## Some Isomorphisms Between Functor Categories

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**Summary.** We define some well known isomorphisms between functor categories: between  $A^{\dot{\bigcirc}(o,m)}$  and A, between  $C^{[A,B]}$  and  $(C^B)^A$ , and between  $[B,C]^A$  and  $[B^A, C^A]$ . Compare [12] and [11]. Unfortunately in this paper "functor" is used in two different meanings, as a lingual function and as a functor between categories.

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The notation and terminology used in this paper are introduced in the following papers: [17], [18], [4], [5], [3], [7], [1], [2], [10], [13], [8], [14], [6], [9], [16], and [15].

## 1. Preliminaries

The scheme *ChoiceD* concerns a non-empty set  $\mathcal{A}$ , a non-empty set  $\mathcal{B}$ , and a binary predicate  $\mathcal{P}$ , and states that:

there exists a function h from  $\mathcal{A}$  into  $\mathcal{B}$  such that for every element a of  $\mathcal{A}$  holds  $\mathcal{P}[a, h(a)]$ 

provided the parameters meet the following requirement:

• for every element a of  $\mathcal{A}$  there exists an element b of  $\mathcal{B}$  such that  $\mathcal{P}[a, b]$ .

Let A, B, C be non-empty sets, and let f be a function from A into  $C^B$ . Then uncurry f is a function from [A, B] into C.

We now state several propositions:

(1) For all non-empty sets A, B, C and for every function f from A into  $C^B$  holds curry uncurry f = f.

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- (2) For all non-empty sets A, B, C and for every function f from A into  $C^B$  and for every element a of A and for every element b of B holds (uncurry  $f)(\langle a, b \rangle) = f(a)(b)$ .
- (3) For an arbitrary x and for every non-empty set A and for all functions f, g from  $\{x\}$  into A such that f(x) = g(x) holds f = g.
- (4) For all non-empty sets A, B and for every element x of A and for every function f from A into B holds  $f(x) \in \operatorname{rng} f$ .
- (5) For all non-empty sets A, B, C and for all functions f, g from A into [B, C] such that  $\pi_1(B \times C) \cdot f = \pi_1(B \times C) \cdot g$  and  $\pi_2(B \times C) \cdot f = \pi_2(B \times C) \cdot g$ holds f = g.

We adopt the following rules: A, B, C will be categories and  $F, F_1, F_2$  will be functors from A to B. The following two propositions are true:

- (6) For every morphism f of A holds  $\operatorname{id}_{\operatorname{cod} f} \cdot f = f$ .
- (7) For every morphism f of A holds  $f \cdot id_{\text{dom } f} = f$ .
- In the sequel o, m will be arbitrary. The following two propositions are true:
- (8) o is an object of  $B^A$  if and only if o is a functor from A to B.
- (9) For every morphism f of  $B^A$  there exist functors  $F_1$ ,  $F_2$  from A to B and there exists a natural transformation t from  $F_1$  to  $F_2$  such that  $F_1$  is naturally transformable to  $F_2$  and dom  $f = F_1$  and cod  $f = F_2$  and  $f = \langle \langle F_1, F_2 \rangle, t \rangle$ .

## 2. The isomorphism between $A^{\dot{\bigcirc}(o,m)}$ and A

Let us consider A, B, and let a be an object of A. The functor  $a \mapsto B$  yields a functor from  $B^A$  to B and is defined by:

(Def.1) for all functors  $F_1$ ,  $F_2$  from A to B and for every natural transformation t from  $F_1$  to  $F_2$  such that  $F_1$  is naturally transformable to  $F_2$  holds  $(a \mapsto B)(\langle \langle F_1, F_2 \rangle, t \rangle) = t(a)$ .

One can prove the following two propositions:

- (10) The objects of  $\dot{\heartsuit}(o,m) = \{o\}$  and the morphisms of  $\dot{\circlearrowright}(o,m) = \{m\}$ .
- (11)  $A^{\dot{\circlearrowright}(o,m)} \cong A.$

3. The isomorphism between  $C^{[A,B]}$  and  $(C^B)^A$ 

Next we state four propositions:

- (12) For every functor F from [A, B] to C and for every object a of A and for every object b of B holds  $F(a, -)(b) = F(\langle a, b \rangle)$ .
- (13) For all objects  $a_1, a_2$  of A and for all objects  $b_1, b_2$  of B holds  $\hom(a_1, a_2) \neq \emptyset$  and  $\hom(b_1, b_2) \neq \emptyset$  if and only if  $\hom(\langle a_1, b_1 \rangle, \langle a_2, b_2 \rangle) \neq \emptyset$ .

