mathcal{E}_{T}^{n}$ , and let f be a map from  $\mathrm{Tdisk}(o, r)$  into  $\mathrm{Tdisk}(o, r)$ . The functor BR-map f yielding a map from  $\mathrm{Tdisk}(o, r)$  into  $\mathrm{Tcircle}(o, r)$  is defined as follows:

- (Def. 5) For every point x of Tdisk(o, r) holds (BR-map f)(x) = HC(x, f). The following propositions are true:
  - (11) Let o be a point of  $\mathcal{E}_{T}^{n}$ , x be a point of Tdisk(o, r), and f be a map from Tdisk(o, r) into Tdisk(o, r). If x is not a fixpoint of f and x is a point of Tcircle(o, r), then (BR-map f)(x) = x.
  - (12) For every continuous map f from  $\operatorname{Tdisk}(o, r)$  into  $\operatorname{Tdisk}(o, r)$  such that f has no fixpoint holds BR-map  $f \upharpoonright \operatorname{Sphere}(o, r) = \operatorname{id}_{\operatorname{Tcircle}(o, r)}$ .
  - (13) Let r be a positive real number, o be a point of  $\mathcal{E}_{T}^{2}$ , and f be a continuous map from  $\mathrm{Tdisk}(o, r)$  into  $\mathrm{Tdisk}(o, r)$ . If f has no fixpoint, then BR-map f is continuous.
  - (14) For every non negative real number r and for every point o of  $\mathcal{E}_{\mathrm{T}}^2$  holds every continuous map from  $\mathrm{Tdisk}(o, r)$  into  $\mathrm{Tdisk}(o, r)$  has a fixpoint.
  - (15) Let r be a non negative real number, o be a point of  $\mathcal{E}_{\mathrm{T}}^2$ , and f be a continuous map from  $\mathrm{Tdisk}(o, r)$  into  $\mathrm{Tdisk}(o, r)$ . Then there exists a point x of  $\mathrm{Tdisk}(o, r)$  such that f(x) = x.

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## Stirling Numbers of the Second Kind

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**Summary.** In this paper we define Stirling numbers of the second kind by cardinality of certain functional classes so that

 $S(n,k) = \{ f \text{ where } f \text{ is function of } n, k : f \text{ is onto increasing} \}$ 

After that we show basic properties of this number in order to prove recursive dependence of Stirling number of the second kind. Consecutive theorems are introduced to prove formula

$$S(n,k) = \frac{1}{k!} \sum_{i=0}^{k-1} (-1)^i \binom{k}{i} (k-i)^n$$

where  $k \leq n$ .

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The papers [18], [9], [21], [14], [23], [6], [24], [2], [3], [8], [10], [1], [22], [7], [11], [20], [16], [19], [4], [5], [13], [12], [17], and [15] provide the terminology and notation for this paper.

For simplicity, we adopt the following convention: k, l, m, n, i, j denote natural numbers, K, N denote non empty subsets of  $\mathbb{N}, K_1, N_1, M_1$  denote subsets of  $\mathbb{N}$ , and X, Y denote sets.

Let us consider k. Then  $\{k\}$  is a subset of N. Let us consider l. Then  $\{k, l\}$  is a subset of N. Let us consider m. Then  $\{k, l, m\}$  is a non empty subset of N.

The following propositions are true:

- (1)  $\min N = \min^* N.$
- (2)  $\min(\min K, \min N) = \min(K \cup N).$
- (3)  $\min(\min^* K_1, \min^* N_1) \le \min^* (K_1 \cup N_1).$

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- (4) If  $\min^* N_1 \notin N_1 \cap K_1$ , then  $\min^* N_1 = \min^* (N_1 \setminus K_1)$ .
- (5)  $\min^*{n} = n \text{ and } \min{n} = n.$
- (6)  $\min^{*}\{n,k\} = \min(n,k) \text{ and } \min\{n,k\} = \min(n,k).$
- (7)  $\min^{*}\{n, k, l\} = \min(n, \min(k, l)).$
- (8) n is a subset of  $\mathbb{N}$ .

Let us consider n. One can verify that every element of n is natural. We now state several propositions:

- (9) If  $N \subseteq n$ , then n-1 is a natural number.
- (10) If  $k \in n$ , then  $k \le n-1$  and n-1 is a natural number.
- (11)  $\min^* n = 0.$
- (12) If  $N \subseteq n$ , then  $\min^* N \leq n 1$ .
- (13) If  $N \subseteq n$  and  $N \neq \{n-1\}$ , then  $\min^* N < n-1$ .
- (14) If  $N_1 \subseteq n$  and n > 0, then  $\min^* N_1 \leq n 1$ .

In the sequel f, g are functions from n into k.

Let us consider n, X, let f be a function from n into X, and let x be a set. Then  $f^{-1}(x)$  is a subset of  $\mathbb{N}$ .

Let us consider X, k, let f be a function from X into k, and let x be a set. Then f(x) is an element of k.

Let us consider X,  $N_1$ , let f be a function from X into  $N_1$ , and let x be a set. One can verify that f(x) is natural.

Let us consider n, k and let f be a function from n into k. We say that f is increasing if and only if:

(Def. 1) n = 0 iff k = 0 and for all l, m such that  $l \in \operatorname{rng} f$  and  $m \in \operatorname{rng} f$  and l < m holds  $\min^*(f^{-1}(\{l\})) < \min^*(f^{-1}(\{m\}))$ .

We now state several propositions:

- (15) If n = 0 and k = 0, then f is onto and increasing.
- (16) If n > 0, then  $\min^*(f^{-1}(\{m\})) \le n 1$ .
- (17) If f is onto, then  $n \ge k$ .
- (18) If f is onto and increasing, then for every m such that m < k holds  $m \leq \min^*(f^{-1}(\{m\})).$
- (19) If f is onto and increasing, then for every m such that m < k holds  $\min^*(f^{-1}(\{m\})) \le (n-k) + m$ .
- (20) If f is onto and increasing and n = k, then  $f = id_n$ .
- (21) If  $f = id_n$  and n > 0, then f is increasing.
- (22) If n = 0 iff k = 0, then there exists a function from n into k which is increasing.
- (23) If n = 0 iff k = 0 and  $n \ge k$ , then there exists a function from n into k which is onto and increasing.

The scheme *Sch1* deals with natural numbers  $\mathcal{A}$ ,  $\mathcal{B}$  and a unary predicate  $\mathcal{P}$ , and states that:

 $\{f; f \text{ ranges over functions from } \mathcal{A} \text{ into } \mathcal{B} : \mathcal{P}[f]\}$  is finite

for all values of the parameters.

In the sequel f is a function from n into k.

One can prove the following propositions:

- (24) For all n, k holds  $\{f : f \text{ is onto and increasing}\}$  is finite.
- (25) For all n, k holds  $\overline{\{f: f \text{ is onto and increasing}\}}$  is a natural number.

