

The Canonical Formulae

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The articles [23], [29], [11], [28], [14], [2], [27], [12], [30], [8], [5], [3], [20], [9], [6], [22], [7], [10], [1], [4], [15], [17], [18], [24], [25], [19], [16], [21], [13], and [26] provide the notation and terminology for this paper.

1. PRELIMINARIES

One can prove the following propositions:

- (1) For every integer i holds i is even iff $i - 1$ is odd.
- (2) For every integer i holds i is odd iff $i - 1$ is even.
- (3) Let X be a trivial set and x be a set. Suppose $x \in X$. Let f be a function from X into X . Then x is a fixpoint of f .

Let A, B, C be sets. Note that every function from A into C^B is function yielding.

One can prove the following three propositions:

- (4) For every function yielding function f holds $\text{Sub}_f \text{rng } f = \text{rng } f$.
- (5) For all sets A, B, x and for every function f such that $x \in A$ and $f \in B^A$ holds $f(x) \in B$.
- (6) For all sets A, B, C such that if $C = \emptyset$, then $B = \emptyset$ or $A = \emptyset$ and for every function f from A into C^B holds $\text{dom}_\kappa f(\kappa) = A \mapsto B$.

Let us note that \emptyset is function yielding.

In the sequel n is a natural number and p, q, r are elements of HP-WFF.

Next we state the proposition

- (7) For every set x holds $\emptyset(x) = \emptyset$.

Let A be a set and let B be a functional set. One can verify that every function from A into B is function yielding.

One can prove the following propositions:

- (8) For every set X and for every subset A of X holds $[0 \mapsto 1, 1 \mapsto 0] \cdot \chi_{A,X} = \chi_{A^c,X}$.
- (9) For every set X and for every subset A of X holds $[0 \mapsto 1, 1 \mapsto 0] \cdot \chi_{A^c,X} = \chi_{A,X}$.
- (10) For all sets a, b, x, y, x', y' such that $a \neq b$ and $[a \mapsto x, b \mapsto y] = [a \mapsto x', b \mapsto y']$ holds $x = x'$ and $y = y'$.
- (11) For all sets a, b, x, y, X, Y such that $a \neq b$ and $x \in X$ and $y \in Y$ holds $[a \mapsto x, b \mapsto y] \in \prod[a \mapsto X, b \mapsto Y]$.
- (12) For every non empty set D and for every function f from 2 into D there exist elements d_1, d_2 of D such that $f = [0 \mapsto d_1, 1 \mapsto d_2]$.
- (13) For all sets a, b, c, d such that $a \neq b$ holds $[a \mapsto c, b \mapsto d] \cdot [a \mapsto b, b \mapsto a] = [a \mapsto d, b \mapsto c]$.
- (14) For all sets a, b, c, d and for every function f such that $a \neq b$ and $c \in \text{dom } f$ and $d \in \text{dom } f$ holds $f \cdot [a \mapsto c, b \mapsto d] = [a \mapsto f(c), b \mapsto f(d)]$.

2. THE CARTESIAN PRODUCT OF FUNCTIONS AND THE FREGE FUNCTION

Let f, g be one-to-one functions. Note that $[f, g]$ is one-to-one.

We now state a number of propositions:

- (15) Let A, B be non empty sets, C, D be sets, f be a function from C into A , and g be a function from D into B . Then $\pi_1(A \times B) \cdot [f, g] = f \cdot \pi_1(C \times D)$.
- (16) Let A, B be non empty sets, C, D be sets, f be a function from C into A , and g be a function from D into B . Then $\pi_2(A \times B) \cdot [f, g] = g \cdot \pi_2(C \times D)$.
- (17) For every function g holds $\emptyset \leftarrow^{\rho} g = \emptyset$.
- (18) For every function yielding function f and for all functions g, h holds $f \leftarrow^{\rho} g \cdot h = (f \cdot h) \leftarrow^{\rho} (g \cdot h)$.
- (19) Let C be a set, A be a non empty set, f be a function from A into $C^{(\emptyset \text{ qua set})}$, and g be a function from A into \emptyset . Then $\text{rng}(f \leftarrow^{\rho} g) = \{\emptyset\}$.
- (20) Let A, B, C be sets such that if $B = \emptyset$, then $A = \emptyset$. Let f be a function from A into C^B and g be a function from A into B . Then $\text{rng}(f \leftarrow^{\rho} g) \subseteq C$.
- (21) For all sets A, B, C such that if $C = \emptyset$, then $B = \emptyset$ or $A = \emptyset$ and for every function f from A into C^B holds $\text{dom } \text{Frege}(f) = B^A$.
- (22) $\text{Frege}(\emptyset) = \{\emptyset\} \mapsto \emptyset$.
- (23) For all sets A, B, C such that if $C = \emptyset$, then $B = \emptyset$ or $A = \emptyset$ and for every function f from A into C^B holds $\text{rng } \text{Frege}(f) \subseteq C^A$.