- (14) Let  $a_1, a_2$  be objects of A. Then for all objects  $b_1, b_2$  of B such that  $\hom(\langle a_1, b_1 \rangle, \langle a_2, b_2 \rangle) \neq \emptyset$  and for every morphism f of A and for every morphism g of B holds  $\langle f, g \rangle$  is a morphism from  $\langle a_1, b_1 \rangle$  to  $\langle a_2, b_2 \rangle$  if and only if f is a morphism from  $a_1$  to  $a_2$  and g is a morphism from  $b_1$  to  $b_2$ .
- (15) For all functors  $F_1$ ,  $F_2$  from [A, B] to C such that  $F_1$  is naturally transformable to  $F_2$  and for every natural transformation t from  $F_1$  to  $F_2$  and for every object a of A holds  $F_1(a, -)$  is naturally transformable to  $F_2(a, -)$  and  $(\operatorname{curry} t)(a)$  is a natural transformation from  $F_1(a, -)$  to  $F_2(a, -)$ .

Let us consider A, B, C, and let F be a functor from [A, B] to C, and let f be a morphism of A. The functor  $\operatorname{curry}(F, f)$  yields a function from the morphisms of B into the morphisms of C and is defined by:

(Def.2)  $\operatorname{curry}(F, f) = (\operatorname{curry} F)(f).$ 

The following two propositions are true:

- (16) For all objects  $a_1$ ,  $a_2$  of A and for all objects  $b_1$ ,  $b_2$  of B and for every morphism f of A and for every morphism g of B such that  $f \in \text{hom}(a_1, a_2)$  and  $g \in \text{hom}(b_1, b_2)$  holds  $\langle f, g \rangle \in \text{hom}(\langle a_1, b_1 \rangle, \langle a_2, b_2 \rangle)$ .
- (17) For every functor F from [A, B] to C and for all objects a, b of A such that  $hom(a, b) \neq \emptyset$  and for every morphism f from a to b holds F(a, -) is naturally transformable to F(b, -) and  $curry(F, f) \cdot$  the id-map of B is a natural transformation from F(a, -) to F(b, -).

Let us consider A, B, C, and let F be a functor from [A, B] to C, and let f be a morphism of A. The functor F(f, -) yielding a natural transformation from F(dom f, -) to F(cod f, -) is defined by:

(Def.3)  $F(f, -) = \operatorname{curry}(F, f) \cdot \operatorname{the id-map} of B.$ 

We now state four propositions:

- (18) For every functor F from [A, B] to C and for every morphism g of A holds  $F(\operatorname{dom} g, -)$  is naturally transformable to  $F(\operatorname{cod} g, -)$ .
- (19) For every functor F from [A, B] to C and for every morphism f of A and for every object b of B holds  $F(f, -)(b) = F(\langle f, id_b \rangle)$ .
- (20) For every functor F from [A, B] to C and for every object a of A holds  $\operatorname{id}_{F(a,-)} = F(\operatorname{id}_a, -).$
- (21) For every functor F from [A, B] to C and for all morphisms g, f of A such that dom  $g = \operatorname{cod} f$  and for every natural transformation t from  $F(\operatorname{dom} f, -)$  to  $F(\operatorname{dom} g, -)$  such that t = F(f, -) holds  $F(g \cdot f, -) = F(g, -) \circ t$ .

Let us consider A, B, C, and let F be a functor from [A, B] to C. The functor export(F) yielding a functor from A to  $C^B$  is defined as follows:

(Def.4) for every morphism f of A holds  $(export(F))(f) = \langle \langle F(\operatorname{dom} f, -), F(\operatorname{cod} f, -) \rangle, F(f, -) \rangle$ .