Let us consider n, k. The functor n block k yields a natural number and is defined by:

(Def. 2)  $n \operatorname{block} k = \overline{\{f : f \text{ is onto and increasing}\}}.$ Next we state several propositions:

- (26)  $n \operatorname{block} n = 1.$
- (27) If  $k \neq 0$ , then 0 block k = 0.
- (28) 0 block k = 1 iff k = 0.
- (29) If n < k, then  $n \operatorname{block} k = 0$ .
- (30)  $n \operatorname{block} 0 = 1 \operatorname{iff} n = 0.$
- (31) If  $n \neq 0$ , then  $n \operatorname{block} 0 = 0$ .
- (32) If  $n \neq 0$ , then n block 1 = 1.
- (33)  $1 \le k$  and  $k \le n$  or k = n iff n block k > 0.

In the sequel x, y denote sets.

Now we present three schemes. The scheme  $Sch^2$  deals with sets  $\mathcal{A}, \mathcal{B}, \mathcal{C}, \mathcal{D}$ , a function  $\mathcal{E}$  from  $\mathcal{A}$  into  $\mathcal{B}$ , and a unary functor  $\mathcal{F}$  yielding a set, and states that:

There exists a function h from  $\mathcal{C}$  into  $\mathcal{D}$  such that  $h \upharpoonright \mathcal{A} = \mathcal{E}$  and

for every x such that  $x \in \mathcal{C} \setminus \mathcal{A}$  holds  $h(x) = \mathcal{F}(x)$ 

provided the parameters satisfy the following conditions:

- For every x such that  $x \in \mathcal{C} \setminus \mathcal{A}$  holds  $\mathcal{F}(x) \in \mathcal{D}$ ,
- $\mathcal{A} \subseteq \mathcal{C}$  and  $\mathcal{B} \subseteq \mathcal{D}$ , and
- If  $\mathcal{B}$  is empty, then  $\mathcal{A}$  is empty.

The scheme *Sch3* deals with sets  $\mathcal{A}$ ,  $\mathcal{B}$ ,  $\mathcal{C}$ ,  $\mathcal{D}$ , a unary functor  $\mathcal{F}$  yielding a set, and a ternary predicate  $\mathcal{P}$ , and states that:

f; f ranges over functions from  $\mathcal{A}$  into  $\mathcal{B}: \mathcal{P}[f, \mathcal{A}, \mathcal{B}] =$ 

$$\{f: f \text{ ranges over functions from } \mathcal{C} \text{ into } \mathcal{D} : \mathcal{P}[f, \mathcal{C}, \mathcal{D}] \}$$

 $\overline{\wedge \operatorname{rng}(f \upharpoonright \mathcal{A}) \subseteq \mathcal{B} \land \bigwedge_x (x \in \mathcal{C} \setminus \mathcal{A} \Rightarrow f(x) = \mathcal{F}(x))\}}$ provided the following requirements are met:

- For every x such that  $x \in \mathcal{C} \setminus \mathcal{A}$  holds  $\mathcal{F}(x) \in \mathcal{D}$ ,
- $\mathcal{A} \subseteq \mathcal{C}$  and  $\mathcal{B} \subseteq \mathcal{D}$ ,
- If  $\mathcal{B}$  is empty, then  $\mathcal{A}$  is empty, and

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• Let f be a function from  $\mathcal{C}$  into  $\mathcal{D}$ . Suppose that for every x such that  $x \in \mathcal{C} \setminus \mathcal{A}$  holds  $\mathcal{F}(x) = f(x)$ . Then  $\mathcal{P}[f, \mathcal{C}, \mathcal{D}]$  if and only if  $\mathcal{P}[f \upharpoonright \mathcal{A}, \mathcal{A}, \mathcal{B}]$ .

The scheme *Sch4* deals with sets  $\mathcal{A}, \mathcal{B}, \mathcal{C}, \mathcal{D}$  and a ternary predicate  $\mathcal{P}$ , and states that:

$\{f; f$	f ranges	over	functions	from	${\cal A}~{ m into}~{\cal A}$	$\mathcal{B}:\mathcal{P}[f,$	$[\mathcal{A},\mathcal{B}]\}=$
$\{f; f$	<sup>r</sup> ranges	over	functions	from	$\mathcal{A} \cup \{\mathcal{C}\}$	$\}$ into ${\cal B}$	$\cup \{\mathcal{D}\}$ :
$\overline{\mathcal{P}[f,}$	$\overline{\mathcal{A}\cup\{\mathcal{C}\}}$	$\mathcal{B} \cup$	$ \{\mathcal{D}\}  \wedge 1$	$\operatorname{rng}(f)$	$(\mathcal{A}) \subseteq \mathcal{B}$	$f \wedge f(\mathcal{C}$	$) = \mathcal{D} \}$

provided the parameters meet the following conditions:

- If  $\mathcal{B}$  is empty, then  $\mathcal{A}$  is empty,
- $\mathcal{C} \notin \mathcal{A}$ , and
- For every function f from  $\mathcal{A} \cup \{\mathcal{C}\}$  into  $\mathcal{B} \cup \{\mathcal{D}\}$  such that  $f(\mathcal{C}) = \mathcal{D}$  holds  $\mathcal{P}[f, \mathcal{A} \cup \{\mathcal{C}\}, \mathcal{B} \cup \{\mathcal{D}\}]$  iff  $\mathcal{P}[f \mid \mathcal{A}, \mathcal{A}, \mathcal{B}]$ .

We now state several propositions:

- (34) For every function f from n + 1 into k + 1 such that f is onto and increasing and  $f^{-1}(\{f(n)\}) = \{n\}$  holds f(n) = k.
- (35) For every function f from n+1 into k such that  $k \neq 0$  and  $f^{-1}(\{f(n)\}) \neq \{n\}$  there exists m such that  $m \in f^{-1}(\{f(n)\})$  and  $m \neq n$ .
- (36) Let f be a function from n into k and g be a function from n + m into k+l. Suppose g is increasing and  $f = g \upharpoonright n$ . Let given i, j. If  $i \in \operatorname{rng} f$  and  $j \in \operatorname{rng} f$  and i < j, then  $\min^*(f^{-1}(\{i\})) < \min^*(f^{-1}(\{j\}))$ .
- (37) Let f be a function from n+1 into k+1. Suppose f is onto and increasing and  $f^{-1}(\{f(n)\}) = \{n\}$ . Then  $\operatorname{rng}(f \upharpoonright n) \subseteq k$  and for every function g from n into k such that  $g = f \upharpoonright n$  holds g is onto and increasing.
- (38) Let f be a function from n + 1 into k and g be a function from n into k. Suppose f is onto and increasing and  $f^{-1}(\{f(n)\}) \neq \{n\}$  and  $f \upharpoonright n = g$ . Then g is onto and increasing.
- (39) Let f be a function from n into k and g be a function from n + 1 into k + m. Suppose f is onto and increasing and  $f = g \upharpoonright n$ . Let given i, j. If  $i \in \operatorname{rng} g$  and  $j \in \operatorname{rng} g$  and i < j, then  $\min^*(g^{-1}(\{i\})) < \min^*(g^{-1}(\{j\}))$ .
- (40) Let f be a function from n into k and g be a function from n + 1 into k + 1. Suppose f is onto and increasing and  $f = g \upharpoonright n$  and g(n) = k. Then g is onto and increasing and  $g^{-1}(\{g(n)\}) = \{n\}$ .
- (41) Let f be a function from n into k and g be a function from n + 1 into k. Suppose f is onto and increasing and  $f = g \upharpoonright n$  and g(n) < k. Then g is onto and increasing and  $g^{-1}(\{g(n)\}) \neq \{n\}$ .