- (24) Let A, B, C be sets such that if $C = \emptyset$, then $B = \emptyset$ or $A = \emptyset$. Let f be a function from A into C^B . Then $\text{Frege}(f)$ is a function from B^A into C^A .

3. ABOUT PERMUTATIONS

Let A be a set. Observe that every permutation of A is one-to-one.

The following proposition is true

- (25) For all sets A, B and for every permutation P of A and for every permutation Q of B holds $[P, Q]$ is permutation-like.

Let A, B be non empty sets, let P be a permutation of A , and let Q be a function from B into B . The functor $P \Rightarrow Q$ yielding a function from B^A into B^A is defined as follows:

- (Def. 1) For every function f from A into B holds $(P \Rightarrow Q)(f) = Q \cdot f \cdot P^{-1}$.

Let A, B be non empty sets, let P be a permutation of A , and let Q be a permutation of B . Observe that $P \Rightarrow Q$ is permutation-like.

Next we state three propositions:

- (26) Let A, B be non empty sets, P be a permutation of A , Q be a permutation of B , and f be a function from A into B . Then $(P \Rightarrow Q)^{-1}(f) = Q^{-1} \cdot f \cdot P$.
- (27) For all non empty sets A, B and for every permutation P of A and for every permutation Q of B holds $(P \Rightarrow Q)^{-1} = P^{-1} \Rightarrow Q^{-1}$.
- (28) Let A, B, C be non empty sets, f be a function from A into C^B , g be a function from A into B , P be a permutation of B , and Q be a permutation of C . Then $((P \Rightarrow Q) \cdot f) \leftrightarrow (P \cdot g) = Q \cdot f \leftrightarrow g$.

4. SET VALUATIONS

A SetValuation is a non-empty many sorted set indexed by \mathbb{N} .

In the sequel V denotes a SetValuation.

Let us consider V . The functor $\text{SetVal } V$ yielding a many sorted set indexed by HP-WFF is defined by the conditions (Def. 2).

- (Def. 2)(i) $(\text{SetVal } V)(\text{VERUM}) = 1$,
- (ii) for every n holds $(\text{SetVal } V)(\text{prop } n) = V(n)$, and
- (iii) for all p, q holds $(\text{SetVal } V)(p \wedge q) = [(\text{SetVal } V)(p), (\text{SetVal } V)(q)]$ and $(\text{SetVal } V)(p \Rightarrow q) = (\text{SetVal } V)(q)^{(\text{SetVal } V)(p)}$.

Let us consider V, p . The functor $\text{SetVal}(V, p)$ is defined as follows:

- (Def. 3) $\text{SetVal}(V, p) = (\text{SetVal } V)(p)$.

Let us consider V, p . One can check that $\text{SetVal}(V, p)$ is non empty.

Next we state four propositions:

- (29) $\text{SetVal}(V, \text{VERUM}) = 1$.
- (30) $\text{SetVal}(V, \text{prop } n) = V(n)$.
- (31) $\text{SetVal}(V, p \wedge q) = [\text{SetVal}(V, p), \text{SetVal}(V, q)]$.
- (32) $\text{SetVal}(V, p \Rightarrow q) = (\text{SetVal}(V, q))^{\text{SetVal}(V, p)}$.

Let us consider V, p, q . Observe that $\text{SetVal}(V, p \Rightarrow q)$ is functional.

Let us consider V, p, q, r . Note that every element of $\text{SetVal}(V, p \Rightarrow (q \Rightarrow r))$ is function yielding.

Let us consider V, p, q, r . One can check that there exists a function from $\text{SetVal}(V, p \Rightarrow q)$ into $\text{SetVal}(V, p \Rightarrow r)$ which is function yielding and there exists an element of $\text{SetVal}(V, p \Rightarrow (q \Rightarrow r))$ which is function yielding.

5. PERMUTING SET VALUATIONS

Let us consider V . A function is called a permutation of V if:

(Def. 4) $\text{dom it} = \mathbb{N}$ and for every n holds $\text{it}(n)$ is a permutation of $V(n)$.

In the sequel P is a permutation of V .

Let us consider V, P . The functor $\text{Perm } P$ yielding a many sorted function from $\text{SetVal } V$ into $\text{SetVal } V$ is defined by the conditions (Def. 5).

- (Def. 5)(i) $(\text{Perm } P)(\text{VERUM}) = \text{id}_1$,
- (ii) for every n holds $(\text{Perm } P)(\text{prop } n) = P(n)$, and
- (iii) for all p, q there exists a permutation p' of $\text{SetVal}(V, p)$ and there exists a permutation q' of $\text{SetVal}(V, q)$ such that $p' = (\text{Perm } P)(p)$ and $q' = (\text{Perm } P)(q)$ and $(\text{Perm } P)(p \wedge q) = [p', q']$ and $(\text{Perm } P)(p \Rightarrow q) = p' \Rightarrow q'$.

