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Binary Arithmetics

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Summary. Formalizes the basic concepts of binary arithmetic and its related operations. We present the definitions for the following logical operators: 'or' and 'xor' (exclusive or) and include in this article some theorems concerning these operators. We also introduce the concept of an n-bit register. Such registers are used in the definition of binary unsigned arithmetic presented in this article. Theorems on the relationships of such concepts to the operations of natural numbers are also given.

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The articles [14], [9], [17], [2], [3], [15], [1], [19], [18], [7], [8], [6], [4], [13], [12], [10], [5], [11], and [16] provide the notation and terminology for this paper.

Let us note that there exists a natural number which is non empty.

Next we state four propositions:

(1) For all natural numbers i, j holds $+_{\mathbb{N}}(i, j) = i + j$.

(2) Let i, n be natural numbers, D be a non empty set, d be an element of D, and z be a n-tuple of D. If $i \in \operatorname{Seg} n$, then $(z^\frown \langle d \rangle)_i = z_i$.

(3) Let n be a natural number, D be a non empty set, d be an element of D, and z be a n-tuple of D. Then $(z^\frown \langle d \rangle)_{n+1} = d$.

(5)[1] For all natural numbers i, n such that $i \in \operatorname{Seg} n$ holds i is non empty.

Let x, y be boolean sets. The functor $x \vee y$ is defined as follows:

(Def. 1) $x \vee y = \neg(\neg x \wedge \neg y)$.

Let us note that the functor $x \vee y$ is commutative.

Let x, y be boolean sets. The functor $x \oplus y$ is defined by:

(Def. 2) $x \oplus y = \neg x \wedge y \vee x \wedge \neg y$.

Let us note that the functor $x \oplus y$ is commutative.

Let x, y be boolean sets. One can check that $x \vee y$ is boolean.

Let x, y be boolean sets. One can verify that $x \oplus y$ is boolean.

Let x, y be elements of *Boolean*. Then $x \vee y$ is an element of *Boolean*. Then $x \oplus y$ is an element of *Boolean*.

In the sequel x, y, z are boolean sets.

Next we state a number of propositions:

[1] The proposition (4) has been removed.

(7)[2] $x \vee \mathit{false} = x$.

(9)[3] $\neg(x \wedge y) = \neg x \vee \neg y$.

(10) $\neg(x \vee y) = \neg x \wedge \neg y$.

(12)[4] $x \wedge y = \neg(\neg x \vee \neg y)$.

(13) $\mathit{true} \oplus x = \neg x$.

(14) $\mathit{false} \oplus x = x$.

(15) $x \oplus x = \mathit{false}$.

(16) $x \wedge x = x$.

(17) $x \oplus \neg x = \mathit{true}$.

(18) $x \vee \neg x = \mathit{true}$.

(19) $x \vee \mathit{true} = \mathit{true}$.

(20) $(x \vee y) \vee z = x \vee (y \vee z)$.

(21) $x \vee x = x$.

(22) $x \wedge (y \vee z) = x \wedge y \vee x \wedge z$.

(23) $x \vee y \wedge z = (x \vee y) \wedge (x \vee z)$.

(24) $x \vee x \wedge y = x$.

(25) $x \wedge (x \vee y) = x$.

(26) $x \vee \neg x \wedge y = x \vee y$.

(27) $x \wedge (\neg x \vee y) = x \wedge y$.

(33)[5] $\mathit{true} \oplus \mathit{false} = \mathit{true}$.

(34) $(x \oplus y) \oplus z = x \oplus (y \oplus z)$.

(35) $x \oplus \neg x \wedge y = x \vee y$.

(36) $x \vee (x \oplus y) = x \vee y$.

(37) $x \vee (\neg x \oplus y) = x \vee \neg y$.

(38) $x \wedge (y \oplus z) = x \wedge y \oplus x \wedge z$.

Let i, j be natural numbers. The functor $i -' j$ yields a natural number and is defined as follows:

(Def. 3) $i -' j = \begin{cases} i - j, \text{ if } i - j \geq 0, \\ 0, \text{ otherwise.} \end{cases}$

The following proposition is true

(39) For all natural numbers i, j holds $(i + j) -' j = i$.

We adopt the following convention: i is a natural number, n is a non empty natural number, and x, y, z_1, z_2 are n-tuples of *Boolean*.

Let n be a natural number and let x be a n-tuple of *Boolean*. The functor $\neg x$ yields a n-tuple of *Boolean* and is defined as follows:

[2] The proposition (6) has been removed.
[3] The proposition (8) has been removed.
[4] The proposition (11) has been removed.
[5] The propositions (28)–(32) have been removed.

(Def. 4) For every i such that $i \in \operatorname{Seg} n$ holds $(\neg x)_i = \neg(x_i)$.

Let n be a non empty natural number and let x, y be n-tuples of *Boolean*. The functor $\operatorname{carry}(x,y)$ yielding a n-tuple of *Boolean* is defined as follows:

(Def. 5) $(\operatorname{carry}(x,y))_1 = \mathit{false}$ and for every i such that $1 \leq i$ and $i < n$ holds $(\operatorname{carry}(x,y))_{i+1} = x_i \wedge y_i \vee x_i \wedge (\operatorname{carry}(x,y))_i \vee y_i \wedge (\operatorname{carry}(x,y))_i$.