In the sequel  $f_1$  denotes a function from n + 1 into k + 1 and f denotes a function from n into k.

We now state the proposition

(42)  $\overline{\{f_1: f_1 \text{ is onto and increasing } \land f_1^{-1}(\{f_1(n)\}) = \{n\}\}} =$ 

 $\{f: f \text{ is onto and increasing}\}.$ 

In the sequel f' is a function from n+1 into k.

The following proposition is true

(43) For every l such that l < k holds  $\frac{\overline{\{f': f' \text{ is onto and increasing} \land f'^{-1}(\{f'(n)\}) \neq \{n\} \land f'(n) = l\}}{\{f: f \text{ is onto and increasing}\}} =$ 

For simplicity, we adopt the following convention: D denotes a non empty set, F, G denote finite 0-sequences of D,  $F_1$  denotes a finite 0-sequence of  $\mathbb{N}$ , bdenotes a binary operation on D, and d,  $d_1$ ,  $d_2$  denote elements of D.

Let us consider D, F, b. Let us assume that b has a unity or len  $F \ge 1$ . The functor  $b \odot F$  yielding an element of D is defined as follows:

(Def. 3)(i)  $b \odot F = \mathbf{1}_b$  if b has a unity and len F = 0,

(ii) there exists a function f from  $\mathbb{N}$  into D such that f(0) = F(0) and for every n such that n + 1 < len F holds f(n + 1) = b(f(n), F(n + 1)) and  $b \odot F = f(\text{len } F - 1)$ , otherwise.

One can prove the following three propositions:

- $(44) \quad b \odot \langle d \rangle = d.$
- (45) If b has a unity or len F > 0, then  $b \odot F \cap \langle d \rangle = b(b \odot F, d)$ .
- (46) If  $F \neq \langle \rangle_D$ , then there exist G, d such that  $F = G \cap \langle d \rangle$ .

The scheme *Sch5* deals with a non empty set  $\mathcal{A}$  and a unary predicate  $\mathcal{P}$ , and states that:

For every finite 0-sequence F of  $\mathcal{A}$  holds  $\mathcal{P}[F]$ 

provided the parameters satisfy the following conditions:

- $\mathcal{P}[\langle \rangle_{\mathcal{A}}]$ , and
- For every finite 0-sequence F of  $\mathcal{A}$  and for every element d of  $\mathcal{A}$  such that  $\mathcal{P}[F]$  holds  $\mathcal{P}[F \cap \langle d \rangle]$ .

Next we state the proposition

(47) If b is associative and if b has a unity or len  $F \ge 1$  and len  $G \ge 1$ , then  $b \odot F \cap G = b(b \odot F, b \odot G)$ .

Let us consider D and let us consider  $d, d_1$ . Then  $\langle d, d_1 \rangle$  is a finite 0-sequence of D. Let us consider  $d_2$ . Then  $\langle d, d_1, d_2 \rangle$  is a finite 0-sequence of D.

The following propositions are true:

- $(48) \quad b \odot \langle d_1, d_2 \rangle = b(d_1, d_2).$
- (49)  $b \odot \langle d, d_1, d_2 \rangle = b(b(d, d_1), d_2).$

Let us consider  $F_1$ . The functor  $\sum F_1$  yields a natural number and is defined by:

(Def. 4)  $\sum F_1 = +_{\mathbb{N}} \odot F_1$ .

Let us consider  $F_1$ , x. Then  $F_1(x)$  is a natural number. One can prove the following propositions:

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- (50) If for every n such that  $n \in \text{dom } F_1$  holds  $F_1(n) \leq k$ , then  $\sum F_1 \leq K$ len  $F_1 \cdot k$ .
- (51) If for every n such that  $n \in \operatorname{dom} F_1$  holds  $F_1(n) \geq k$ , then  $\sum F_1 \geq k$ len  $F_1 \cdot k$ .
- (52) If len  $F_1 > 0$  and there exists x such that  $x \in \text{dom } F_1$  and  $F_1(x) = k$ , then  $\sum F_1 \ge k$ .
- (53)  $\sum F_1 = 0$  iff len  $F_1 = 0$  or for every n such that  $n \in \operatorname{dom} F_1$  holds  $F_1(n) = 0.$
- (54) For every function f and for every n holds  $\bigcup \operatorname{rng}(f \upharpoonright n) \cup f(n) =$  $\bigcup \operatorname{rng}(f{\upharpoonright}(n+1)).$

Now we present three schemes. The scheme Sch6 deals with a non empty set  $\mathcal{A}$ , a natural number  $\mathcal{B}$ , and a binary predicate  $\mathcal{P}$ , and states that:

There exists a finite 0-sequence p of  $\mathcal{A}$  such that dom  $p = \mathcal{B}$  and for every k such that  $k \in \mathcal{B}$  holds  $\mathcal{P}[k, p(k)]$ 

provided the parameters have the following property:

• For every k such that  $k \in \mathcal{B}$  there exists an element x of A such that  $\mathcal{P}[k, x]$ .

The scheme Sch7 deals with a non empty set  $\mathcal{A}$  and a finite 0-sequence  $\mathcal{B}$  of  $\mathcal{A}$ , and states that:

There exists a finite 0-sequence  $C_1$  of  $\mathbb{N}$  such that dom  $C_1 = \operatorname{dom} \mathcal{B}$ and for every i such that  $i \in \operatorname{dom} C_1$  holds  $C_1(i) = \overline{\mathcal{B}(i)}$  and  $\bigcup \operatorname{rng} \mathcal{B} = \sum C_1$ 

provided the following requirements are met:

- For every i such that  $i \in \text{dom } \mathcal{B}$  holds  $\mathcal{B}(i)$  is finite, and
- For all i, j such that  $i \in \operatorname{dom} \mathcal{B}$  and  $j \in \operatorname{dom} \mathcal{B}$  and  $i \neq j$  holds  $\mathcal{B}(i)$  misses  $\mathcal{B}(j)$ .

The scheme Sch8 deals with finite sets  $\mathcal{A}$ ,  $\mathcal{B}$ , a set  $\mathcal{C}$ , a function  $\mathcal{D}$  from card  $\mathcal{B}$  into  $\mathcal{B}$ , and a unary predicate  $\mathcal{P}$ , and states that:

- There exists a finite 0-sequence F of  $\mathbb{N}$  such that
  - $\operatorname{dom} F = \operatorname{card} \mathcal{B},$ (i)
  - $\overline{\{g; g \text{ ranges over functions from } \mathcal{A} \text{ into } \mathcal{B} : \mathcal{P}[g]\}} = \sum F,$ (ii) and

(iii) for every *i* such that  $i \in \operatorname{dom} F$  holds  $F(i) = \{g; g \text{ ranges over functions from } \mathcal{A} \text{ into } \mathcal{B} : \mathcal{P}[g] \land g(\mathcal{C}) = \mathcal{D}(i)\}$ 

provided the parameters have the following properties:

- $\mathcal{D}$  is onto and one-to-one,
- $\mathcal{B}$  is non empty, and
- $\mathcal{C} \in \mathcal{A}$ .