Let us consider V, P, p . The functor $\text{Perm}(P, p)$ yields a function from $\text{SetVal}(V, p)$ into $\text{SetVal}(V, p)$ and is defined by:

(Def. 6) $\text{Perm}(P, p) = (\text{Perm } P)(p)$.

Next we state four propositions:

- (33) $\text{Perm}(P, \text{VERUM}) = \text{id}_{\text{SetVal}(V, \text{VERUM})}$.
- (34) $\text{Perm}(P, \text{prop } n) = P(n)$.
- (35) $\text{Perm}(P, p \wedge q) = [\text{Perm}(P, p), \text{Perm}(P, q)]$.
- (36) For every permutation p' of $\text{SetVal}(V, p)$ and for every permutation q' of $\text{SetVal}(V, q)$ such that $p' = \text{Perm}(P, p)$ and $q' = \text{Perm}(P, q)$ holds $\text{Perm}(P, p \Rightarrow q) = p' \Rightarrow q'$.

Let us consider V, P, p . One can check that $\text{Perm}(P, p)$ is permutation-like.

We now state four propositions:

- (37) For every function g from $\text{SetVal}(V, p)$ into $\text{SetVal}(V, q)$ holds $(\text{Perm}(P, p \Rightarrow q))(g) = \text{Perm}(P, q) \cdot g \cdot (\text{Perm}(P, p))^{-1}$.
- (38) For every function g from $\text{SetVal}(V, p)$ into $\text{SetVal}(V, q)$ holds $(\text{Perm}(P, p \Rightarrow q))^{-1}(g) = (\text{Perm}(P, q))^{-1} \cdot g \cdot \text{Perm}(P, p)$.
- (39) For all functions f, g from $\text{SetVal}(V, p)$ into $\text{SetVal}(V, q)$ such that $f = (\text{Perm}(P, p \Rightarrow q))(g)$ holds $\text{Perm}(P, q) \cdot g = f \cdot \text{Perm}(P, p)$.
- (40) Let given V, P be a permutation of V , and x be a set. Suppose x is a fixpoint of $\text{Perm}(P, p)$. Let f be a function. If f is a fixpoint of $\text{Perm}(P, p \Rightarrow q)$, then $f(x)$ is a fixpoint of $\text{Perm}(P, q)$.

6. CANONICAL FORMULAE

Let us consider p . We say that p is canonical if and only if:

- (Def. 7) For every V there exists a set x such that for every permutation P of V holds x is a fixpoint of $\text{Perm}(P, p)$.

Let us observe that VERUM is canonical.

Next we state several propositions:

- (41) $p \Rightarrow (q \Rightarrow p)$ is canonical.
- (42) $(p \Rightarrow (q \Rightarrow r)) \Rightarrow ((p \Rightarrow q) \Rightarrow (p \Rightarrow r))$ is canonical.
- (43) $p \wedge q \Rightarrow p$ is canonical.
- (44) $p \wedge q \Rightarrow q$ is canonical.
- (45) $p \Rightarrow (q \Rightarrow p \wedge q)$ is canonical.
- (46) If p is canonical and $p \Rightarrow q$ is canonical, then q is canonical.
- (47) If $p \in \text{HP_TAUT}$, then p is canonical.

Let us observe that there exists an element of HP-WFF which is canonical.

7. PSEUDO-CANONICAL FORMULAE

Let us consider p . We say that p is pseudo-canonical if and only if:

- (Def. 8) For every V and for every permutation P of V holds there exists a set which is a fixpoint of $\text{Perm}(P, p)$.

Let us observe that every element of HP-WFF which is canonical is also pseudo-canonical.

One can prove the following propositions:

- (48) $p \Rightarrow (q \Rightarrow p)$ is pseudo-canonical.
- (49) $(p \Rightarrow (q \Rightarrow r)) \Rightarrow ((p \Rightarrow q) \Rightarrow (p \Rightarrow r))$ is pseudo-canonical.

- (50) $p \wedge q \Rightarrow p$ is pseudo-canonical.
- (51) $p \wedge q \Rightarrow q$ is pseudo-canonical.
- (52) $p \Rightarrow (q \Rightarrow p \wedge q)$ is pseudo-canonical.
- (53) If p is pseudo-canonical and $p \Rightarrow q$ is pseudo-canonical, then q is pseudo-canonical.
- (54) Let given p, q , given V , and P be a permutation of V . Suppose there exists a set which is a fixpoint of $\text{Perm}(P, p)$ and there exists no set which is a fixpoint of $\text{Perm}(P, q)$. Then $p \Rightarrow q$ is not pseudo-canonical.
- (55) $((\text{prop } 0 \Rightarrow \text{prop } 1) \Rightarrow \text{prop } 0) \Rightarrow \text{prop } 0$ is not pseudo-canonical.

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