We now state several propositions:

- (22) For every functor F from [A, B] to C and for every morphism f of A holds  $(export(F))(f) = \langle \langle F(\operatorname{dom} f, -), F(\operatorname{cod} f, -) \rangle, F(f, -) \rangle$ .
- (23) For all functors  $F_1$ ,  $F_2$  from A to B such that  $F_1$  is transformable to  $F_2$ and for every transformation t from  $F_1$  to  $F_2$  and for every object a of Aholds  $t(a) \in \text{hom}(F_1(a), F_2(a))$ .
- (24) For every functor F from [A, B] to C and for every object a of A holds (export(F))(a) = F(a, -).
- (25) For every functor F from [A, B] to C and for every object a of A holds (export(F))(a) is a functor from B to C.
- (26) For all functors  $F_1$ ,  $F_2$  from [A, B] to C such that  $export(F_1) = export(F_2)$  holds  $F_1 = F_2$ .
- (27) Let  $F_1$ ,  $F_2$  be functors from [A, B] to C. Suppose  $F_1$  is naturally transformable to  $F_2$ . Let t be a natural transformation from  $F_1$  to  $F_2$ . Then export $(F_1)$  is naturally transformable to export $(F_2)$  and there exists a natural transformation G from export $(F_1)$  to export $(F_2)$  such that for every function s from [ the objects of A, the objects of B ] into the morphisms of C such that t = s and for every object a of A holds  $G(a) = \langle \langle (export(F_1))(a), (export(F_2))(a) \rangle$ ,  $(eurry s)(a) \rangle$ .

Let us consider A, B, C, and let  $F_1$ ,  $F_2$  be functors from [A, B] to C satisfying the condition:  $F_1$  is naturally transformable to  $F_2$ . Let t be a natural transformation from  $F_1$  to  $F_2$ . The functor export(t) yielding a natural transformation from export $(F_1)$  to export $(F_2)$  is defined as follows:

(Def.5) for every function s from [the objects of A, the objects of B] into the morphisms of C such that t = s and for every object a of A holds  $(export(t))(a) = \langle \langle (export(F_1))(a), (export(F_2))(a) \rangle, (curry s)(a) \rangle$ .

We now state several propositions:

- (28) For every functor F from [A, B] to C holds  $id_{export(F)} = export(id_F)$ .
- (29) For all functors  $F_1$ ,  $F_2$ ,  $F_3$  from [A, B] to C such that  $F_1$  is naturally transformable to  $F_2$  and  $F_2$  is naturally transformable to  $F_3$  and for every natural transformation  $t_1$  from  $F_1$  to  $F_2$  and for every natural transformation  $t_2$  from  $F_2$  to  $F_3$  holds export $(t_2 \circ t_1) = \text{export}(t_2) \circ \text{export}(t_1)$ .
- (30) For all functors  $F_1$ ,  $F_2$  from [A, B] to C such that  $F_1$  is naturally transformable to  $F_2$  and for all natural transformations  $t_1$ ,  $t_2$  from  $F_1$  to  $F_2$  such that  $export(t_1) = export(t_2)$  holds  $t_1 = t_2$ .
- (31) For every functor G from A to  $C^B$  there exists a functor F from [A, B] to C such that  $G = \operatorname{export}(F)$ .
- (32) For all functors  $F_1$ ,  $F_2$  from [A, B] to C such that  $export(F_1)$  is naturally transformable to  $export(F_2)$  and for every natural transformation t from  $export(F_1)$  to  $export(F_2)$  holds  $F_1$  is naturally transformable to  $F_2$  and there exists a natural transformation u from  $F_1$  to  $F_2$  such that t = export(u).

Let us consider A, B, C. The functor  $export_{A,B,C}$  yields a functor from  $C^{[A,B]}$  to  $(C^B)^A$  and is defined by:

(Def.6) for all functors  $F_1$ ,  $F_2$  from [A, B] to C such that  $F_1$  is naturally transformable to  $F_2$  and for every natural transformation t from  $F_1$  to  $F_2$  holds  $export_{A,B,C}(\langle\langle F_1, F_2 \rangle, t \rangle) = \langle\langle export(F_1), export(F_2) \rangle, export(t) \rangle.$ 

Next we state two propositions:

- (33) **export**<sub>A,B,C</sub> is an isomorphism.
- $(34) C^{[A,B]} \cong (C^B)^A.$ 
  - 4. The isomorphism between  $[B, C]^A$  and  $[B^A, C^A]$

We now state the proposition

(35) For all functors  $F_1$ ,  $F_2$  from A to B and for every functor G from B to C such that  $F_1$  is naturally transformable to  $F_2$  and for every natural transformation t from  $F_1$  to  $F_2$  holds  $G \cdot t = G \cdot t$  qua a function.