One can prove the following propositions:

(55) 
$$k \cdot (n \operatorname{block} k) = \{ f' : f' \text{ is onto and increasing } \land f'^{-1}(\{ f'(n) \}) \neq \{ n \} \}.$$
- (56) (n+1) block $(k+1) = (k+1) \cdot (n \operatorname{block}(k+1)) + (n \operatorname{block} k).$
- (57) If  $n \ge 1$ , then  $n \operatorname{block} 2 = \frac{1}{2} \cdot (2^n 2)$ .
- (58) If  $n \ge 2$ , then n block  $3 = \frac{1}{6} \cdot ((3^n 3 \cdot 2^n) + 3)$ .
- (59) If  $n \ge 3$ , then  $n \operatorname{block} 4 = \frac{1}{24} \cdot (((4^n 4 \cdot 3^n) + 6 \cdot 2^n) 4).$
- (60) 3! = 6 and 4! = 24.
- (61)  $\binom{n}{1} = n$  and  $\binom{n}{2} = \frac{n \cdot (n-1)}{2}$  and  $\binom{n}{3} = \frac{n \cdot (n-1) \cdot (n-2)}{6}$  and  $\binom{n}{4} = \frac{n \cdot (n-1) \cdot (n-2)}{24}$ .
- (62) (n+1) block  $n = \binom{n+1}{2}$ .
- (63) (n+2) block  $n = 3 \cdot \binom{n+2}{4} + \binom{n+2}{3}$ .
- (64) For every function F and for every y holds  $\operatorname{rng}(F \upharpoonright (\operatorname{dom} F \setminus F^{-1}(\{y\}))) =$ rng  $F \setminus \{y\}$  and for every x such that  $x \neq y$  holds  $(F \mid (\operatorname{dom} F \setminus$  $F^{-1}(\{y\})))^{-1}(\{x\}) = F^{-1}(\{x\}).$

(65) If 
$$\overline{X} = k + 1$$
 and  $x \in X$ , then  $\overline{X \setminus \{x\}} = k$ .

The scheme *Sch9* concerns a unary predicate  $\mathcal{P}$  and a binary predicate  $\mathcal{Q}$ , and states that:

For every function F such that rng F is finite holds  $\mathcal{P}[F]$ 

provided the following conditions are met:

- $\mathcal{P}[\emptyset]$ , and
- For every function F such that for every x such that  $x \in \operatorname{rng} F$ and  $\mathcal{Q}[x, F]$  holds  $\mathcal{P}[F \upharpoonright (\operatorname{dom} F \setminus F^{-1}(\{x\}))]$  holds  $\mathcal{P}[F]$ .

We now state several propositions:

- (66) For every subset N of N such that N is finite there exists k such that for every n such that  $n \in N$  holds  $n \leq k$ .
- (67) Let given X, Y, x, y. Suppose if Y is empty, then X is empty and  $x \notin X$ . Let F be a function from X into Y. Then there exists a function G from  $X \cup \{x\}$  into  $Y \cup \{y\}$  such that  $G \upharpoonright X = F$  and G(x) = y.
- (68) Let given X, Y, x, y such that if Y is empty, then X is empty. Let F be a function from X into Y and G be a function from  $X \cup \{x\}$  into  $Y \cup \{y\}$ such that  $G \upharpoonright X = F$  and G(x) = y. Then
  - if F is onto, then G is onto, and (i)
  - if  $y \notin Y$  and F is one-to-one, then G is one-to-one. (ii)
- (69) Let N be a finite subset of N. Then there exists a function  $O_1$  from N into card N such that  $O_1$  is bijective and for all n, k such that  $n \in \text{dom } O_1$ and  $k \in \text{dom } O_1$  and n < k holds  $O_1(n) < O_1(k)$ .
- (70) Let X, Y be finite sets and F be a function from X into Y. If card X =card Y, then F is onto iff F is one-to-one.
- (71) Let F, G be functions and given y. Suppose  $y \in \operatorname{rng}(G \cdot F)$  and G is one-to-one. Then there exists x such that  $x \in \text{dom } G$  and  $x \in \text{rng } F$  and  $G^{-1}(\{y\}) = \{x\}$  and  $F^{-1}(\{x\}) = (G \cdot F)^{-1}(\{y\}).$

Let us consider  $N_1$ ,  $K_1$  and let f be a function from  $N_1$  into  $K_1$ . We say that f is increasing if and only if:

(Def. 5) For all l, m such that  $l \in \operatorname{rng} f$  and  $m \in \operatorname{rng} f$  and l < m holds  $\min^*(f^{-1}(\{l\})) < \min^*(f^{-1}(\{m\})).$ 

The following four propositions are true:

- (72) For every function F from  $N_1$  into  $K_1$  such that F is increasing holds  $\min^* \operatorname{rng} F = F(\min^* \operatorname{dom} F).$
- (73) Let F be a function from  $N_1$  into  $K_1$ . Suppose rng F is finite. Then there exists a function I from  $N_1$  into  $K_1$  and there exists a permutation P of rng F such that  $F = P \cdot I$  and rng F = rng I and I is increasing.
- (74) Let F be a function from  $N_1$  into  $K_1$ . Suppose rng F is finite. Let  $I_1$ ,  $I_2$  be functions from  $N_1$  into  $M_1$  and  $P_1$ ,  $P_2$  be functions. Suppose that  $P_1$  is one-to-one and  $P_2$  is one-to-one and rng  $I_1 = \operatorname{rng} I_2$  and rng  $I_1 = \operatorname{dom} P_1$  and dom  $P_1 = \operatorname{dom} P_2$  and  $F = P_1 \cdot I_1$  and  $F = P_2 \cdot I_2$  and  $I_1$  is increasing and  $I_2$  is increasing. Then  $P_1 = P_2$  and  $I_1 = I_2$ .
- (75) Let F be a function from  $N_1$  into  $K_1$ . Suppose rng F is finite. Let  $I_1$ ,  $I_2$  be functions from  $N_1$  into  $K_1$  and  $P_1$ ,  $P_2$  be permutations of rng F. Suppose  $F = P_1 \cdot I_1$  and  $F = P_2 \cdot I_2$  and rng  $F = \operatorname{rng} I_1$  and rng  $F = \operatorname{rng} I_2$  and  $I_1$  is increasing and  $I_2$  is increasing. Then  $P_1 = P_2$  and  $I_1 = I_2$ .

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# Limit of Sequence of Subsets

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**Summary.** A concept of "limit of sequence of subsets" is defined here. This article contains the following items: 1. definition of the superior sequence and the inferior sequence of sets, 2. definition of the superior limit and the inferior limit of sets, and additional properties for the sigma-field of sets, 3. definition of the limit value of a convergent sequence of sets, and additional properties for the sigma-field of sets.