Let n be a natural number and let x be a n-tuple of *Boolean*. The functor $\operatorname{Binary}(x)$ yields a n-tuple of $\mathbb{N}$ and is defined as follows:

(Def. 6) For every i such that $i \in \operatorname{Seg} n$ holds $(\operatorname{Binary}(x))_i = (x_i = \mathit{false} \to 0, 2^{i-'1})$.

Let n be a natural number and let x be a n-tuple of *Boolean*. The functor $\operatorname{Absval}(x)$ yields a natural number and is defined as follows:

(Def. 7) $\operatorname{Absval}(x) = +_{\mathbb{N}} \circledast \operatorname{Binary}(x)$.

Let us consider n, x, y. The functor $x+y$ yielding a n-tuple of *Boolean* is defined by:

(Def. 8) For every i such that $i \in \operatorname{Seg} n$ holds $(x+y)_i = x_i \oplus y_i \oplus (\operatorname{carry}(x,y))_i$.

Let us consider n, z_1, z_2. The functor $\operatorname{add\_ovfl}(z_1,z_2)$ yielding an element of *Boolean* is defined as follows:

(Def. 9) $\operatorname{add\_ovfl}(z_1,z_2) = (z_1)_n \wedge (z_2)_n \vee (z_1)_n \wedge (\operatorname{carry}(z_1,z_2))_n \vee (z_2)_n \wedge (\operatorname{carry}(z_1,z_2))_n$.

The scheme *Ind from 1* concerns a unary predicate $\mathcal{P}$, and states that:

For every non empty natural number k holds $\mathcal{P}[k]$

provided the parameters meet the following requirements:

- $\mathcal{P}[1]$, and
- For every non empty natural number k such that $\mathcal{P}[k]$ holds $\mathcal{P}[k+1]$.

Let us consider n, z_1, z_2. We say that z_1 and z_2 are summable if and only if:

(Def. 10) $\operatorname{add\_ovfl}(z_1,z_2) = \mathit{false}$.

We now state three propositions:

(40) For every 1-tuple z_1 of *Boolean* holds $z_1 = \langle \mathit{false} \rangle$ or $z_1 = \langle \mathit{true} \rangle$.

(41) For every 1-tuple z_1 of *Boolean* such that $z_1 = \langle \mathit{false} \rangle$ holds $\operatorname{Absval}(z_1) = 0$.

(42) For every 1-tuple z_1 of *Boolean* such that $z_1 = \langle \mathit{true} \rangle$ holds $\operatorname{Absval}(z_1) = 1$.

Let n_1 be a non empty natural number, let n_2 be a natural number, let D be a non empty set, let z_1 be a n_1-tuple of D, and let z_2 be a n_2-tuple of D. Then $z_1 \frown z_2$ is a n_1+n_2-tuple of D.

Let D be a non empty set and let d be an element of D. Then $\langle d \rangle$ is a 1-tuple of D.

We now state several propositions:

(43) Let z_1, z_2 be n-tuples of *Boolean*, d_1, d_2 be elements of *Boolean*, and i be a natural number. If $i \in \operatorname{Seg} n$, then $(\operatorname{carry}(z_1 \frown \langle d_1 \rangle, z_2 \frown \langle d_2 \rangle))_i = (\operatorname{carry}(z_1,z_2))_i$.

(44) For all n-tuples z_1, z_2 of *Boolean* and for all elements d_1, d_2 of *Boolean* holds $\operatorname{add\_ovfl}(z_1,z_2) = (\operatorname{carry}(z_1 \frown \langle d_1 \rangle, z_2 \frown \langle d_2 \rangle))_{n+1}$.

(45) For all n-tuples z_1, z_2 of *Boolean* and for all elements d_1, d_2 of *Boolean* holds $z_1 \frown \langle d_1 \rangle + z_2 \frown \langle d_2 \rangle = (z_1+z_2) \frown \langle d_1 \oplus d_2 \oplus \operatorname{add\_ovfl}(z_1,z_2) \rangle$.

(46) For every n-tuple z of *Boolean* and for every element d of *Boolean* holds $\operatorname{Absval}(z \frown \langle d \rangle) = \operatorname{Absval}(z) + (d = \mathit{false} \to 0, 2^n)$.

(47) For every n and for all n-tuples z_1, z_2 of *Boolean* holds $\operatorname{Absval}(z_1+z_2) + (\operatorname{add\_ovfl}(z_1,z_2) = \mathit{false} \to 0, 2^n) = \operatorname{Absval}(z_1) + \operatorname{Absval}(z_2)$.

(48) For all n-tuples z_1, z_2 of *Boolean* such that z_1 and z_2 are summable holds $\operatorname{Absval}(z_1+z_2) = \operatorname{Absval}(z_1) + \operatorname{Absval}(z_2)$.

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