We now define two new functors. Let us consider A, B. Then  $\pi_1(A \times B)$  is a functor from [A, B] to A. Then  $\pi_2(A \times B)$  is a functor from [A, B] to B. Let us consider A, B, C, and let F be a functor from A to B, and let G be a functor from A to C. Then  $\langle F, G \rangle$  is a functor from A to [B, C]. Let F be a functor from A to [B, C]. The functor  $\pi_1 \cdot F$  yielding a functor from A to B is defined as follows:

(Def.7) 
$$\pi_1 \cdot F = \pi_1(B \times C) \cdot F.$$

The functor  $\pi_2 \cdot F$  yielding a functor from A to C is defined by:

(Def.8)  $\pi_2 \cdot F = \pi_2(B \times C) \cdot F.$ 

The following two propositions are true:

- (36) For every functor F from A to B and for every functor G from A to C holds  $\pi_1 \cdot \langle F, G \rangle = F$  and  $\pi_2 \cdot \langle F, G \rangle = G$ .
- (37) For all functors F, G from A to [B, C] such that  $\pi_1 \cdot F = \pi_1 \cdot G$  and  $\pi_2 \cdot F = \pi_2 \cdot G$  holds F = G.

We now define two new functors. Let us consider A, B, C, and let  $F_1, F_2$  be functors from A to [B, C], and let t be a natural transformation from  $F_1$  to  $F_2$ . The functor  $\pi_1 \cdot t$  yielding a natural transformation from  $\pi_1 \cdot F_1$  to  $\pi_1 \cdot F_2$  is defined as follows:

(Def.9)  $\pi_1 \cdot t = \pi_1(B \times C) \cdot t.$ 

The functor  $\pi_2 \cdot t$  yielding a natural transformation from  $\pi_2 \cdot F_1$  to  $\pi_2 \cdot F_2$  is defined as follows:

(Def.10)  $\pi_2 \cdot t = \pi_2(B \times C) \cdot t.$ 

We now state several propositions:

- (38) For all functors F, G from A to [B, C] such that F is naturally transformable to G holds  $\pi_1 \cdot F$  is naturally transformable to  $\pi_1 \cdot G$  and  $\pi_2 \cdot F$  is naturally transformable to  $\pi_2 \cdot G$ .
- (39) For all functors  $F_1$ ,  $F_2$ ,  $G_1$ ,  $G_2$  from A to [B, C] such that  $F_1$  is naturally transformable to  $F_2$  and  $G_1$  is naturally transformable to  $G_2$ and for every natural transformation s from  $F_1$  to  $F_2$  and for every natural transformation t from  $G_1$  to  $G_2$  such that  $\pi_1 \cdot s = \pi_1 \cdot t$  and  $\pi_2 \cdot s = \pi_2 \cdot t$ holds s = t.
- (40) For every functor F from A to [B, C] holds  $\operatorname{id}_{\pi_1 \cdot F} = \pi_1 \cdot (\operatorname{id}_F)$  and  $\operatorname{id}_{\pi_2 \cdot F} = \pi_2 \cdot (\operatorname{id}_F)$ .
- (41) For all functors F, G, H from A to [B, C] such that F is naturally transformable to G and G is naturally transformable to H and for every natural transformation s from F to G and for every natural transformation t from G to H holds  $\pi_1 \cdot (t^\circ s) = \pi_1 \cdot t^\circ \pi_1 \cdot s$  and  $\pi_2 \cdot (t^\circ s) = \pi_2 \cdot t^\circ \pi_2 \cdot s$ .
- (42) For every functor F from A to B and for every functor G from A to C and for all objects a, b of A such that  $hom(a, b) \neq \emptyset$  and for every morphism f from a to b holds  $\langle F, G \rangle(f) = \langle F(f), G(f) \rangle$ .
- (43) For every functor F from A to B and for every functor G from A to C and for every object a of A holds  $\langle F, G \rangle(a) = \langle F(a), G(a) \rangle$ .
- (44) For all functors  $F_1$ ,  $G_1$  from A to B and for all functors  $F_2$ ,  $G_2$  from A to C such that  $F_1$  is transformable to  $G_1$  and  $F_2$  is transformable to  $G_2$  holds  $\langle F_1, F_2 \rangle$  is transformable to  $\langle G_1, G_2 \rangle$ .