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The notation and terminology used here are introduced in the following papers: [9], [1], [13], [2], [10], [6], [11], [4], [12], [14], [8], [7], [3], and [5].

For simplicity, we adopt the following rules:  $n, m, k, k_1, k_2$  denote natural numbers, x, X, Y, Z denote sets, A denotes a subset of  $X, B, A_1, A_2, A_3$  denote sequences of subsets of  $X, S_1$  denotes a  $\sigma$ -field of subsets of X, and  $S, S_2, S_3, S_4$  denote sequences of subsets of  $S_1$ .

Next we state a number of propositions:

- (1) For every function f from  $\mathbb{N}$  into Y and for every n holds  $\{f(k) : n \leq k\} \neq \emptyset$ .
- (2) For every function f from N into Y holds  $f(n+m) \in \{f(k) : n \le k\}$ .
- (3) For every function f from  $\mathbb{N}$  into Y holds  $\{f(k_1) : n \le k_1\} = \{f(k_2) : n+1 \le k_2\} \cup \{f(n)\}.$
- (4) Let f be a function from  $\mathbb{N}$  into Y. Then for every  $k_1$  holds  $x \in f(n+k_1)$  if and only if for every Z such that  $Z \in \{f(k_2) : n \leq k_2\}$  holds  $x \in Z$ .
- (5) For every non empty set Y and for every function f from N into Y holds  $x \in \operatorname{rng} f$  iff there exists n such that x = f(n).
- (6) For every non empty set Y and for every function f from N into Y holds rng  $f = \{f(k)\}.$

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- (7) For every non empty set Y and for every function f from N into Y holds  $\operatorname{rng}(f \uparrow k) = \{f(n) : k \leq n\}.$
- (8)  $x \in \bigcap \operatorname{rng} B$  iff for every *n* holds  $x \in B(n)$ .
- (9) Intersection  $B = \bigcap \operatorname{rng} B$ .
- (10) Intersection  $B \subseteq \bigcup B$ .
- (11) If for every n holds B(n) = A, then  $\bigcup B = A$ .
- (12) If for every n holds B(n) = A, then Intersection B = A.
- (13) If B is constant, then  $\bigcup B =$  Intersection B.
- (14) If B is constant and the value of B = A, then for every n holds  $\bigcup \{B(k) : n \le k\} = A$ .
- (15) If B is constant and the value of B = A, then for every n holds  $\bigcap \{B(k) : n \le k\} = A$ .
- (16) Let given X, B and f be a function. Suppose dom  $f = \mathbb{N}$  and for every n holds  $f(n) = \bigcap \{B(k) : n \leq k\}$ . Then f is a sequence of subsets of X.
- (17) Let X be a set, B be a sequence of subsets of X, and f be a function. Suppose dom  $f = \mathbb{N}$  and for every n holds  $f(n) = \bigcup \{B(k) : n \leq k\}$ . Then f is a function from  $\mathbb{N}$  into  $2^X$ .

Let us consider X, B. We say that B is monotone if and only if:

(Def. 1) B is non-decreasing or non-increasing.

Let B be a function. The inferior setsequence B yields a function and is defined by the conditions (Def. 2).

- (Def. 2)(i) dom (the inferior setsequence B) =  $\mathbb{N}$ , and
  - (ii) for every *n* holds (the inferior setsequence B) $(n) = \bigcap \{B(k) : n \le k\}$ .

Let X be a set and let B be a sequence of subsets of X. Then the inferior sets equence B is a sequence of subsets of X.

Let B be a function. The superior setsequence B yields a function and is defined by the conditions (Def. 3).

(Def. 3)(i) dom (the superior setsequence B) =  $\mathbb{N}$ , and

(ii) for every *n* holds (the superior setsequence B) $(n) = \bigcup \{B(k) : n \le k\}$ .

Let X be a set and let B be a sequence of subsets of X. Then the superior sets equence B is a sequence of subsets of X.

Next we state several propositions:

- (18) (The inferior sets equence B)(0) = Intersection B.
- (19) (The superior setsequence B)(0) =  $\bigcup B$ .
- (20)  $x \in (\text{the inferior sets equence } B)(n)$  iff for every k holds  $x \in B(n+k)$ .
- (21)  $x \in (\text{the superior sets} equence B)(n)$  iff there exists k such that  $x \in B(n+k)$ .
- (22) (The inferior set sequence B)(n) = (the inferior set sequence B) $(n+1) \cap B(n)$ .

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- (23) (The superior sets equence B)(n) = (the superior sets equence B) $(n+1) \cup B(n)$ .
- (24) The inferior sets equence B is non-decreasing.
- (25) The superior sets equence B is non-increasing.
- (26) The inferior sets equence B is monotone and the superior sets equence B is monotone.

Let X be a set and let A be a sequence of subsets of X. Observe that the inferior sets equence A is non-decreasing.

Let X be a set and let A be a sequence of subsets of X. Observe that the superior sets equence A is non-increasing.

The following propositions are true:

- (27) Intersection  $B \subseteq (\text{the inferior setsequence } B)(n)$ .
- (28) (The superior sets equence B) $(n) \subseteq \bigcup B$ .
- (29) For all B, n holds  $\{B(k) : n \le k\}$  is a family of subsets of X.
- (30)  $\bigcup B = (\text{Intersection Complement } B)^c$ .
- (31) (The inferior setsequence B)(n) = (the superior setsequence Complement B) $(n)^{c}$ .
- (32) (The superior setsequence B)(n) = (the inferior setsequence Complement B) $(n)^{c}$ .
- (33) Complement (the inferior sets equence B) = the superior sets equence Complement B.
- (34) Complement (the superior sets equence B) = the inferior sets equence Complement B.
- (35) Suppose that for every n holds  $A_3(n) = A_1(n) \cup A_2(n)$ . Let given n. Then (the inferior setsequence  $B(n) \cup$  (the inferior setsequence  $A_2(n) \subseteq$  (the inferior setsequence  $A_3(n)$ .
- (36) Suppose that for every n holds  $A_3(n) = A_1(n) \cap A_2(n)$ . Let given n. Then (the inferior setsequence  $A_3(n) =$  (the inferior setsequence  $A_1(n) \cap$  (the inferior setsequence  $A_2(n)$ .
- (37) Suppose that for every n holds  $A_3(n) = A_1(n) \cup A_2(n)$ . Let given n. Then (the superior setsequence  $A_3(n) =$  (the superior setsequence  $A_1(n) \cup$  (the superior setsequence  $A_2(n)$ .
- (38) Suppose that for every n holds  $A_3(n) = A_1(n) \cap A_2(n)$ . Let given n. Then (the superior setsequence  $A_3(n) \subseteq$  (the superior setsequence  $A_1(n) \cap$  (the superior setsequence  $A_2(n)$ .
- (39) If B is constant and the value of B = A, then for every n holds (the inferior setsequence B)(n) = A.
- (40) If B is constant and the value of B = A, then for every n holds (the superior setsequence B)(n) = A.

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- (41) If B is non-decreasing, then  $B(n) \subseteq$  (the superior setsequence B)(n+1).
- (42) If B is non-decreasing, then (the superior sets equence B)(n) = (the superior sets equence B)(n + 1).
- (43) If B is non-decreasing, then (the superior setsequence B) $(n) = \bigcup B$ .
- (44) If B is non-decreasing, then Intersection (the superior sets equence B) =  $\bigcup B$ .
- (45) If B is non-decreasing, then  $B(n) \subseteq$  (the inferior setsequence B)(n+1).
- (46) If B is non-decreasing, then (the inferior setsequence B)(n) = B(n).
- (47) If B is non-decreasing, then the inferior sets equence B = B.
- (48) If B is non-increasing, then (the superior sets equence B) $(n+1) \subseteq B(n)$ .
- (49) If B is non-increasing, then (the superior sets equence B)(n) = B(n).
- (50) If B is non-increasing, then the superior sets equence B = B.
- (51) If B is non-increasing, then (the inferior setsequence B) $(n+1) \subseteq B(n)$ .
- (52) If B is non-increasing, then (the inferior setsequence B)(n) = (the inferior setsequence B)(n + 1).
- (53) If B is non-increasing, then (the inferior sets equence B)(n) = Intersection B.
- (54) If B is non-increasing, then  $\bigcup$  (the inferior setsequence B) = Intersection B.

Let X be a set and let B be a sequence of subsets of X. Then  $\liminf B$  can be characterized by the condition:

(Def. 4)  $\liminf B = \bigcup$  (the inferior setsequence B).

Let X be a set and let B be a sequence of subsets of X. Then  $\limsup B$  can be characterized by the condition:

(Def. 5)  $\limsup B =$ Intersection (the superior setsequence B).

Let X be a set and let B be a sequence of subsets of X. We introduce  $\lim B$  as a synonym of  $\limsup B$ .

Next we state a number of propositions:

- (55) Intersection  $B \subseteq \liminf B$ .
- (56)  $\liminf B = \lim (\text{the inferior setsequence } B).$
- (57)  $\limsup B = \lim (\text{the superior sets equence } B).$
- (58)  $\limsup B = (\liminf \operatorname{Complement} B)^{c}$ .
- (59) If B is constant and the value of B = A, then B is convergent and  $\lim B = A$  and  $\lim \inf B = A$  and  $\limsup B = A$ .
- (60) If B is non-decreasing, then  $\limsup B = \bigcup B$ .
- (61) If B is non-decreasing, then  $\liminf B = \bigcup B$ .
- (62) If B is non-increasing, then  $\limsup B = \operatorname{Intersection} B$ .
- (63) If B is non-increasing, then  $\liminf B = \operatorname{Intersection} B$ .

- (64) If B is non-decreasing, then B is convergent and  $\lim B = \bigcup B$ .
- (65) If B is non-increasing, then B is convergent and  $\lim B = \text{Intersection } B$ .
- (66) If B is monotone, then B is convergent.

Let X be a set, let  $S_1$  be a  $\sigma$ -field of subsets of X, and let S be a sequence of subsets of  $S_1$ . Let us observe that S is constant if and only if:

(Def. 6) There exists an element A of  $S_1$  such that for every n holds S(n) = A.

Let X be a set, let  $S_1$  be a  $\sigma$ -field of subsets of X, and let S be a sequence of subsets of  $S_1$ . Then the inferior sets equence S is a sequence of subsets of  $S_1$ .

Let X be a set, let  $S_1$  be a  $\sigma$ -field of subsets of X, and let S be a sequence of subsets of  $S_1$ . Then the superior sets equence S is a sequence of subsets of  $S_1$ .

The following propositions are true:

- (67)  $x \in \limsup S$  iff for every *n* there exists *k* such that  $x \in S(n+k)$ .
- (68)  $x \in \liminf S$  iff there exists n such that for every k holds  $x \in S(n+k)$ .
- (69) Intersection  $S \subseteq \liminf S$ .
- (70)  $\limsup S \subseteq \bigcup S$ .
- (71)  $\liminf S \subseteq \limsup S.$

Let X be a set, let  $S_1$  be a  $\sigma$ -field of subsets of X, and let S be a sequence of subsets of  $S_1$ . The functor  $S^{\mathbf{c}}$  yields a sequence of subsets of  $S_1$  and is defined by:

(Def. 7)  $S^{\mathbf{c}} = \text{Complement } S.$ 

Next we state a number of propositions:

- (72)  $\liminf S = (\limsup(S^{\mathbf{c}}))^{\mathbf{c}}.$
- (73)  $\limsup S = (\liminf (S^{\mathbf{c}}))^{\mathbf{c}}.$
- (74) If for every n holds  $S_4(n) = S_2(n) \cup S_3(n)$ , then  $\liminf S_2 \cup \liminf S_3 \subseteq \liminf S_4$ .
- (75) If for every n holds  $S_4(n) = S_2(n) \cap S_3(n)$ , then  $\liminf S_4 = \liminf S_2 \cap \liminf S_3$ .
- (76) If for every *n* holds  $S_4(n) = S_2(n) \cup S_3(n)$ , then  $\limsup S_4 = \limsup S_2 \cup \limsup S_3$ .
- (77) If for every *n* holds  $S_4(n) = S_2(n) \cap S_3(n)$ , then  $\limsup S_4 \subseteq \limsup S_2 \cap \limsup S_3$ .
- (78) If S is constant and the value of S = A, then S is convergent and  $\lim S = A$  and  $\lim \inf S = A$  and  $\limsup S = A$ .
- (79) If S is non-decreasing, then  $\limsup S = \bigcup S$ .
- (80) If S is non-decreasing, then  $\liminf S = \bigcup S$ .
- (81) If S is non-decreasing, then S is convergent and  $\lim S = \bigcup S$ .
- (82) If S is non-increasing, then  $\limsup S = \operatorname{Intersection} S$ .
- (83) If S is non-increasing, then  $\liminf S = \operatorname{Intersection} S$ .

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- If S is non-increasing, then S is convergent and  $\lim S = \text{Intersection } S$ . (84)
- (85) If S is monotone, then S is convergent.

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# The Properties of Supercondensed Sets, Subcondensed Sets and Condensed Sets

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**Summary.** We formalized the article "New concepts in the theory of topological space – supercondensed set, subcondensed set, and condensed set" by Yoshinori Isomichi [4]. First we defined supercondensed, subcondensed, and condensed sets and then gradually, defining other attributes such as regular open set or regular closed set, we formalized all the theorems and remarks that one can find in Isomichi's article.

In the last section, the classification of subsets of a topological space is given, depending on the inclusion relation between the interior of the closure and the closure of the interior of a given subset.

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The notation and terminology used in this paper are introduced in the following papers: [10], [11], [1], [6], [8], [9], [7], [12], [2], [3], and [5].

# 1. Preliminaries

In this paper T denotes a topological space and A, B denote subsets of T. Let D be a non trivial set. Note that ADTS(D) is non trivial.

One can check that there exists a topological space which is anti-discrete, non trivial, non empty, and strict.