Let us consider A, B, C, and let  $F_1$ ,  $G_1$  be functors from A to B, and let  $F_2$ ,  $G_2$  be functors from A to C satisfying the condition:  $F_1$  is transformable to  $G_1$  and  $F_2$  is transformable to  $G_2$ . Let  $t_1$  be a transformation from  $F_1$  to  $G_1$ , and let  $t_2$  be a transformation from  $F_2$  to  $G_2$ . The functor  $\langle t_1, t_2 \rangle$  yielding a transformation from  $\langle F_1, F_2 \rangle$  to  $\langle G_1, G_2 \rangle$  is defined as follows:

(Def.11) 
$$\langle t_1, t_2 \rangle = \langle t_1, t_2 \rangle.$$

One can prove the following propositions:

- (45) For all functors  $F_1$ ,  $G_1$  from A to B and for all functors  $F_2$ ,  $G_2$  from A to C such that  $F_1$  is transformable to  $G_1$  and  $F_2$  is transformable to  $G_2$  and for every transformation  $t_1$  from  $F_1$  to  $G_1$  and for every transformation  $t_2$  from  $F_2$  to  $G_2$  and for every object a of A holds  $\langle t_1, t_2 \rangle \langle a \rangle = \langle t_1(a), t_2(a) \rangle$ .
- (46) For all functors  $F_1$ ,  $G_1$  from A to B and for all functors  $F_2$ ,  $G_2$  from A to C such that  $F_1$  is naturally transformable to  $G_1$  and  $F_2$  is naturally transformable to  $G_2$  holds  $\langle F_1, F_2 \rangle$  is naturally transformable to  $\langle G_1, G_2 \rangle$ .

Let us consider A, B, C, and let  $F_1$ ,  $G_1$  be functors from A to B, and let  $F_2$ ,  $G_2$  be functors from A to C satisfying the conditions:  $F_1$  is naturally transformable to  $G_1$  and  $F_2$  is naturally transformable to  $G_2$ . Let  $t_1$  be a natural transformation from  $F_1$  to  $G_1$ , and let  $t_2$  be a natural transformation from  $F_2$ to  $G_2$ . The functor  $\langle t_1, t_2 \rangle$  yielding a natural transformation from  $\langle F_1, F_2 \rangle$  to  $\langle G_1, G_2 \rangle$  is defined as follows: (Def.12)  $\langle t_1, t_2 \rangle = \langle t_1, t_2 \rangle.$ 

Next we state the proposition

(47) For all functors  $F_1$ ,  $G_1$  from A to B and for all functors  $F_2$ ,  $G_2$  from A to C such that  $F_1$  is naturally transformable to  $G_1$  and  $F_2$  is naturally transformable to  $G_2$  and for every natural transformation  $t_1$  from  $F_1$  to  $G_1$  and for every natural transformation  $t_2$  from  $F_2$  to  $G_2$  holds  $\pi_1 \langle t_1, t_2 \rangle = t_1$  and  $\pi_2 \cdot \langle t_1, t_2 \rangle = t_2$ .

Let us consider A, B, C. The functor **distribute**<sub>A,B,C</sub> yielding a functor from  $[B, C]^A$  to  $[B^A, C^A]$  is defined by:

(Def.13) for all functors  $F_1$ ,  $F_2$  from A to [B, C] such that  $F_1$  is naturally transformable to  $F_2$  and for every natural transformation t from  $F_1$  to  $F_2$ holds **distribute**<sub>A,B,C</sub>( $\langle\langle F_1, F_2 \rangle, t \rangle$ ) =  $\langle\langle\langle \pi_1 \cdot F_1, \pi_1 \cdot F_2 \rangle, \pi_1 \cdot t \rangle, \langle\langle \pi_2 \cdot F_1, \pi_2 \cdot F_2 \rangle, \pi_2 \cdot t \rangle$ .

One can prove the following two propositions:

- (48) **distribute**<sub>A,B,C</sub> is an isomorphism.
- $(49) \quad [B, C]^A \cong [B^A, C^A].$

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