One can prove the following propositions:

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- (1) Int  $\overline{\operatorname{Int} A} \cap \operatorname{Int} \overline{\operatorname{Int} B} = \operatorname{Int} \overline{\operatorname{Int} (A \cap B)}$ .
- (2) Int  $\overline{A \cup B} = \operatorname{Int} \overline{A} \cup \operatorname{Int} \overline{B}$ .

# 2. Connections between Supercondensed, Condensed, and Subcondensed Sets

Let T be a topological structure and let A be a subset of T. We say that A is supercondensed if and only if:

(Def. 1) Int  $\overline{A} = \text{Int } A$ .

We say that A is subcondensed if and only if:

(Def. 2)  $\overline{\operatorname{Int} A} = \overline{A}$ .

Next we state two propositions:

- (3) If A is closed, then A is supercondensed.
- (4) If A is open, then A is subcondensed.

Let T be a topological space and let A be a subset of T. Let us observe that A is condensed if and only if:

(Def. 3)  $\overline{\operatorname{Int} A} = \overline{A}$  and  $\operatorname{Int} \overline{A} = \operatorname{Int} A$ .

We now state the proposition

(5) A is condensed iff A is subcondensed and supercondensed.

Let T be a topological space. One can verify that every subset of T which is condensed is also subcondensed and supercondensed and every subset of Twhich is subcondensed and supercondensed is also condensed.

Let T be a topological space. Observe that there exists a subset of T which is condensed, subcondensed, and supercondensed.

One can prove the following propositions:

- (6) If A is supercondensed, then  $A^{c}$  is subcondensed.
- (7) If A is subcondensed, then  $A^{c}$  is supercondensed.
- (8) A is supercondensed iff  $\operatorname{Int} \overline{A} \subseteq A$ .
- (9) A is subcondensed iff  $A \subseteq \overline{\operatorname{Int} A}$ .

Let T be a topological space. Note that every subset of T which is subcondensed is also semi-open and every subset of T which is semi-open is also subcondensed.

We now state the proposition

(10) A is condensed iff  $\operatorname{Int} \overline{A} \subseteq A$  and  $A \subseteq \overline{\operatorname{Int} A}$ .

3. Regular Open and Regular Closed Sets

Let T be a topological structure and let A be a subset of T. We introduce A is regular open as a synonym of A is open condensed.

Let T be a topological structure and let A be a subset of T. We introduce A is regular closed as a synonym of A is closed condensed.

The following proposition is true

(11) For every topological space T holds  $\Omega_T$  is regular open and  $\Omega_T$  is regular closed.

Let T be a topological space. Note that  $\Omega_T$  is regular open and regular closed.

We now state the proposition

(12) For every topological space X holds  $\emptyset_X$  is regular open and  $\emptyset_X$  is regular closed.

Let T be a topological space. One can verify that  $\emptyset_T$  is regular open and regular closed.

The following propositions are true:

 $(14)^2$  Int  $\emptyset_T = \emptyset_T$ .

(15) If A is regular open, then  $A^{c}$  is regular closed.

Let T be a topological space. Observe that there exists a subset of T which is regular open and regular closed.

Let T be a topological space and let A be a regular open subset of T. Observe that  $A^{c}$  is regular closed.

One can prove the following proposition

(16) If A is regular closed, then  $A^{c}$  is regular open.

Let T be a topological space and let A be a regular closed subset of T. One can check that  $A^{c}$  is regular open.

Let T be a topological space. Note that every subset of T which is regular open is also open and every subset of T which is regular closed is also closed.

Next we state the proposition

(17) Int  $\overline{A}$  is regular open and  $\overline{\operatorname{Int} A}$  is regular closed.

Let T be a topological space and let A be a subset of T. Observe that  $\operatorname{Int} \overline{A}$  is regular open and  $\overline{\operatorname{Int} A}$  is regular closed.

Next we state two propositions:

(18) A is regular open iff A is supercondensed and open.

(19) A is regular closed iff A is subcondensed and closed.

Let T be a topological space. One can check the following observations:

\* every subset of T which is regular open is also condensed and open,

<sup>&</sup>lt;sup>2</sup>The proposition (13) has been removed.

- \* every subset of T which is condensed and open is also regular open,
- \* every subset of T which is regular closed is also condensed and closed, and

\* every subset of T which is condensed and closed is also regular closed. One can prove the following two propositions:

- (20) A is condensed iff there exists B such that B is regular open and  $B \subseteq A$  and  $A \subseteq \overline{B}$ .
- (21) A is condensed iff there exists B such that B is regular closed and Int  $B \subseteq A$  and  $A \subseteq B$ .

# 4. Boundaries and Borders

Let T be a topological structure and let A be a subset of T. We introduce Bound A as a synonym of Fr A.

Let T be a topological structure and let A be a subset of T. Then  $\operatorname{Fr} A$  can be characterized by the condition:

(Def. 4) Fr  $A = \overline{A} \setminus \text{Int } A$ .

One can prove the following proposition

(22) Fr A is closed.

Let T be a topological space and let A be a subset of T. Observe that  $\operatorname{Fr} A$  is closed.

One can prove the following proposition

(23) A is condensed iff  $\operatorname{Fr} A = \overline{\operatorname{Int} A} \setminus \operatorname{Int} \overline{A}$  and  $\operatorname{Fr} A = \overline{\operatorname{Int} A} \cap \operatorname{Int}(A^c)$ .

Let T be a topological structure and let A be a subset of T. The functor Border A yields a subset of T and is defined by:

(Def. 5) Border A =Int Fr A.

One can prove the following proposition

(24) Border A is regular open and Border  $A = \operatorname{Int} \overline{A} \setminus \overline{\operatorname{Int} A}$  and Border  $A = \operatorname{Int} \overline{A} \cap \operatorname{Int} \overline{A^c}$ .

Let T be a topological space and let A be a subset of T. One can verify that Border A is regular open.

One can prove the following two propositions:

- (25) A is supercondensed iff Int A is regular open and Border A is empty.
- (26) A is subcondensed iff A is regular closed and Border A is empty.

Let T be a topological space and let A be a subset of T. One can verify that Border Border A is empty.

The following proposition is true

(27) A is condensed iff  $\operatorname{Int} A$  is regular open and  $\overline{A}$  is regular closed and Border A is empty.

# 5. Auxiliary Theorems about Intervals

Next we state a number of propositions:

- (28) For every subset A of  $\mathbb{R}^1$  and for every real number a such that  $A = [-\infty, a]$  holds Int  $A = [-\infty, a]$ .
- (29) For every subset A of  $\mathbb{R}^1$  and for every real number a such that  $A = [a, +\infty]$  holds Int  $A = [a, +\infty]$ .
- (30) For every subset A of  $\mathbb{R}^1$  and for all real numbers a, b such that  $A = ]-\infty, a] \cup ]a, b[_{\mathbb{IQ}} \cup [b, +\infty[$  holds  $\overline{A} =$  the carrier of  $\mathbb{R}^1$ .
- (31) For every subset A of  $\mathbb{R}^1$  and for all real numbers a, b such that  $A = ]a, b[_{\mathbb{Q}}]$  holds Int  $A = \emptyset$ .
- (32) For every subset A of  $\mathbb{R}^1$  and for all real numbers a, b such that  $A = ]a, b[_{\mathbb{IQ}} \text{ holds Int } A = \emptyset.$
- (33) For all real numbers a, b holds  $]-\infty, a] \setminus ]-\infty, b[=[b, a].$
- (34) For all real numbers a, b such that a < b holds  $[b, +\infty)$  misses  $]-\infty, a[$ .
- (35) For all real numbers a, b such that  $a \ge b$  holds  $]a, b[_{\mathbb{IQ}} = \emptyset$ .
- (36) For all real numbers a, b holds  $]a, b[_{\mathbb{IQ}} \subseteq [a, +\infty[.$
- (37) For every subset A of  $\mathbb{R}^1$  and for all real numbers a, b, c such that  $A = ]-\infty, a[\cup]b, c[\mathbb{Q} \text{ and } a < b \text{ and } b < c \text{ holds Int } A = ]-\infty, a[.$
- (38) For all real numbers a, b holds [a, b] misses  $]b, +\infty[$ .
- (39) For every real number b holds  $[b, +\infty[ \setminus ]b, +\infty[ = \{b\}.$
- (40) For all real numbers a, b such that a < b holds  $[a, b] = [a, +\infty[ \ ]b, +\infty[.$
- (41) For all real numbers a, b such that a < b holds  $\mathbb{R} = ]-\infty, a[ \cup [a, b] \cup ]b, +\infty[.$
- (42) For all real numbers a, b holds  $]a, b[=]a, +\infty[\setminus [b, +\infty[.$
- (43) For all real numbers a, b, c such that b < c and c < a holds  $]-\infty, a[ \setminus [b, c] = ]-\infty, b[ \cup ]c, a[.$
- (44) For every subset A of  $\mathbb{R}^1$  and for all real numbers a, b, c such that  $A = [-\infty, a] \cup [b, c]$  and a < b and b < c holds Int  $A = [-\infty, a] \cup [b, c]$ .

# 6. Classification of Subsets

Let A, B be sets. We introduce A and B are  $\subseteq$ -incomparable as an antonym of A and B are  $\subseteq$ -comparable.

We now state the proposition

(45) For all sets A, B holds A and B are  $\subseteq$ -incomparable or  $A \subseteq B$  or  $B \subset A$ . Let us consider T, A. We say that A is of the 1<sup>st</sup> class if and only if:

(Def. 6) Int  $\overline{A} \subseteq \overline{\operatorname{Int} A}$ .

We say that A is of the  $2^{nd}$  class if and only if:

(Def. 7)  $\overline{\operatorname{Int} A} \subset \operatorname{Int} \overline{A}$ .

We say that A is of the  $3^{rd}$  class if and only if:

(Def. 8)  $\overline{\operatorname{Int} A}$  and  $\operatorname{Int} \overline{A}$  are  $\subseteq$ -incomparable.

The following proposition is true

- (46) A is of the  $1^{st}$  class, or of the  $2^{nd}$  class, or of the  $3^{rd}$  class.
  - Let T be a topological space. One can verify the following observations:
  - \* every subset of T which is of the 1<sup>st</sup> class is also non of the 2<sup>nd</sup> class and non of the 3<sup>rd</sup> class,
  - \* every subset of T which is of the  $2^{nd}$  class is also non of the  $1^{st}$  class and non of the  $3^{rd}$  class, and
  - \* every subset of T which is of the  $3^{rd}$  class is also non of the  $1^{st}$  class and non of the  $2^{nd}$  class.

One can prove the following proposition

(47) A is of the 1<sup>st</sup> class iff Border A is empty.

Let T be a topological space. Note that every subset of T which is supercondensed is also of the 1<sup>st</sup> class and every subset of T which is subcondensed is also of the 1<sup>st</sup> class.

Let T be a topological space. We say that T has subsets of the 1<sup>st</sup> class if and only if:

(Def. 9) There exists a subset of T which is of the 1<sup>st</sup> class.

We say that T has subsets of the 2<sup>nd</sup> class if and only if:

(Def. 10) There exists a subset of T which is of the 2<sup>nd</sup> class.

We say that T has subsets of the  $3^{rd}$  class if and only if:

(Def. 11) There exists a subset of T which is of the  $3^{rd}$  class.

Let T be an anti-discrete non empty topological space. Note that every subset of T which is proper and non empty is also of the  $2^{nd}$  class.

Let T be an anti-discrete non trivial non empty strict topological space. Observe that there exists a subset of T which is of the  $2^{nd}$  class.

One can verify that there exists a topological space which is non empty, strict, and non trivial and has subsets of the  $1^{st}$  class and subsets of the  $2^{nd}$  class and there exists a topological space which is non empty and strict and has subsets of the  $3^{rd}$  class.

Let us consider T. Observe that there exists a subset of T which is of the  $1^{st}$  class.

Let T be a topological space with subsets of the  $2^{nd}$  class. One can verify that there exists a subset of T which is of the  $2^{nd}$  class.

Let T be a topological space with subsets of the  $3^{rd}$  class. Observe that there exists a subset of T which is of the  $3^{rd}$  class.

The following propositions are true:

- (48) A is of the 1<sup>st</sup> class iff  $A^c$  is of the 1<sup>st</sup> class.
- (49) A is of the  $2^{nd}$  class iff  $A^c$  is of the  $2^{nd}$  class.
- (50) A is of the  $3^{\rm rd}$  class iff  $A^{\rm c}$  is of the  $3^{\rm rd}$  class.

Let us consider T and let A be an of the 1<sup>st</sup> class subset of T. Observe that  $A^{c}$  is of the 1<sup>st</sup> class.

Let T be a topological space with subsets of the  $2^{nd}$  class and let A be an of the  $2^{nd}$  class subset of T. Note that  $A^c$  is of the  $2^{nd}$  class.

Let T be a topological space with subsets of the  $3^{rd}$  class and let A be an of the  $3^{rd}$  class subset of T. Note that  $A^c$  is of the  $3^{rd}$  class.

Next we state four propositions:

- (51) If A is of the 1<sup>st</sup> class, then  $\operatorname{Int} \overline{A} = \operatorname{Int} \overline{\operatorname{Int} A}$  and  $\overline{\operatorname{Int} A} = \operatorname{Int} \overline{A}$ .
- (52) If  $\operatorname{Int} \overline{A} = \operatorname{Int} \overline{\operatorname{Int} A}$  or  $\overline{\operatorname{Int} A} = \overline{\operatorname{Int} \overline{A}}$ , then A is of the 1<sup>st</sup> class.
- (53) Suppose A is of the 1<sup>st</sup> class and B is of the 1<sup>st</sup> class. Then  $\operatorname{Int} \overline{A} \cap \operatorname{Int} \overline{B} = \operatorname{Int} \overline{A \cap B}$  and  $\operatorname{Int} \overline{A} \cup \operatorname{Int} \overline{B} = \operatorname{Int} (A \cup B)$ .
- (54) Suppose A is of the 1<sup>st</sup> class and B is of the 1<sup>st</sup> class. Then  $A \cup B$  is of the 1<sup>st</sup> class and  $A \cap B$  is of the 1<sup>st</sup> class.

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